

THE CARTOON GUIDE TO

COMPUTER SCIENCE



by Larry Gonick

Also available by Larry Gonick, with Mark Wheelis

THE CARTOON GUIDE TO GENETICS

THE CARTOON GUIDE TO
**COMPUTER
SCIENCE**

Larry Gonick



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LARRY GONICK, THE OVEREDUCATED CARTOONIST, HOLDS TWO DEGREES IN MATHEMATICS FROM HARVARD. HE HAS WORKED AS A FORTRAN PROGRAMMER, AND SOME OF HIS BEST FRIENDS ARE IN THE COMPUTER BUSINESS. HE LIVES IN SAN FRANCISCO WITH HIS WIFE AND DAUGHTER, WHO WOULD LIKE TO FIND SOME CARTOON PROCESSING SOFTWARE TO IMPROVE HIS PRODUCTIVITY.



COMPUTER SCIENCE

LEARNING HAS NEVER BEEN SO EASY OR SO MUCH FUN

Here are the elements of computer science illustrated, simplified, and humor-coated so that you understand them at once. Use this book to lighten up that serious course you are taking or to penetrate the fog of that equally serious textbook you are trying to follow. Read it to gain both an overview and an inner view of that computer you are learning to use. Or if you feel the computer revolution is passing you by, let it give you a point of entry. It won't make a programmer out of you, but it will put you well on the way to computer literacy.

In these pages you'll meet Charles Babbage and his analytical engine, which was never built, and Ada Augusta, Lady Lovelace, who programmed it nevertheless. You'll also meet George Boole, whose algebra underlies the design of circuitry. You'll learn about binary numbers, computer components and architecture, software, programming languages from machine language to BASIC, and special computer applications—cryptography, artificial intelligence, and others you may not have heard of.

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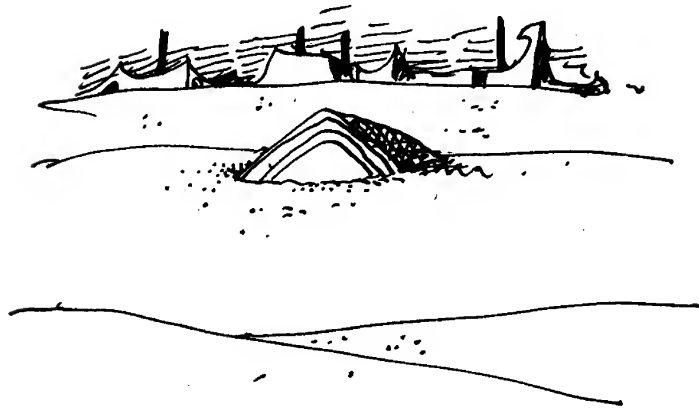
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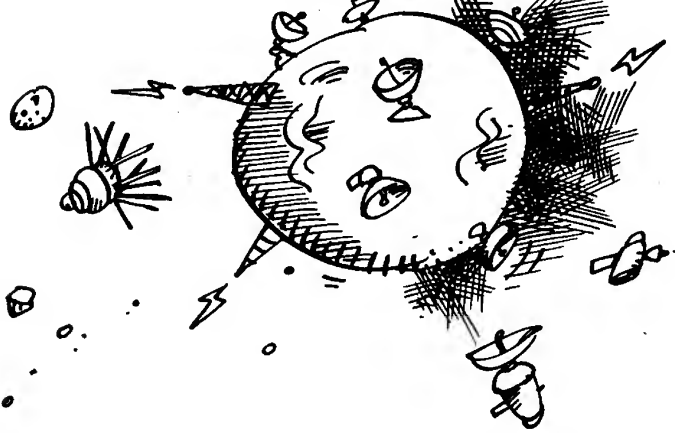
CONTENTS

PART I. THE AGES OF INFORMATION	1
WHAT IS INFORMATION?	7
THE EVOLUTION OF THE COMPUTER	14
PART II. LOGICAL SPAGHETTI	87
THE INFORMATION PROCESSOR	90
THE PROCESSING UNIT	97
MEMORY	151
GETTING EVERYTHING UNDER CONTROL	169
PART III. SOFTWARE	185
TURING MACHINES	190
ALGORITHMS	195
BASIC B.A.S.I.C.	207
SOFTWARE SURVEY	221
IN CONCLUSION	237
BIBLIOGRAPHY	242
INDEX	243

PART I
THE AGES OF
INFORMATION



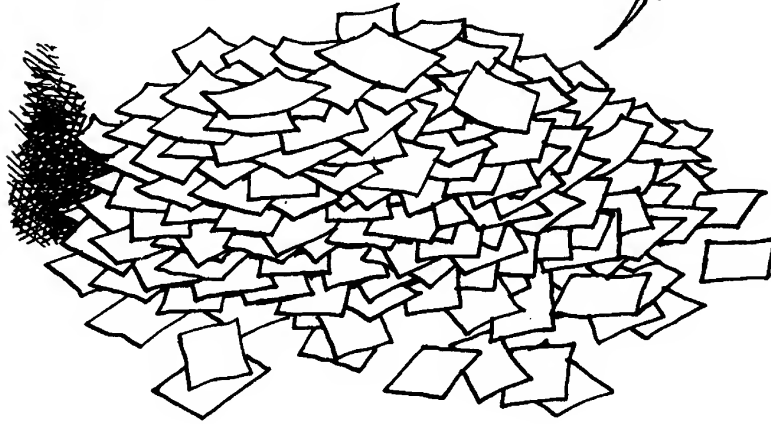
WE LIVE IN
THE AGE OF
EXCESS INFORMATION.
THANKS TO THE
TECHNOLOGICAL
MIRACLES OF THE
TWENTIETH CENTURY,
WE CITIZENS OF
EARTH ENJOY
INSTANT ACCESS
TO MORE
INFORMATION THAN
ANYONE CAN
POSSIBLY
MANAGE!





CLEARLY, THE AGE DEMANDS A PIECE
OF TECHNOLOGY SOLELY DEVOTED
TO STORING, CLASSIFYING,
SORTING, COMPARING, COMBINING,
AND DISPLAYING INFORMATION
AT HIGH SPEED!

THAT, AND
A SHOVEL...

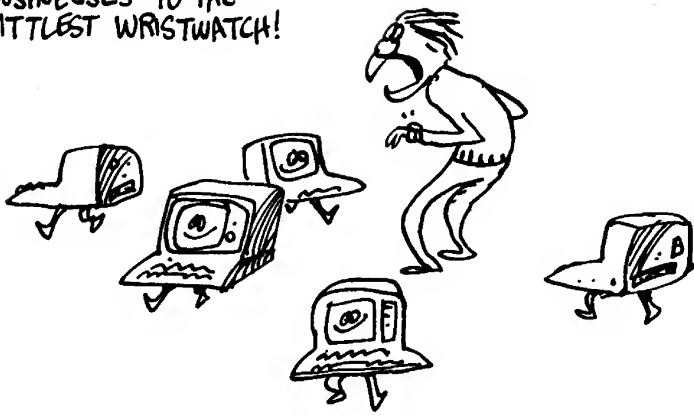


THAT PIECE OF EQUIPMENT IS THE **COMPUTER**.



THIS EXPLAINS WHY COMPUTERS ARE POPPING UP WHEREVER INFORMATION COUNTS, FROM THE BIGGEST BUSINESSES TO THE LITTLEST WRISTWATCH!

IT'S ENOUGH TO MAKE YOU PARANOID!



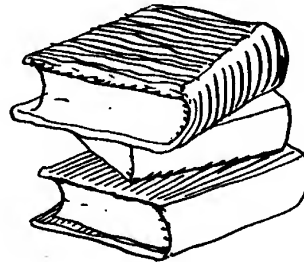
AND IT ALSO ACCOUNTS FOR THE FACT THAT BEFORE YOU CAN UNDERSTAND COMPUTERS, IT HELPS TO KNOW SOMETHING ABOUT INFORMATION FIRST—SUCH AS, FOR EXAMPLE, WHAT IT IS...

WHAT IS IT?
IT'S...
IT'S... AH...
UM... ER...
WHAT A STUPID QUESTION!



What is information?

IN THE EVERYDAY SENSE OF THE WORD, "INFORMATION" MEANS FACTS: THE SORT OF STUFF THAT FILLS NON FICTION BOOKS, AND CAN ONLY BE EXPRESSED IN WORDS.



IN THE WORLD OF COMPUTERS, HOWEVER, THE TERM HAS A MUCH BROADER MEANING.



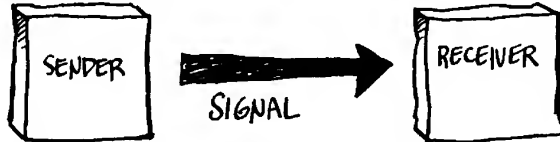
THE MODERN DEFINITION COMES FROM **CLAUDE SHANNON**, A BELL LABS ENGINEER, AMATEUR UNICYCLIST, AND FOUNDER OF THE SCIENCE OF **INFORMATION THEORY**.

MEEP!



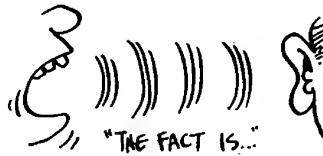
SHANNON ALSO BUILT AN ELECTRIC "MOUSE" THAT COULD BE PROGRAMMED TO RUN MAZES!

ACCORDING TO SHANNON, INFORMATION IS PRESENT WHENEVER A SIGNAL IS TRANSMITTED FROM ONE PLACE TO ANOTHER.

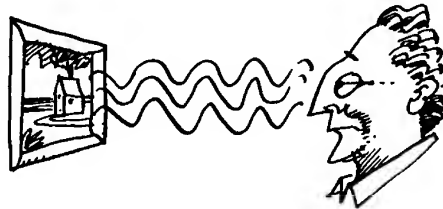


IT DOESN'T MATTER WHAT KIND OF SIGNAL IT IS.
FOR EXAMPLE:

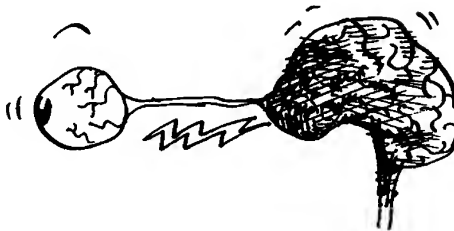
THE SIGNAL MAY BE
IN THE FORM OF
WORDS, THE MOST
FAMILIAR KIND
OF INFORMATION...



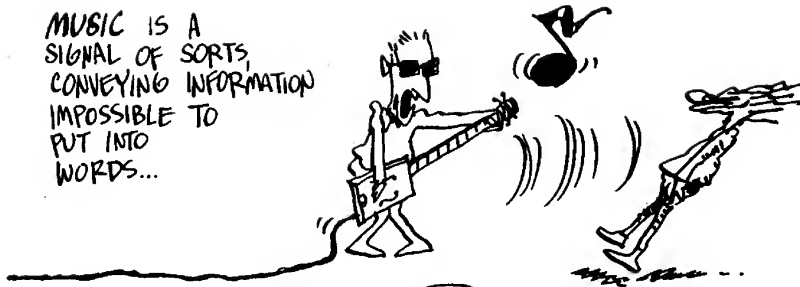
...BUT A PICTURE
ALSO SENDS A
SIGNAL, IN THE
FORM OF LIGHT
WAVES, TO OUR EYES.
IT LOOKS AS IF PICTURES
CONVEY INFORMATION!



FURTHERMORE, OUR
EYE SENDS A
PATTERN OF ELECTRIC
IMPULSES UP THE
OPTIC NERVE TO THE
BRAIN. THAT SIGNAL
CARRIES INFORMATION,
TOO!!



MUSIC IS A
SIGNAL OF SORTS,
CONVEYING INFORMATION
IMPOSSIBLE TO
PUT INTO
WORDS...



FOR THAT MATTER,
A PUNCH IN THE
MOUTH IS NOT
WITHOUT ITS
INFORMATION VALUE!



SO YOU SEE...
INFORMATION COMES
IN MANY FORMS:
VERBAL, VISUAL,
MUSICAL, ETC, ETC ETC...
ALL OF WHICH
CAN BE HANDLED
BY COMPUTERS.
WHY, A COMPUTER
CAN DELIVER A
HYDROGEN BOMB,
NOT JUST A
PUNCH IN THE
MOUTH !!

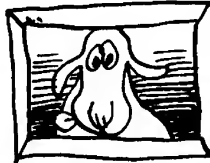
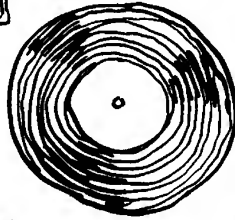


ALL THOSE SIGNALS, INCLUDING A PUNCH IN THE MOUTH,
CAN BE RECORDED IN SOME WAY... SUGGESTING THAT
INFORMATION CAN BE STORED AS WELL AS TRANSMITTED
AND RECEIVED...

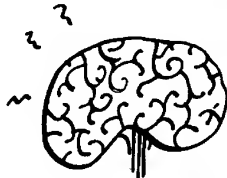


IN BOOKS...

ON AUDIO AND
VIDEO DISKS...

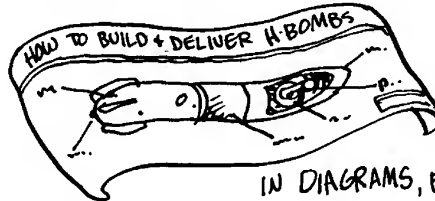


IN PAINTINGS
OR DRAWINGS...



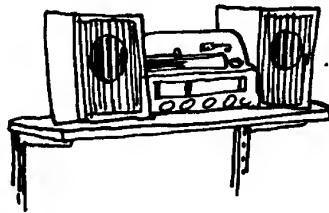
IN THE HUMAN
MEMORY...

ON TAPE...



IN DIAGRAMS, ETC!

THE POINT OF THIS
IS TO TRANSMIT THE
SAME MESSAGE MANY
TIMES...



AND OF COURSE,
ITEMS OF INFORMATION
CAN BE COMBINED
IN VARIOUS WAYS.



AND YOU CAN VERB ANY WORD IN THE LANGUAGE!

WE REFER TO THE STORAGE, TRANSMISSION, COMBINATION, AND COMPARISON OF MESSAGES AS

INFORMATION PROCESSING.

(ALTHOUGH THE COMPUTER INDUSTRY IS GUILTY OF TURNING MANY NOUNS INTO VERBS - ACCESS, INPUT, INTERFACE - "PROCESS" WAS ALREADY A VERB, THANKS TO THE FOOD BUSINESS...)

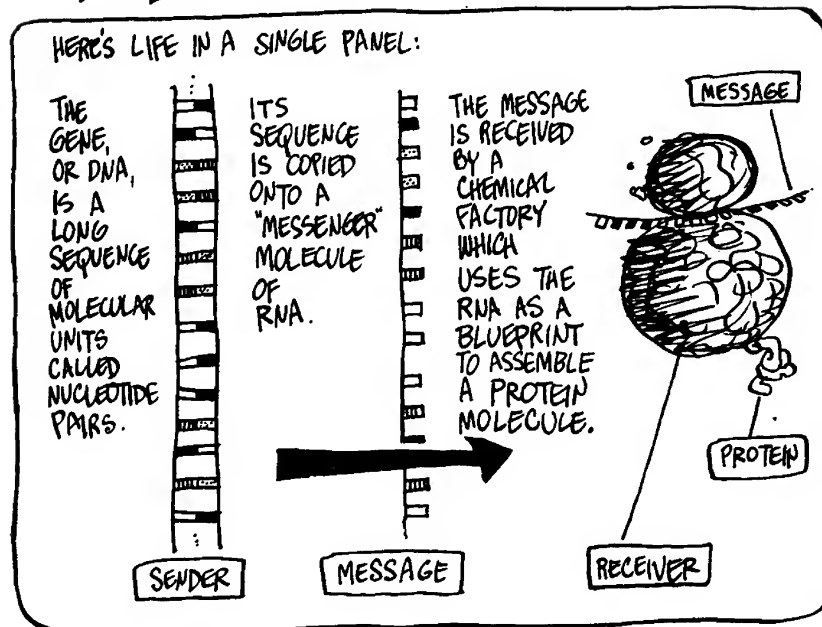
PROCESSED CHEESE

PROCESSED HAM

PROCESSED INFORMATION

TO APPRECIATE THE POWER OF INFORMATION,
CONSIDER ANOTHER EVERYDAY EXAMPLE:

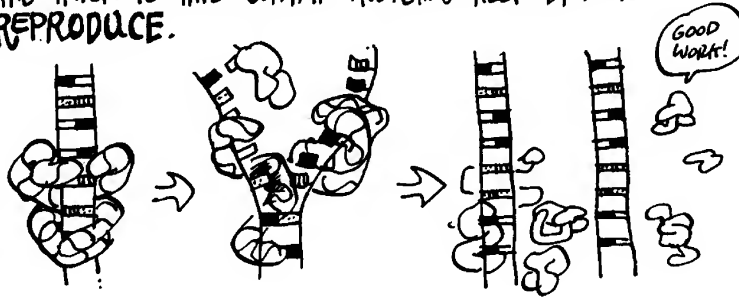
LIFE ITSELF.



IN OTHER WORDS, THE
PROTEIN IS BUILT
ACCORDING TO
INFORMATION
STORED IN THE GENE.



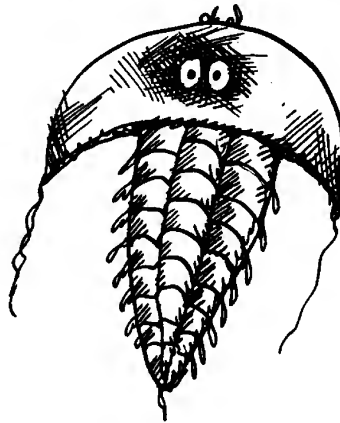
THE TRICK IS THIS: CERTAIN PROTEINS HELP DNA TO REPRODUCE.



WHAT HAPPENS THEN? IF DNA ENCODES PROTEINS THAT HELP DNA TO REPRODUCE, THEN MORE OF THOSE PROTEINS WILL BE BUILT, MORE DNA WILL BE COPIED...ETC! MOREOVER, IF THE DNA ENCODES OTHER PROTEINS WHICH PROTECT IT IN VARIOUS WAYS, AND OTHERS TO ATTACK AND DESTROY RIVAL DNA AND PROTEINS...

THEN THAT DNA-PROTEIN SYSTEM WILL REPRODUCE ITSELF AGAIN AND AGAIN — AND THAT'S WHAT YOU CALL A LIFE FORM.

SO LIFE ITSELF IS A MOLECULAR INFORMATION PROCESSOR, WHICH HAS BEEN RUNNING AUTOMATICALLY FOR OVER 3 BILLION YEARS !!



COMPUTERS SHOULD LAST SO LONG!

The Evolution of the Computer

IT MAY BE GOING TOO FAR TO SAY THAT COMPUTERS HAVE BEEN EVOLVING FROM THE BEGINNING...



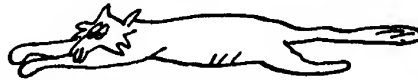
BUT FROM EARLY TIMES, LIFE FORMS HAVE BEEN INCREASING THEIR INFORMATION-PROCESSING ABILITIES. EVEN AN AMOeba RECEIVES CHEMICAL SIGNALS TELLING IT WHERE THE FOOD IS!

OTHERWISE, I'D HAVE TO EAT EVERYTHING!



ALL THE **SENSES** ARE
WAYS OF RECEIVING SIGNALS:

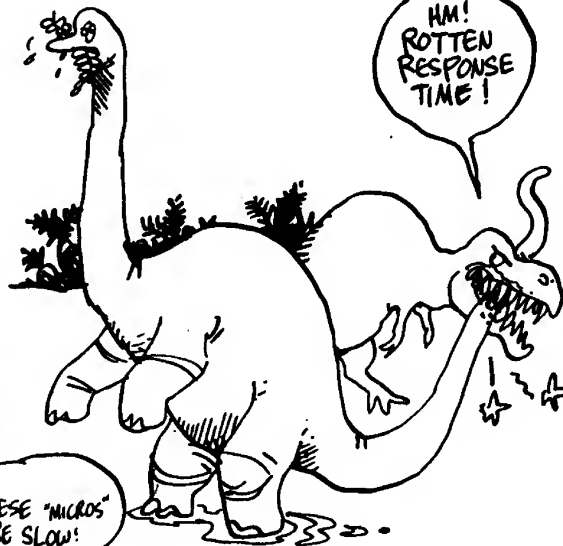
BUT THEN
WE CALL IT
PURR-CENING!



THE EYES PERCEIVE A RANGE OF ELECTROMAGNETIC RAYS;
THE EARS RESPOND TO PRESSURE IN THE AIR; THE NOSE
REACTS TO VARIOUS MOLECULES; SO DO THE TASTE BUDS;
AND THE SENSE OF TOUCH IS A WAY OF RECEIVING A
PUNCH IN THE MOUTH!

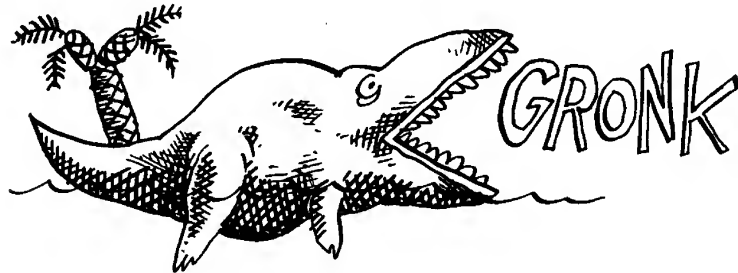
SENSORY
IMPRESSIONS
ARE
TRANSMITTED
ELECTRICALLY
ALONG THE
NERVES
AND
COORDINATED
BY THE
BRAIN —
NATURE'S
FIRST ATTEMPT
TO BUILD
A COMPUTER !!

HM!
ROTTEN
RESPONSE
TIME!

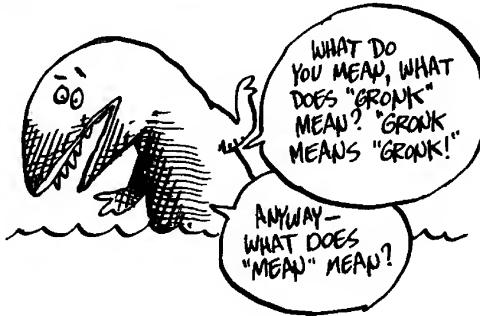


THESE "MICROS"
ARE SLOW!

BESIDES TRANSMITTING INFORMATION WITHIN THEIR OWN BODIES, ANIMALS ALSO SENT MESSAGES TO EACH OTHER:



NOTE AGAIN: THESE DO NOT NECESSARILY CONVEY INFORMATION THAT CAN BE EXPRESSED IN WORDS!



ALSO:
THESE MESSAGES ARE NOT ALWAYS IN THE FORM OF SOUNDS. DOGS COMMUNICATE BY WAGGING THEIR TAILS, AND BEES CAN DESCRIBE THE PRECISE LOCATION OF A FLOWER BY "DANCING."



WHEN HUMANS BEGAN COMMUNICATING, THEY PROBABLY WEREN'T MUCH DIFFERENT FROM ANY OTHER ANIMAL.

"GRONK"



BUT AS THE BRAIN INCREASED IN SIZE AND "COMPUTING POWER," LANGUAGE BECAME MORE EXPRESSIVE.

The reason?

PEOPLE COULD REMEMBER AND USE MORE WORDS. THE MORE WORDS THEY USED, THE GREATER THE NUMBER OF POSSIBLE MESSAGES — WHICH IS ANOTHER WAY OF SAYING THEY COULD SEND MORE INFORMATION.

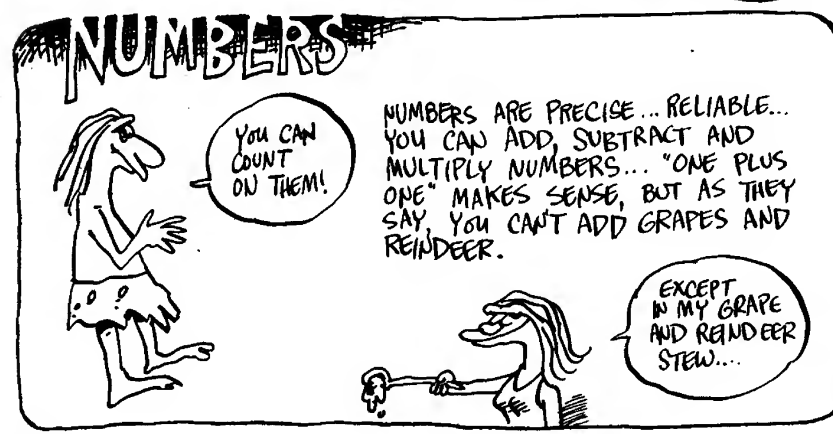
THE SKY IS BLUE...

THE SKY IS BLUE AND FLECKED WITH CLOUDS...

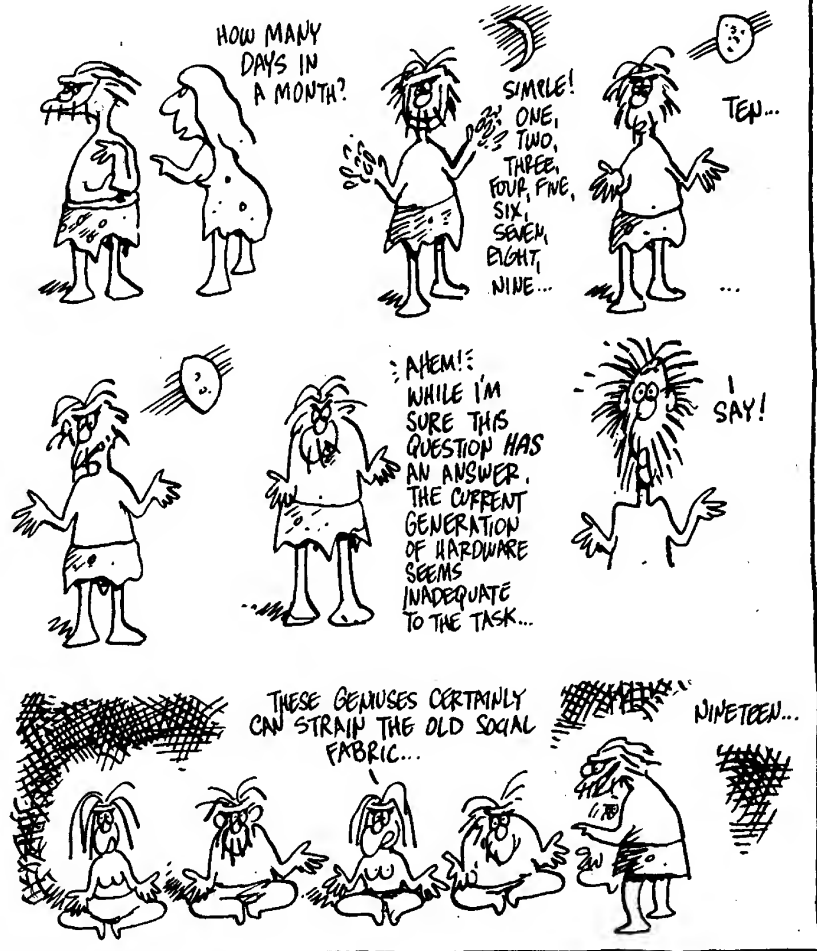


THE SKY, CLEARING AFTER YESTERDAY'S RAIN, IS BLUE AND FLECKED WITH CLOUDS.





NUMBERS ARE ALSO UNIQUE IN THAT YOU "DO THEM" ON YOUR FINGERS, WHILE OTHER PARTS OF LANGUAGE HAPPEN MAINLY IN YOUR HEAD... YES, COUNTING HAS BEEN **DIGITAL** * FROM THE BEGINNING!



* "DIGIT" MEANS FINGER!

NOW, HAVING COUNTED,
WAS THERE SOME
WAY TO SAVE
THE RESULT?

YES!!
AFTER
COMPUTATION,
AMPUTATION!

YOU'RE
MAD!



EXCUSE ME...
I HAVE TO GO PUT
MY HANDS TO BETTER
USE...

WASH THE
WOLF...

SHARPEN
ROCKS...

37



LIKE OTHER ANIMALS, A HUMAN AT FIRST COULD ONLY
RETAIN INFORMATION IN THE BRAIN, WHICH HAD
A LIMITED CAPACITY. (STILL DOES!!) SO THE HUMAN
INVENTED DEVICES TO STORE INFORMATION EXTERNALLY.

THE EARLIEST
KNOWN EXAMPLES
OF EXTERNAL
STORAGE ARE
ABOUT 20,000 YEARS
OLD, LIKE THIS
TALLY BONE,
APPARENTLY USED
TO COUNT THE
DAYS OF THE MONTH.

NOW I
CAN KEEP
TRACK OF MY
INTERNAL
STORAGE!



AROUND THE SAME TIME, CAVE DWELLERS WERE BEGINNING TO STORE ANOTHER KIND OF INFORMATION AS WELL: THEY PAINTED REALISTIC ANIMALS ON THE WALLS OF THEIR CAVES — NO ONE KNOWS WHY!



SEVERAL THOUSAND YEARS LATER, THE SUMERIANS DEvised A SYSTEM TO REPRESENTING THEIR ENTIRE LANGUAGE IN PICTURES:



CALL IT "VISI-TALK!"

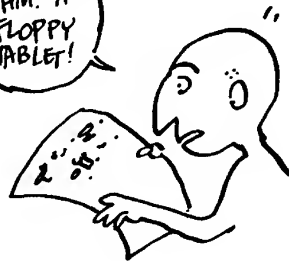
AND SO WRITING WAS BORN!

UNTIL SOMEONE CAN IMPROVE ON LANGUAGE ITSELF, WRITING WILL BE THE ULTIMATE HUMAN SYSTEM OF INFORMATION STORAGE. IT'S NEARLY UNIVERSAL! PEOPLE ALL AROUND THE WORLD INVENTED SYMBOL SYSTEMS TO ENCODE SPOKEN LANGUAGE. OF COURSE, TECHNIQUES VARIED FROM PLACE TO PLACE...



THE SUMERIANS WROTE ON CLAY TABLETS, WHILE THE EGYPTIANS USED SOFT PAPYRUS.

HM! A FLOPPY TABLET!

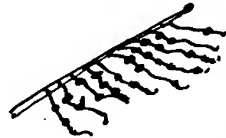


CHINESE WRITING BEGAN WITH MESSAGES TO THE GODS INKED ON TORTOISE SHELLS.

THEY DIDN'T ASK THE GOD OF TORTOISES!

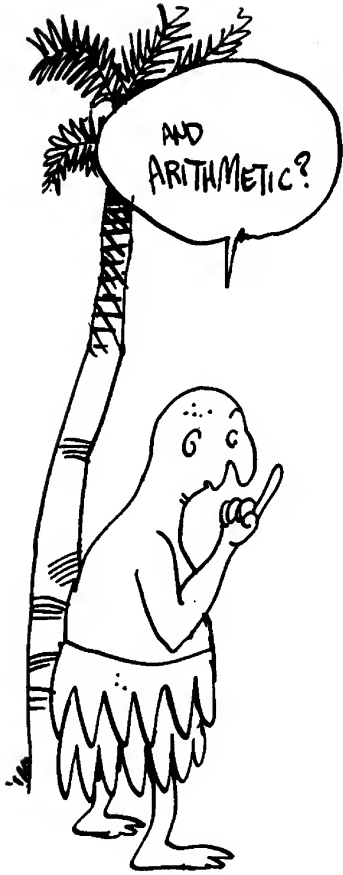


THE INCAS USED A SYSTEM OF KNOTTED CORDS.



GREAT! NOW THAT WE'VE STORED ALL THAT INFORMATION, HOW DO WE FIND IT AGAIN?

WE'LL RETURN TO THAT POINT LATER!



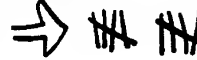
ALL THE EARLY CIVILIZATIONS HAD WAYS OF REPRESENTING NUMBERS THAT WERE FAR ADVANCED OVER THE STONE AGE TALLY BONE, ON WHICH THE NUMBER IS SIMPLY MADE BY PILING UP 1's. NOT TOO USEFUL...



I CALL THIS NUMBER "SMERG."



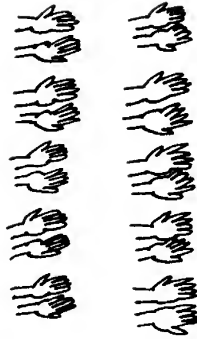
SOMETIME BETWEEN TALLY BONE AND CIVILIZATION, PEOPLE BECAME ACCUSTOMED TO COUNTING BY FIVES AND TENS — FOR AN OBVIOUS REASON: IT WAS HANDY.



LET'S CALL
TEN A "HANDFUL"
AND DO SOME
COUNTING.
FIRST COME
SOME NUMBERS
LIKE
"TWO HANDFULS
AND THREE."



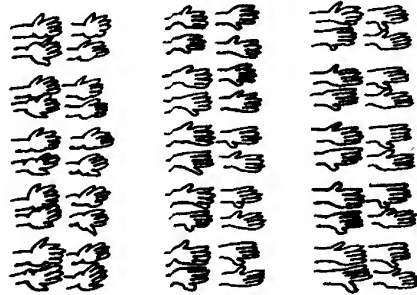
AFTER A
WHILE,
YOU REACH
A
HANDFUL
OF
HANDFULS
(TEN TENS,
OR A
HUNDRED).



THEN COMES A HANDFUL OF HANDFULS AND ONE...

" " " " " TWO
" " " " " A HANDFUL...

EVENTUALLY, YOU'RE
LOOKING AT
NUMBERS LIKE
THIS:

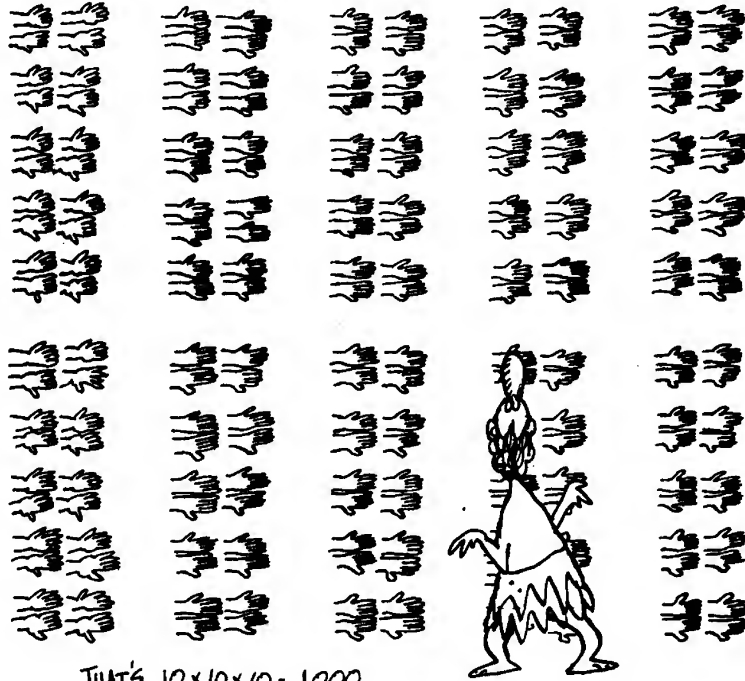


OR: "THREE HANDFULS OF
HANDFULS, FOUR HANDFULS,
AND SEVEN."



A MOUTH-
FUL!

AND THEN — A HANDFUL OF HANDFULS OF HANDFULS.



THAT'S $10 \times 10 \times 10 = 1000$.

NEXT COMES

TEN THOUSAND...

A HUNDRED THOUSAND...

A THOUSAND THOUSAND...

TEN THOUSAND THOUSAND...

EACH OF WHICH IS
A HANDFUL OF
THE ONE BEFORE!



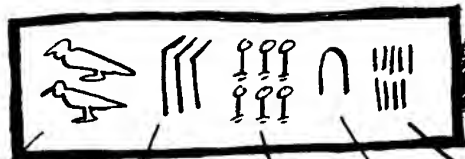
THIS IS
GETTING OUT
OF HAND!

THE ANCIENTS FOUND TWO BASIC WAYS TO TRANSLATE THIS INTO WRITING:

ONE, THE EGYPTIAN SYSTEM, USED A DIFFERENT SYMBOL FOR EACH NEW HANDFUL.

| = ONE ∩ = TEN 9 = HUNDRED
 I = THOUSAND [= TEN THOUSAND [bird] = HUNDRED THOUSAND

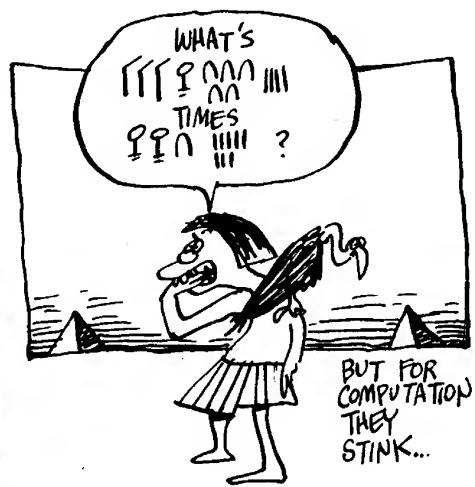
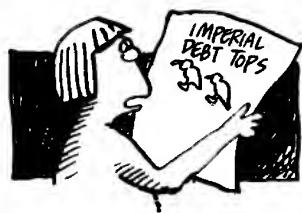
THEN YOU JUST PILE THEM UP:



TWO HUNDRED THOUSANDS THREE TEN THOUSANDS SIX THOUSANDS ONE TEN NINE UNITS

OR 236,019

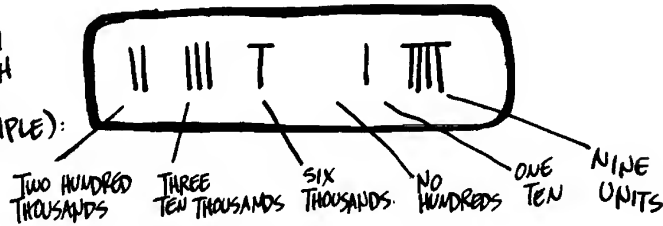
ASIDE FROM HAVING A CERTAIN GRAPHIC CHARM, THESE NUMERALS ARE VERY EASY TO READ, ONCE YOU'RE USED TO THEM (JUST AS "3 BILLION" READS QUICKER THAN "3000 000 000").



ON THE OTHER HAND, THE CHINESE USED THE POSITION OF NUMERALS TO INDICATE THEIR VALUE. FIRST THEY COUNTED FROM ONE TO NINE:

					T	TT	TTT	TTTT
1	2	3	4	5	6	7	8	9

FROM WHICH (FOR EXAMPLE):

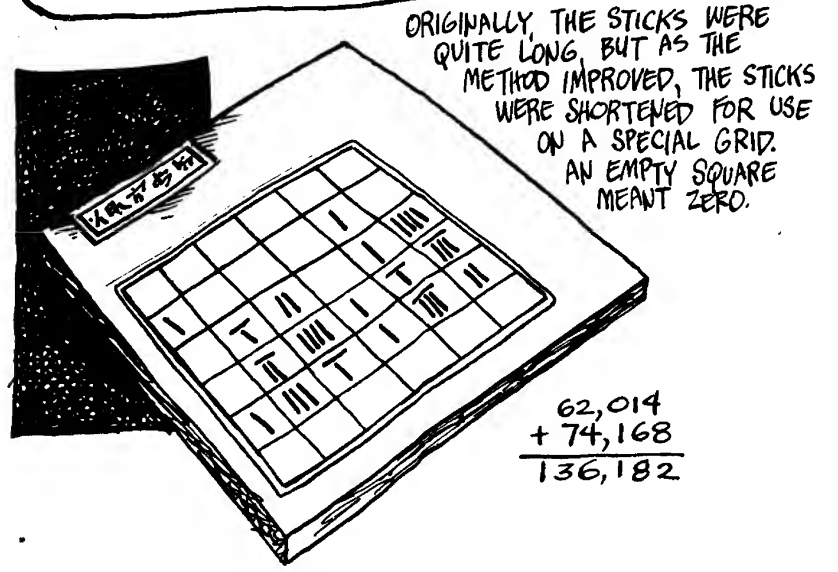
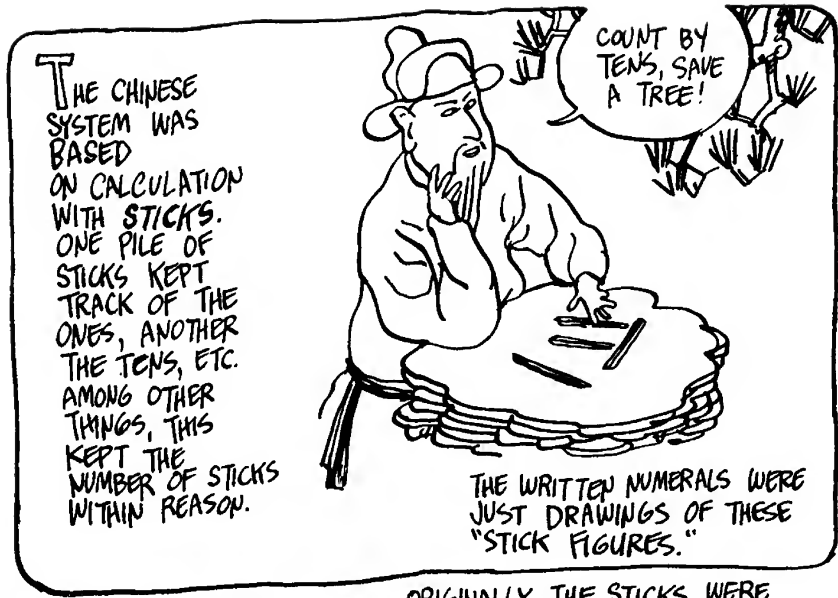


OR 236,019.

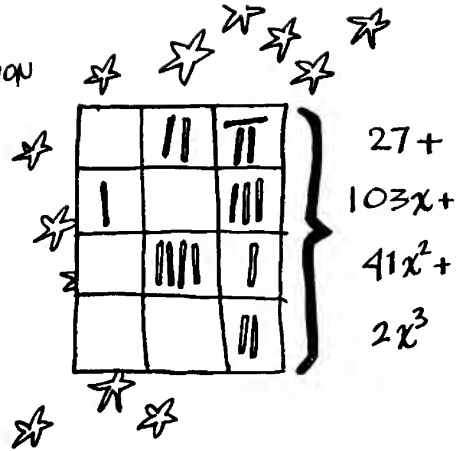
EXCEPT FOR THE UNFAMILIAR NUMERALS, THIS SYSTEM IS NEARLY THE SAME AS OUR OWN.

THE ONLY DIFFERENCE IS THAT IT LACKED A SYMBOL FOR ZERO. THE CHINESE JUST LEFT A BLANK INSTEAD.

IN PRACTICE, THIS WAS MUCH LESS OF A PROBLEM THAN IT MIGHT HAVE BEEN, BECAUSE THE CHINESE DID NOT CALCULATE ON PAPER !!!

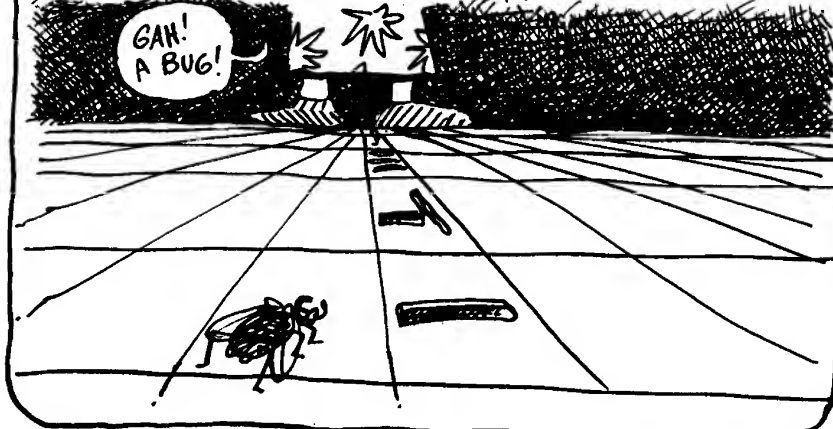


BESIDES ADDITION, SUBTRACTION, MULTIPLICATION AND DIVISION, THIS SWAN-PAN, OR "ARITHMETIC TABLE," WAS ALSO APPLIED TO ALGEBRA AND THE SOLUTION OF EQUATIONS. ENTRIES IN THE SQUARES BECAME THE COEFFICIENTS OF ALGEBRAIC EXPRESSIONS.



THIS TECHNIQUE HAD THE PICTURESQUE NAME OF "THE METHOD OF THE CELESTIAL ELEMENT."

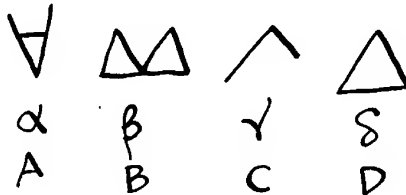
AFTER BORROWING THE DESIGN, THE JAPANESE USED IT TO CALCULATE π (PI) TO 50 DECIMAL PLACES. ONE JAPANESE MATHEMATICIAN WAS SAID TO HAVE CONVERTED A WHOLE ROOM INTO A "MAINFRAME" SWAN-PAN.



MEANWHILE, BACK AT THE MEDITERRANEAN,
THEY HAD MADE TWO GREAT INVENTIONS: THE

ALPHABET & ABACUS.

THE ALPHABET
RANKS AS ONE
OF THE GREAT
IDEAS IN THE
HISTORY OF
INFORMATION.



Before

THE ALPHABET, A
SEPARATE SYMBOL
WAS NEEDED
FOR EVERY WORD
(OR EVERY SYLLABLE,
IN SOME CASES).
TO LEARN WRITING,
ONE HAD TO
MEMORIZE THOUSANDS
OF SYMBOLS.

WE CHINESE
ARE STILL
SADDLED
WITH
PICTOGRAMS!



After

DECOMPOSING LANGUAGE
INTO MORE BASIC SOUNDS,
THE NUMBER OF SYMBOLS
WAS REDUCED TO FEWER
THAN 30. NOW, ANY
IDIOT COULD LEARN
TO READ!

WHEREAS
PREVIOUSLY,
ONLY IDIOTS
WITH LEISURE
COULD LEARN...



THERE'S A LESS OBVIOUS ADVANTAGE OF THE ALPHABET,
BUT NO LESS IMPORTANT:

alphabetical order.

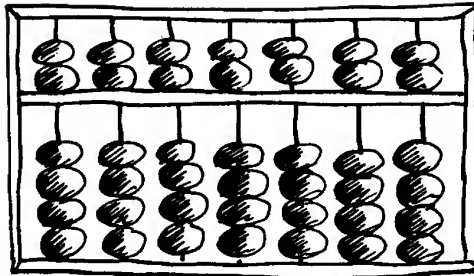


WITH THOUSANDS OF PICTOGRAMS, ANY FILING SYSTEM HAS TO BE COMPLICATED, BUT GIVEN THE ORDER OF AN ALPHABET, YOU CAN PUT WORDS IN ORDER, TOO. IMAGINE USING A PHONE BOOK, DICTIONARY, OR LIBRARY WITHOUT ALPHABETICAL ORDER!

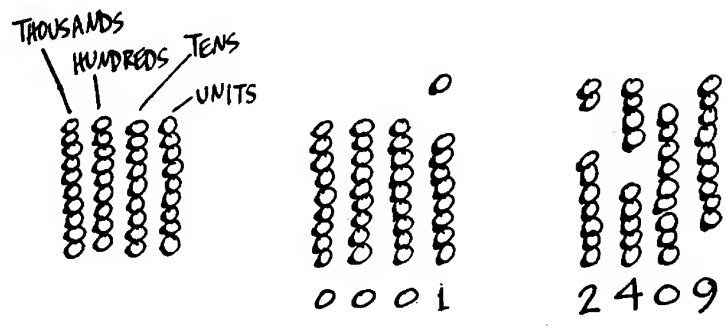


COMPUTERS SPEND A GOOD PART OF THEIR TIME JUST PUTTING THINGS IN ORDER!

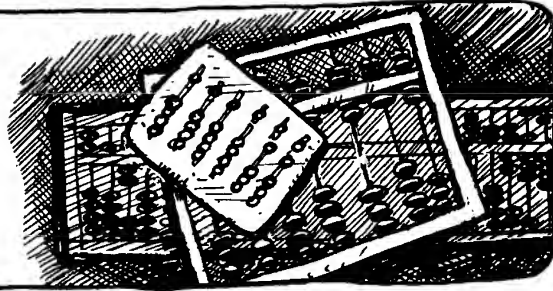
THE ABACUS, ORIGINALLY A PRODUCT OF THE MIDDLE EAST, IS A FULL-BLOWN HAND-HELD DECIMAL CALCULATOR.



LIKE THE ALPHABET, THE ABACUS WAS SIMPLE, SYSTEMATIC AND SPEEDY. IN ITS SIMPLEST FORM, AN ABACUS WAS JUST A FEW COLUMNS OF PEBBLES. A PEBBLE IN A GIVEN COLUMN IS WORTH TEN PEBBLES IN THE COLUMN TO ITS IMMEDIATE RIGHT. NUMBERS ARE ENTERED BY PUSHING UP PEBBLES.



THE ABACUS HAS SEEN MANY INCARNATIONS AND BEEN USED IN MOST PARTS OF THE OLD WORLD.



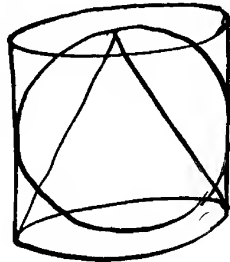
WE KNOW FROM PICTURES THAT THE ANCIENT GREEKS HAD THE ABACUS, BUT THEIR MATHEMATICIANS NEVER DISCUSSED IT. (GREEK INTELLECTUALS LOOKED DOWN ON THE WORK OF THE HANDS...)



THE TOOL OF SHOPKEEPERS!

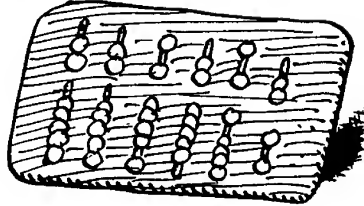


 THIS MAY HAVE BEEN WHY GREEK MATHEMATICIANS CONCENTRATED ON GEOMETRY...



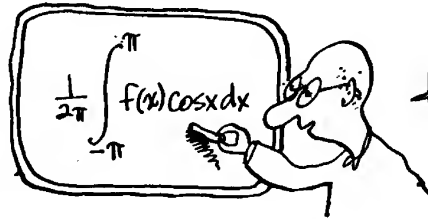
The Romans

ALSO USED THE ABACUS. THEIRS CONSISTED OF MARBLES SLIDING ON A GROOVED BRONZE PLATE:



IT CONTRIBUTED A COUPLE OF MATHEMATICAL WORDS TO ENGLISH:

IN LATIN, **CALX** MEANT LIMESTONE OR MARBLE...
So



FROM WHICH COMES "CHALK!"

So **CALCULUS**

WAS AN ABACUS PEBBLE... AND DOING ARITHMETIC WAS **CALCULATION.**

THE ROMANS DID NOT CALCULATE WITH ROMAN NUMERALS!!

WHAT'S MXVIII TIMES CLXVI?

BEATS ME! I LOST MY MARBLES...





...AND FELL...
ROME WAS SACKED...
CHRISTIANITY ROSE
FROM ITS
ASHES... CLASSICAL
LEARNING VANISHED
IN THE WEST...
AND ONLY A FEW
MATH PROBLEMS
REMAINED
LEGITIMATE, LIKE
COMPUTING THE
DATE OF EASTER...
OR HOW MANY
ANGELS FIT ON
THE HEAD OF
A PIN...



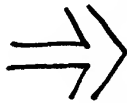
+++IN SUM?++++

ANCIENT TIMES
WERE REALLY THE

AGE OF
CALCULATORS.



ET TO?
ET THREE!



ALTHOUGH THE ANCIENTS HAD WAYS OF WRITING
NUMBERS, THEY RARELY CALCULATED IN
WRITING.

THIS IS NOT SO EASY TO APPRECIATE FOR THOSE OF US
WHO WERE RAISED ON PENCIL AND PAPER.

SO THE NEXT
TIME YOU
HEAR SOMEONE
COMPLAIN THAT
ELECTRONIC
CALCULATORS
ARE
RUNNING
ARITHMETIC...

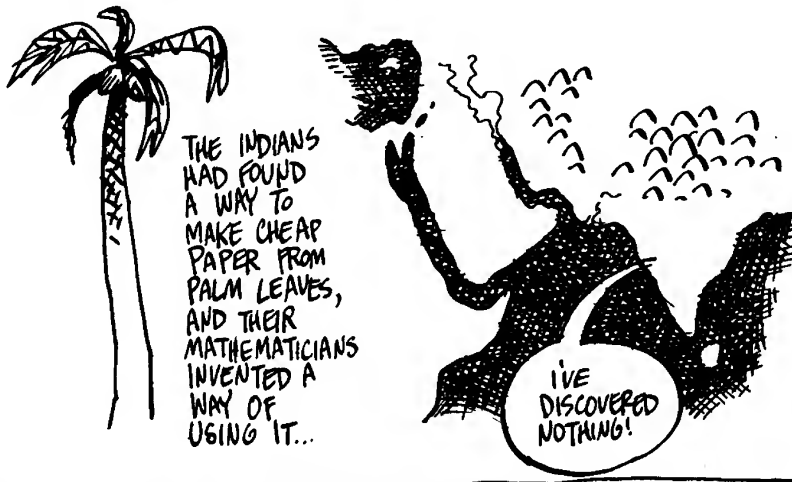
HOW CAN
WE REMEMBER
OUR MULTIPLICATION
TABLES?



...SIMPLY REPLY THAT
PEOPLE SURVIVED WITH
CALCULATORS FOR
MORE THAN 4000 YEARS!!

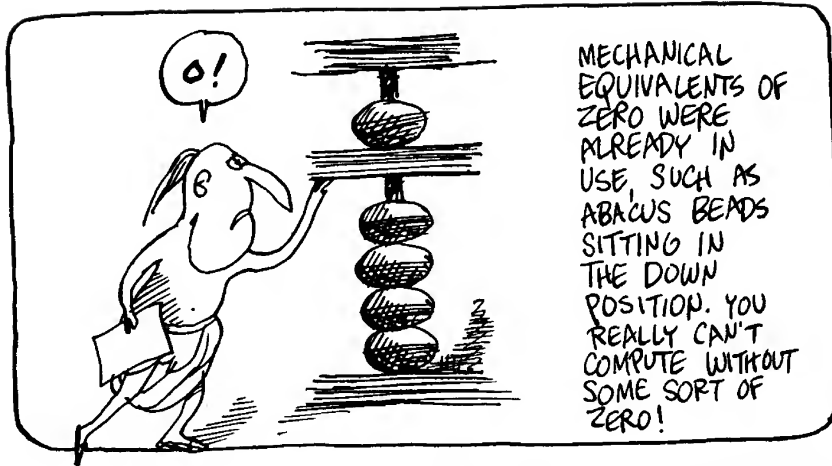
Much ado about NOTHING

AS FAR AS CALCULATION GOES, THE AGE OF PAPER BEGAN IN INDIA, ABOUT 650 A.D.

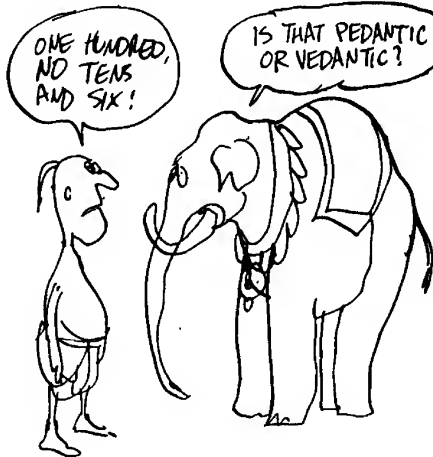


TO DO SO, THEY DEVISED A SYMBOL FOR ZERO!





WHY HADN'T ANYONE THOUGHT OF PUTTING IT IN WRITING BEFORE? MAYBE BECAUSE WRITING WAS FOR REPRESENTING SPOKEN LANGUAGE, AND NOBODY SAYS -



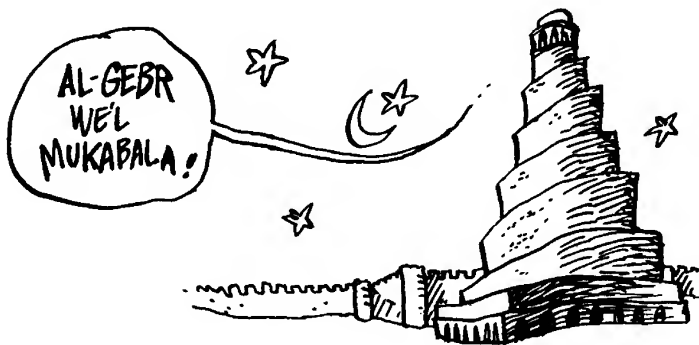
BUT FOR SOME REASON, THE HINDUS INVENTED A WRITTEN ZERO!

१ २ ३ ४ ५ ६ ७ ८ ९ ०

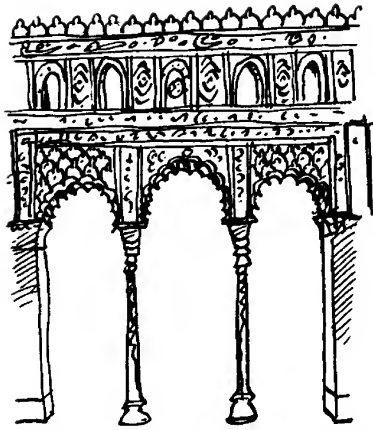
THE INDIAN MATH WAS PICKED UP BY THE ARABS, WHO SPREAD IT ALL THE WAY WEST TO SPAIN.



AROUND THE YEAR 830, A PERSIAN SCHOLAR WROTE THE STANDARD TEXT BOOK ON THE SUBJECT. HIS NAME WAS MOHAMMED IBN MUSA ABU DJEFAR, BUT HE WAS KNOWN AS AL-KHWARISMI. AND THE SUBJECT OF HIS BOOK?



OR ALGEBRA, FOR SHORT.



BY THE 1100'S, MUSLIM CIVILIZATION HAD GROWN SO MAGNIFICENT THAT THE EUROPEANS WERE BEGINNING TO WONDER...



WHY SHOULD GOOD CHRISTIANS REMAIN IN THE DARK AGES?

A FEW INTREPID INFIDELS WENT TO LIVE AMONG THE ARABS, LEARNED THEIR LANGUAGE, SNUCK INTO THEIR UNIVERSITIES, AND TRANSLATED THEIR CLASSICS INTO LATIN.



IN AL-KHAWARISMI'S BOOK THEY FOUND THE INDIAN NUMERALS.

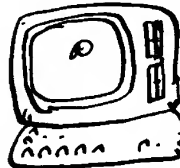


AL-KHAWARISMI
AL-KARISMI
ALGARISMI
ALGORISMI

PRONOUNCED OFTEN ENOUGH, THE MATHEMATICIAN'S NAME WAS TURNED INTO

ALGORISM

WHICH IS WHAT THE EUROPEANS CALLED THE NEW SYSTEM OF CALCULATION.



YETH, VERY NITHE...

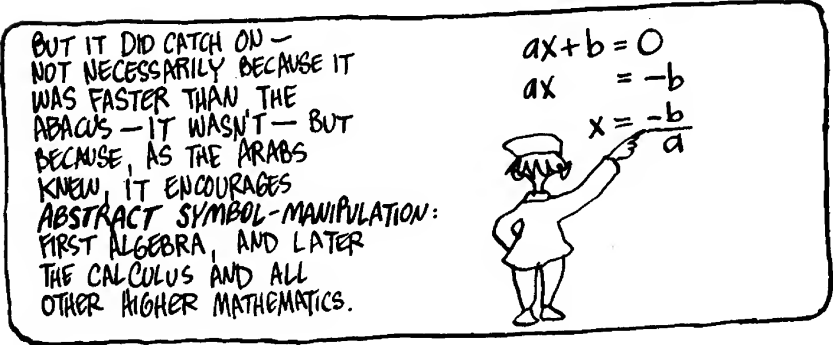
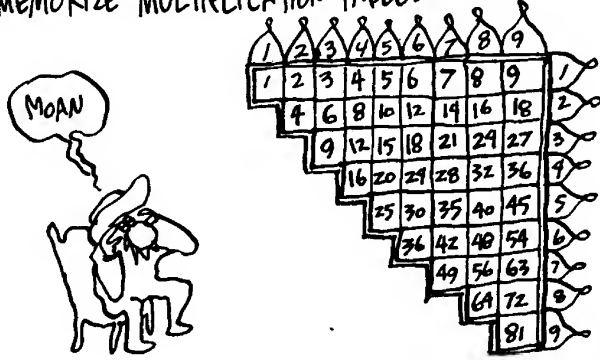
FROM THE SAME ROOT COMES

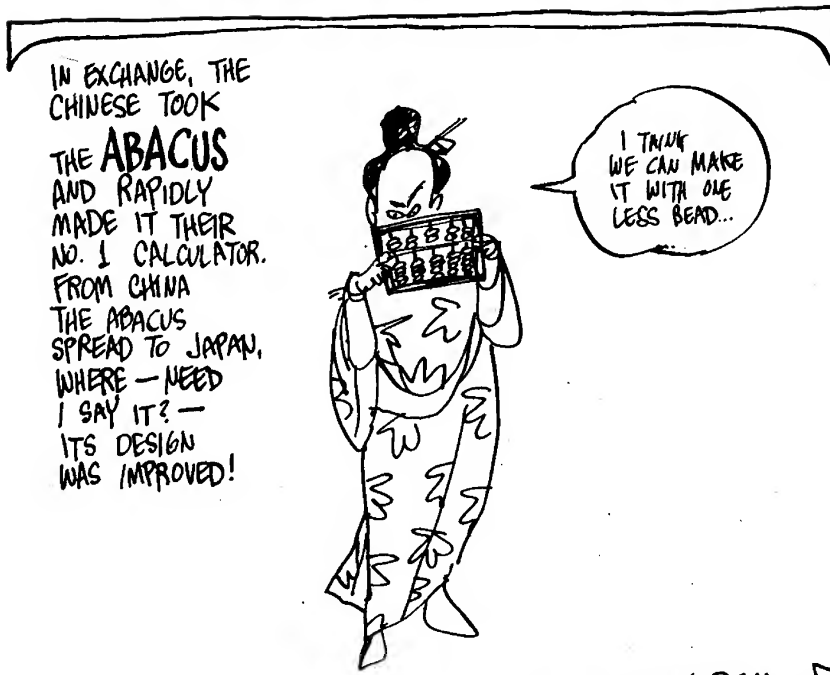
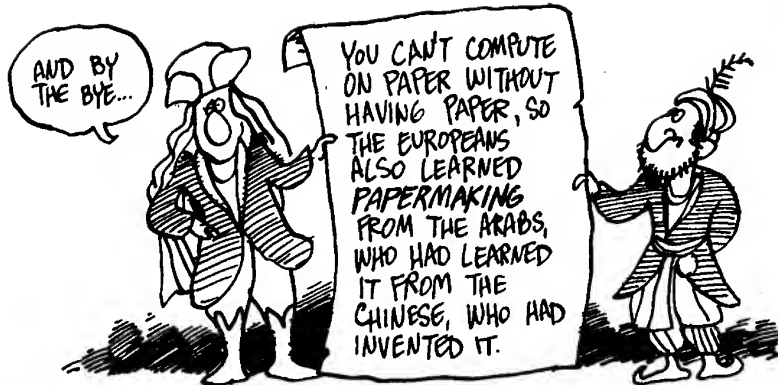
ALGORITHM,

A COMPUTER WORD WE'LL EXPLORE IN A BIT...



... AND EVERYONE AGREED IT WAS A PAIN TO MEMORIZE MULTIPLICATION TABLES..





BUT BACK TO ALGORISM... ⇒

WHILE EUROPEAN
SCHOLARS WERE
TRANSLATING THE CLASSICS
IN ARAB LIBRARIES,
THE **CRUSADERS**
WERE DOING THEIR
BEST TO DESTROY
ISLAMIC CIVILIZATION.

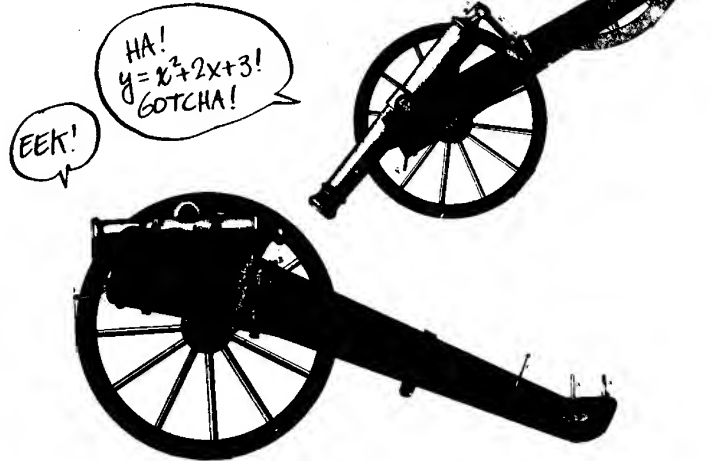


THIS DOUBLE-PRONGED ACTION OF TRANSLATION AND DESTRUCTION
LED TO THE GROWTH OF EUROPEAN LEARNING AND POWER KNOWN
AS:

The RENAISSANCE.



MILITARY ADVANCES OFTEN
GO HAND IN HAND WITH
MATHEMATICAL ONES.



IN THE 1500'S, NICCOLO TARTAGLIA (1499-1559) COMPUTED
THE PATHS OF CANNONBALLS (AN IMPORTANT PROBLEM IN THE
LATER HISTORY OF COMPUTERS, AS WE'LL SEE).

JUST OVER A
CENTURY LATER,
ISAAC NEWTON
UNIFIED THE MOTIONS
OF CANNONBALLS
AND PLANETS WITH
THE THEORY OF
GRAVITATION, ONE
OF THE AGE OF
PENCIL AND PAPER'S
CROWNING GLORIES.



HOWEVER,
THE THEORY
INTRODUCED
SOME
REAL
COMPU-
TATIONAL
HORRORS...

THE WORST WAS THE **THREE BODY PROBLEM**, WHICH ASKS FOR A MATHEMATICAL DESCRIPTION OF THE MOTIONS OF THREE BODIES — SUN, EARTH, AND MOON, FOR EXAMPLE — ACTING UNDER THE INFLUENCE OF GRAVITY. THIS TURNS OUT TO BE INCREDIBLY DIFFICULT AND TEDIOUS!

WE'RE REACHING THE LIMITS OF PAPER!



SO A NUMBER OF SCIENTISTS BEGAN THINKING AGAIN ABOUT CALCULATION BY MACHINE...

John
NAPIER, (1550-1617),

A HALF-MAD SCOT MOST
FAMOUS FOR LOGARITHMS,
DEvised "NAPIER'S BONES."



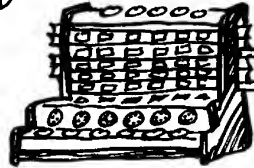
THESE
WERE
SIMPLY
MULTIPLI
CATION TABLES
ON A STICK.

THE FIRST REAL MACHINE WAS
BUILT BY Wilhelm

SCHICKARD (1592-
1635).

IT COULD ADD, SUBTRACT,
MULTIPLY, AND DIVIDE...
BUT WAS LOST IN THE
30-YEAR WAR.

SCHICKARD
HIMSELF
DIED OF
PLAGUE
AND
COULDN'T
DEFEND
HIS PRIORITY,
SO...



Blaise
PASCAL (1623-1662)

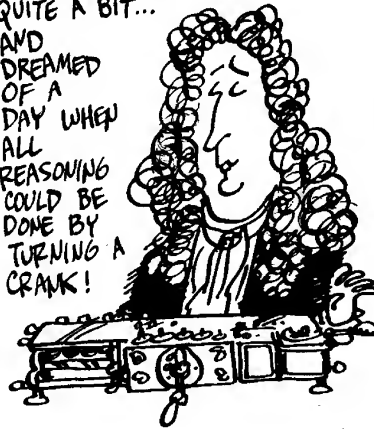
USUALLY GETS CREDIT
FOR BUILDING
THE FIRST
CALCULATOR.
HIS
"PASCALINE"
COULD ONLY
ADD AND
SUBTRACT.



Gottfried Wilhelm
LEIBNIZ (1646-1716)

IMPROVED PASCAL'S DESIGN
QUITE A BIT...

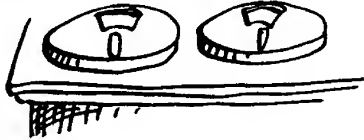
AND
DREAMED
OF A
DAY WHEN
ALL
REASONING
COULD BE
DONE BY
TURNING A
CRANK!



DURING THE 1700's,
MORE SUCH MACHINES
WERE BUILT, BUT
ALL FELL FAR
SHORT OF BEING
ANYTHING LIKE
A GENERAL-PURPOSE
COMPUTER.



FOR EXAMPLE: IN EVERY CASE, THE USER ENTERED
NUMBERS BY SETTING A ROW OF WHEELS OR KNOBS ...



... AND THEN
TURNED THE
APPROPRIATE
CRANK TO ADD
OR MULTIPLY.

ANOTHER WAY
OF SAYING THE
SAME THING:



THE **INPUT**
CONSISTED ONLY OF THE
NUMBERS TO BE COMBINED.

AS WILL BE PLAIN
SOON ENOUGH, AN
ALL-PURPOSE COMPUTER
MUST ALSO BE ABLE TO
DO MORE: IT MUST READ
INSTRUCTIONS
ABOUT WHAT TO DO
WITH THOSE NUMBERS!!

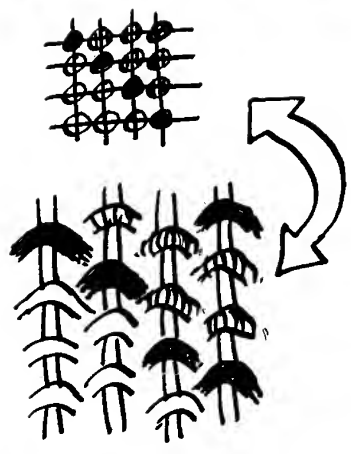


WELL, I
WAS ONLY
TRYING TO
MAKE AN
ADDING
MACHINE...

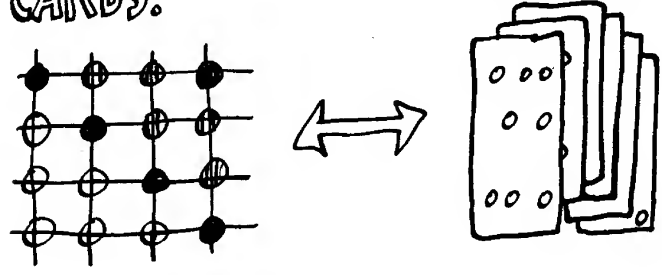
THE GERM OF THIS IDEA CAME NOT FROM THE LAB OR A SCIENTIST'S STUDY, BUT THE SOOTY FACTORIES OF THE

INDUSTRIAL REVOLUTION

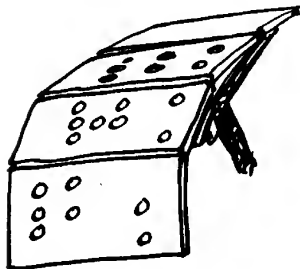
YOU MAY NEVER HAVE THOUGHT OF A **WEAVER'S LOOM** AS AN INFORMATION PROCESSOR, AND YET: IT TRANSLATES AN ABSTRACT DESIGN INTO A PATTERN OF COLORS, CREATED BY LOOPING OVER EACH COLORED THREAD AT THE APPROPRIATE PLACE.



IN THE MID-1700'S, A SYSTEM WAS INVENTED FOR REPRESENTING THESE PATTERNS ON **PUNCHED CARDS**.



WITH AN OLD-FASHIONED
HANDLOOM, THE WEAVER
READ THE CARDS,
BUT IN 1801, JOSEPH MARIE
JACQUARD
INVENTED A POWER
LOOM WITH AN
AUTOMATIC CARD READER.



IN WENT THE CARDS, OUT CAME THE CLOTH...



ACROSS THE ENGLISH CHANNEL, JACQUARD'S IDEA SET OFF A CHAIN REACTION IN THE BRAIN OF

CHARLES BABBAGE

(1792-1871),

WHO HAS BECOME KNOWN AS THE "FATHER OF THE COMPUTER."



I SEE NO FAMILY RESEMBLANCE!

FOR SEVERAL YEARS BABBAGE, A CAMBRIDGE MATH PROFESSOR, HAD BEEN WORKING ON A LARGE MECHANICAL CALCULATOR HE CALLED "THE DIFFERENCE ENGINE."



MY ORIGINAL BRAINSTORM...

IT WOULD HAVE COMPUTED MATHEMATICAL TABLES, IF THE INVENTOR HAD EVER BEEN ABLE TO FINISH IT.

IN 1821,
BABBAGE
APPLIED
TO THE
ROYAL SOCIETY
FOR FUNDS
TO BUILD
THE DIFFERENCE
ENGINE,
AND THEY
GAVE HIM
A SIZABLE
GRANT.



HE HIRED A MASTER
MACHINIST AND
WENT TO WORK...
BUT BABBAGE
COULDN'T RESIST
THINKING UP NEW
IMPROVEMENTS IN
THE MIDST OF
PRODUCTION!

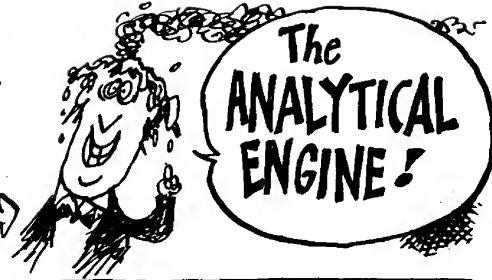


JOLLY
GOOD SHOW,
OLD CHAP!!! NOW
TRY IT AGAIN,
ACCORDING TO
THE NEW SPECS!

MEANWHILE, HIS
HYPERACTIVE MIND
KEPT TURNING TO
NEW PROJECTS:
LIFE INSURANCE TABLES,
LIGHTHOUSE SIGNALS,
GLASS CUTTING, AND
EVEN VOLCANOES.
(HE MIKED INTO A
LIVE ONE!)



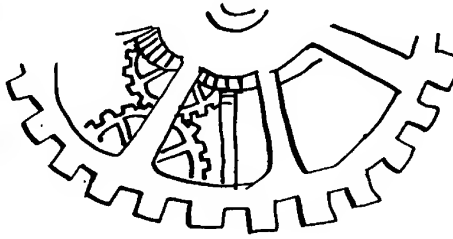
THAT'S HOW MATTERS STOOD WHEN JACQUARD'S PUNCHED CARDS SET OFF BABBAGE'S NEW BRAINSTORM, A MACHINE HE CALLED:



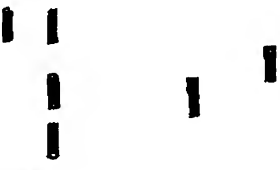

BECAUSE IT SO CLOSELY RESEMBLED A COMPUTER, LET'S TAKE A CLOSER LOOK AT THE ANALYTICAL ENGINE, AS BABBAGE IMAGINED IT. ITS COMPONENTS INCLUDED—

THE MILL:

AT THE ENGINE'S HEART WOULD BE A GREAT NUMBER-CRUNCHER, AN ADDING MACHINE ACCURATE TO 50 DECIMAL PLACES. BABBAGE CALLED THIS THE MILL.



HOW DID IT KNOW WHAT TO DO?



THE INSTRUCTIONS TO THE MILL WOULD BE READ IN ON PUNCHCARDS.

THAT IS, THE PUNCHCARDS CONVEYED NOT ONLY THE NUMBERS TO BE CRUNCHED BUT ALSO THE PATTERN OF CRUNCHING!!



SO THE MACHINE WOULD NEED A SPECIAL
CARD-READING **INPUT** DEVICE.

TO RETAIN NUMBERS FOR FUTURE REFERENCE, BABBAGE ENVISIONED A MEMORY UNIT, OR **STORE**. THIS WAS TO BE A BANK OF 1000 "REGISTERS," EACH A DEVICE CAPABLE OF STORING ONE 50-DIGIT NUMBER. THESE NUMBERS COULD EITHER BE INPUT FROM THE CARDS OR THE RESULT OF COMPUTATIONS IN THE MILL.

FINALLY, THE **OUTPUT**. BABBAGE DESIGNED THE WORLD'S FIRST AUTOMATED TYPE SETTER TO PRINT THE RESULTS OF COMPUTATIONS.

A PUNCHCARD COULD DO ONE OF THE FOLLOWING THINGS:



INPUT A NUMBER TO THE STORE



INPUT A NUMBER TO THE MILL



MOVE A NUMBER FROM THE MILL TO THE STORE



MOVE A NUMBER FROM THE STORE TO THE MILL

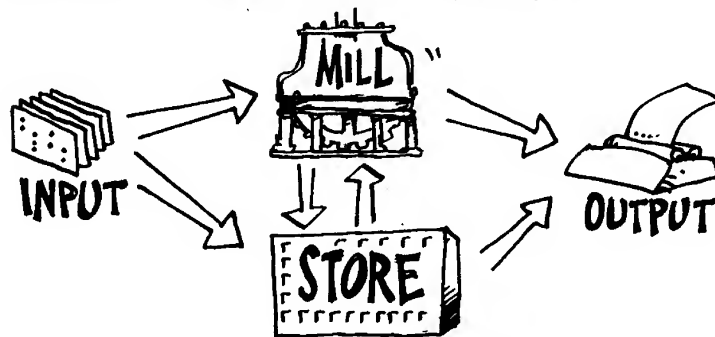


INSTRUCT THE MILL TO PERFORM AN OPERATION

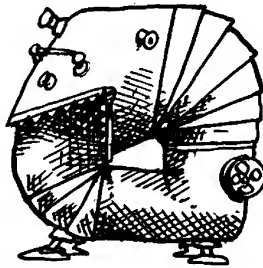


OUTPUT A NUMBER FROM EITHER STORE OR MILL

WHICH MAY BE SUMMARIZED IN THIS DIAGRAM:

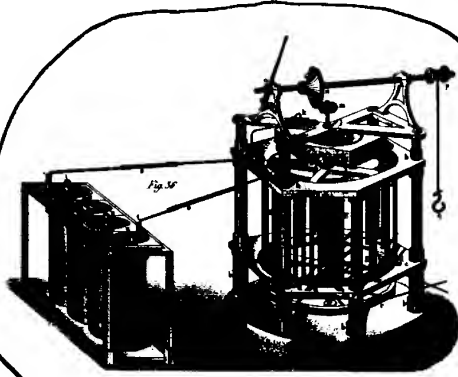


IN PARTICULAR, A RESULT FROM THE MILL COULD BE STORED FOR FUTURE REFERENCE, THEN RETURNED TO THE MILL WHEN NEEDED. AS BABBAGE PUT IT, THE ANALYTICAL ENGINE COULD "EAT ITS OWN TAIL." VERY FLEXIBLE!



YOU HAVE TO BE FLEXIBLE TO EAT YOUR OWN TAIL...

SO FAR, THESE IDEAS WERE STILL ON THE DRAWING BOARD. NOW BABBAGE BEGAN LOOKING FOR SYMPATHETIC SOULS WHO COULD HELP PUT HIS PLANS INTO OPERATION.

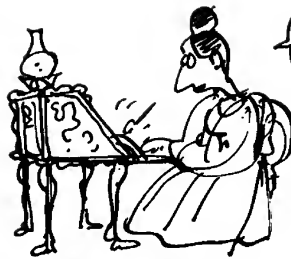


THE MOST SYMPATHETIC WAS

ADA AUGUSTA,

LADY LOVELACE, DAUGHTER OF THE POET LORD BYRON AND AN ENTHUSIASTIC AMATEUR MATHEMATICIAN. IF CHARLES BABBAGE IS THE COMPUTER'S FATHER, ADA LOVELACE IS ITS MOTHER!

ADA BECAME THE FIRST PROGRAMMER: SHE WROTE OUT ACTUAL SEQUENCES OF INSTRUCTIONS FOR THE ANALYTICAL ENGINE...



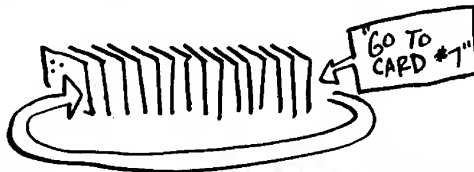
TOO BAD IT DOESN'T EXIST YET...

SHE INVENTED THE SUBROUTINE: A SEQUENCE OF INSTRUCTIONS WHICH CAN BE USED AGAIN AND AGAIN IN MANY CONTEXTS.

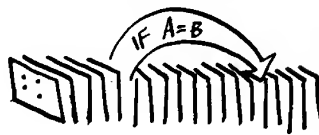


WE CAN HAVE A WHOLE LIBRARY OF THEM!

SHE RECOGNIZED THE VALUE OF LOOPING: THERE SHOULD BE AN INSTRUCTION THAT BACKS UP THE CARD READER TO A SPECIFIED CARD, SO THAT THE SEQUENCE IT INITIATES CAN BE EXECUTED REPEATEDLY.

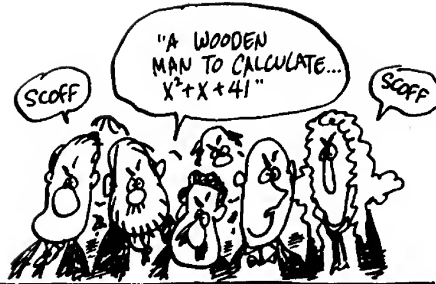


AND SHE DREAMED UP THE CONDITIONAL JUMP: THE CARD READER COULD "JUMP" TO ANOTHER CARD IF SOME CONDITION IS SATISFIED.

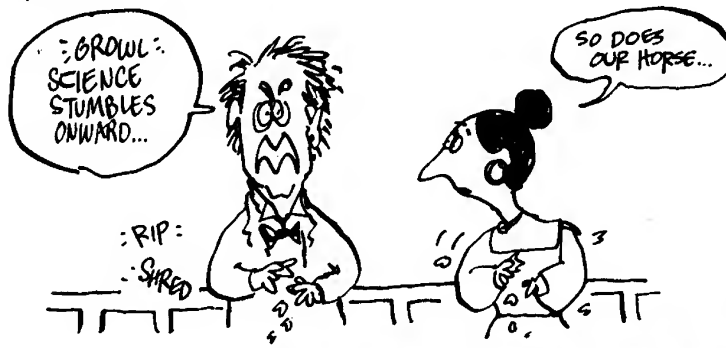


IT CAN MAKE DECISIONS!

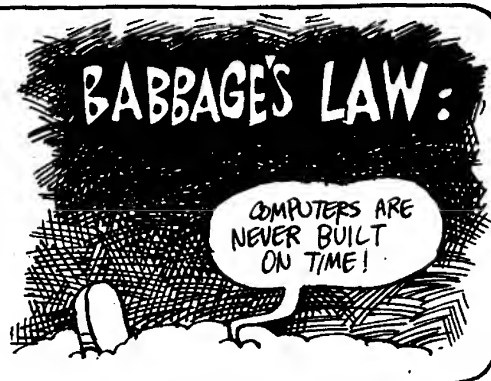
NOT BAD FOR A MACHINE
THAT NEVER EXISTED...
THE GOVERNMENT
REFUSED TO SUPPORT IT,
IN VIEW OF BABBAGE'S
TRACK RECORD WITH
THE DIFFERENCE ENGINE.
THEY CALLED IT:



DESPERATE FOR FUNDS, BABBAGE COOKED UP A "SCIENTIFIC"
RACETRACK BETTING SCHEME — AND SQUANDERED ADA'S
FORTUNE.



THE STORY ENDED
UNHAPPILY: ADA
DIED YOUNG...
AND BABBAGE
NEVER FINISHED
THE ANALYTICAL
ENGINE, WHICH
BECAME THE
FIRST EXAMPLE OF—



THE ILL-STARRED INVENTORS WERE AHEAD OF THEIR TIME. NOTHING EQUIVALENT TO THE ANALYTICAL ENGINE EXISTED UNTIL THE 1940'S.

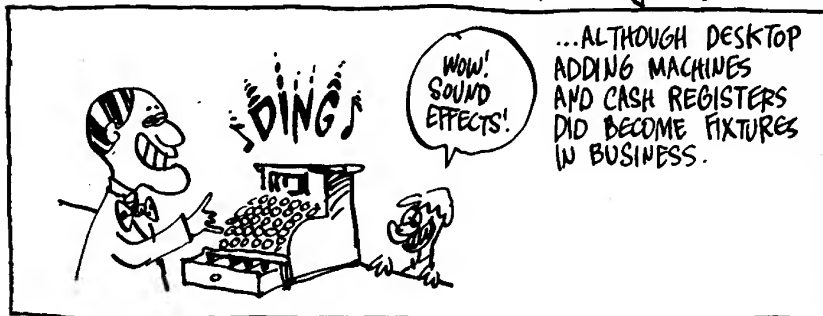
ALMOST AS LATE AS OUR HORSE...



IN THE MEANTIME, MATTERS PROGRESSED IN TWO DIRECTIONS:

ON THE ONE HAND WERE MECHANICAL CALCULATORS: SEVERAL ENGINEERS BUILT BABBAGE-INSPIRED DIFFERENCE ENGINES. FOR SOME REASON, THESE NEVER CAUGHT ON...

YOU DON'T WANT TO COMPUTE x^2+x+41 IN YOUR VERY OWN LIVING ROOM?



...ALTHOUGH DESKTOP ADDING MACHINES AND CASH REGISTERS DID BECOME FIXTURES IN BUSINESS.

ON THE OTHER HAND WERE
THE PUNCHCARD MACHINES,
BEGINNING WITH THE CENSUS
TABULATORS DESIGNED BY

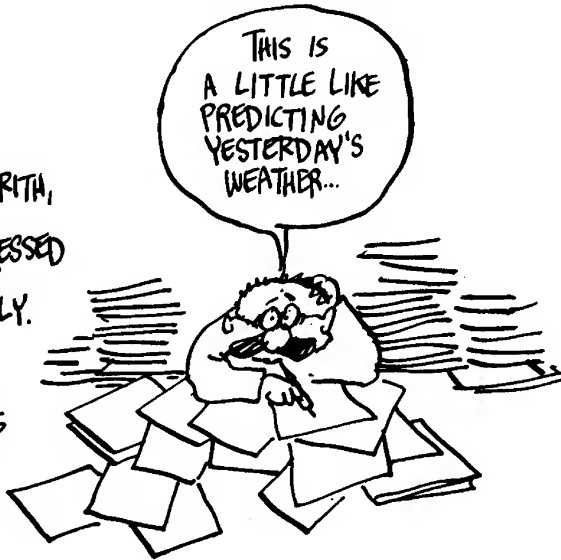
HERMAN HOLLERITH (1860-1929)

INSPIRED, AS BABBAGE HAD BEEN, BY
THE JACQUARD LOOM, HOLLERITH
INVENTED A MACHINE PURELY
FOR ACCUMULATING AND
CLASSIFYING INFORMATION.

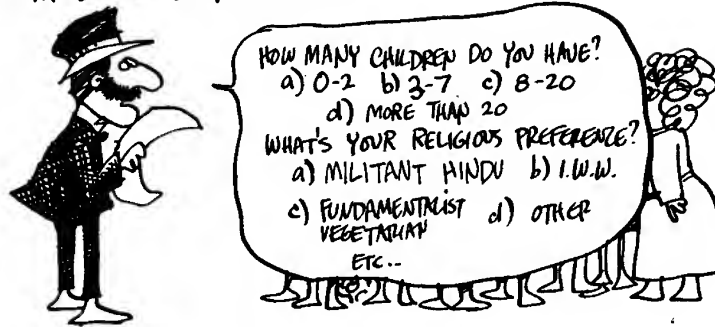


BECAUSE THIS WAS A NEW SORT OF JOB FOR A MACHINE—
AND THE KIND FOR WHICH COMPUTERS ARE IDEALLY SUITED—
LET'S TAKE A CLOSER LOOK.

BEFORE HOLLERITH,
THE CENSUS
BUREAU PROCESSED
ALL DATA BY
HAND... SLOWLY.
THE 1880
CENSUS TOOK
7½ YEARS
TO ANALYZE!



THEN AS NOW, THE CENSUS FORM CONSISTED OF A SERIES OF MULTIPLE CHOICE QUESTIONS...



FROM THIS, ONE WANTED TO FIND:

THE TOTAL NUMBER OF CITIZENS...

HOW MANY HAD 0-2 CHILDREN...

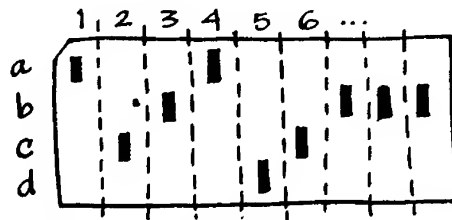
HOW MANY WERE MILITANT HINDUS...
ETC!

AS WELL AS SUCH THINGS AS:



HOW MANY FUNDAMENTALIST VEGETARIANS HAVE MORE THAN 20 CHILDREN?

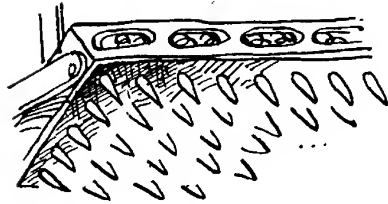
HOLLERITH PROPOSED TO PUT EACH PERSON'S RESPONSES ON A SINGLE PUNCHED CARD THE SIZE OF AN 1880 DOLLAR BILL. TO OVER-



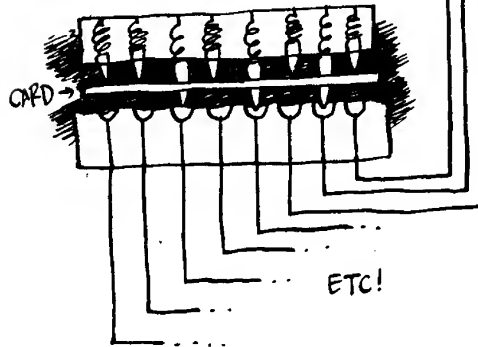
SIMPLIFY SLIGHTLY, EACH COLUMN REPRESENTED ONE QUESTION. THE HOLE IN A GIVEN COLUMN INDICATED THE ANSWER TO THAT QUESTION.

THIS CARD SHOWS RESPONSES OF 1-a, 2-c, 3-b, 4-a, 5-d, ETC...

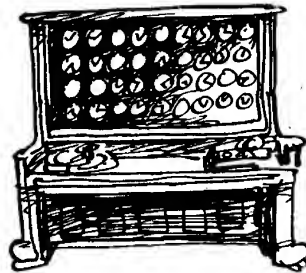
THE CARDS WERE "READ"
BY A DEVICE CONSISTING
OF A GRID OF LITTLE
PINS MOUNTED ON
SPRINGS AND WIRED
ELECTRICALLY.



WHEN BROUGHT INTO
CONTACT WITH THE CARD,
ONLY THOSE PINS LYING OVER
A HOLE WOULD PASS THROUGH.
EACH OF THESE DIPPED
INTO A SMALL CUP OF
MERCURY, COMPLETING AN
ELECTRICAL CIRCUIT.



EACH CUP WAS
WIRED TO A COUNTER,
WHICH ADVANCED
EACH TIME AN ELECTRIC
PULSE ARRIVED.



AND SO THE
RUNNING TOTALS
OF EVERY POSSIBLE
RESPONSE WERE
CONTINUOUSLY
DISPLAYED!



DOES
IT COUNT
UNEMPLOYED
CENSUS
WORKERS?

THE TABULATOR ALSO HELPED ANSWER QUESTIONS SUCH AS: "HOW MANY PEOPLE WHO ANSWERED 2-a ALSO ANSWERED 3-c?"

MEANING: HOW MANY MILITANT HINDUS LIVE IN KANSAS?

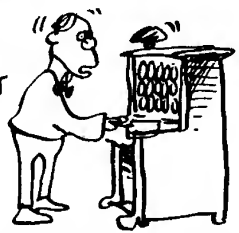
HERE'S HOW:



FIRST ARRANGE A BELL TO RING WHENEVER A CARD WITH 2-a IS ENTERED.



THEY RUN THROUGH ALL CARDS, PULLING OUT ALL THOSE THAT RING THE BELL.



THIS CREATES A STACK OF ALL THE MILITANT HINDU CARDS. RUN THESE THROUGH THE TABULATOR AGAIN.



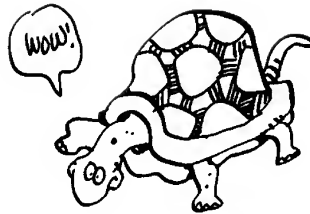
THE MACHINE THEN SHOWS ALL THE TOTALS FOR MILITANT HINDUS.



THIS SORT OF JOB - ANALYZING AND COMPARING LARGE AMOUNTS OF INFORMATION - IS NOW KNOWN AS:

DATA PROCESSING.

THE HOLLERITH TABULATOR CUT THE DATA-PROCESSING TIME FOR THE 1890 CENSUS BY TWO THIRDS, TO 2 1/2 YEARS. THIS SOUNDS LONG NOW, BUT AT THE TIME, IT WAS IMPRESSIVE!!



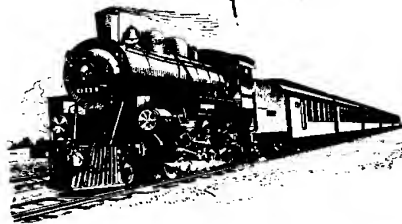
HOLLERITH FOUNDED A COMPANY TO MANUFACTURE HIS CARD-OPERATED DATA PROCESSORS, AND HE FOUND A NUMBER OF TAKERS:

A RAILROAD COMPANY USED THE TABULATOR FOR AUDITING FREIGHT STATISTICS...

A TOOL MANUFACTURER TURNED IT TO COMPILING COSTS, ANALYZING PAYROLL, AND MANAGING INVENTORY...

A WHOLESALE HOUSE NEEDED IT TO KEEP TRACK OF MERCHANDISE, SALES, SALESMEN, CUSTOMERS, ETC ETC ETC...

THIS COMPANY IS GOING SOMEWHERE!



SO HOLLERITH'S COMPANY DID FAIRLY WELL... LATER, IT GOT INTO COMPUTERS, TOO... AND DID WELL... YOU MAY HAVE HEARD OF IT... TODAY IT'S CALLED

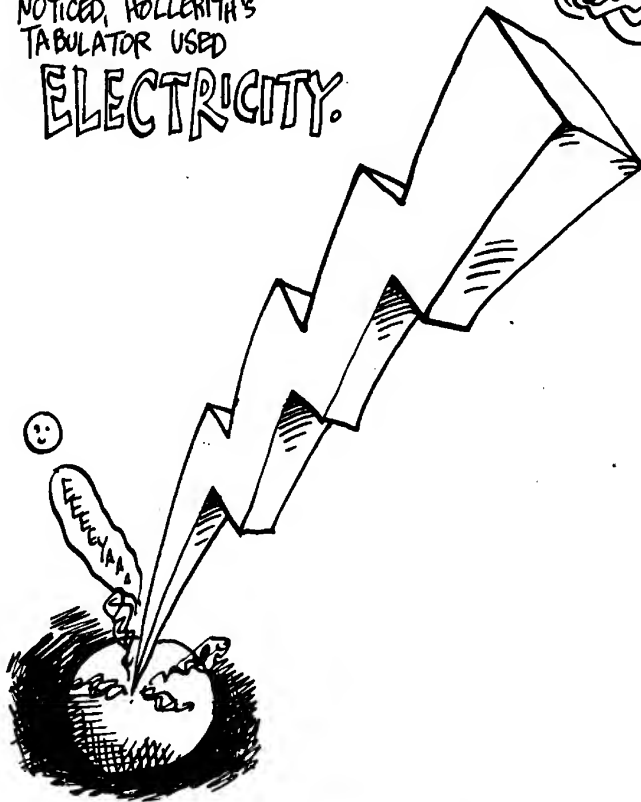
IBM

SO BIG, IT DOESN'T FIT IN THE PANEL!



IN CASE YOU HADN'T
NOTICED, HOLLERITH'S
TABULATOR USED

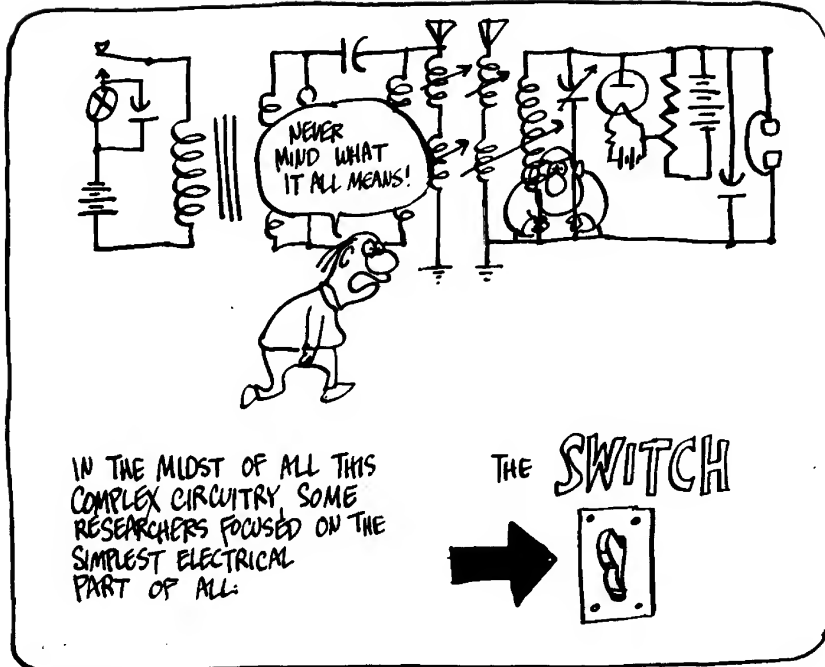
ELECTRICITY.



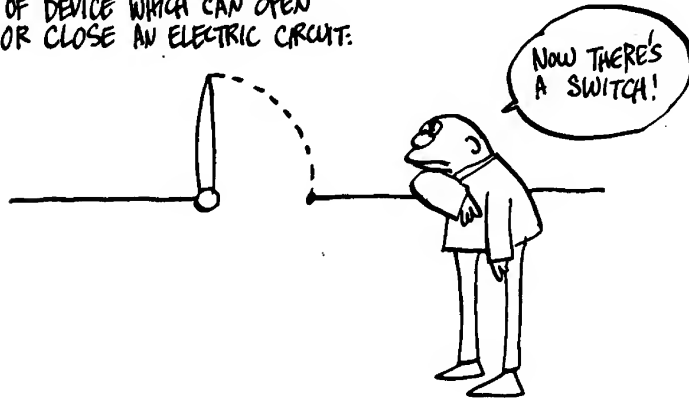
THIS BRINGS US TO THE 20TH CENTURY
AND ITS ELECTRIC MARVELS, RADIO, TELEPHONE,
THE LIGHT BULB, WHICH ALL PLAY A
ROLE IN THE FINAL EPISODES OF
COMPUTER EVOLUTION...

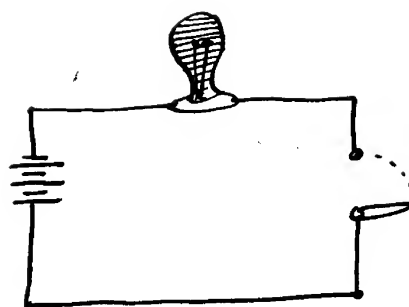
HELLO?
I'D LIKE TO
REPORT MY
THUNDER
STOLEN...





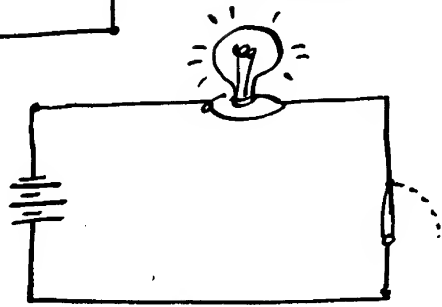
A SWITCH IS ANY KIND OF DEVICE WHICH CAN OPEN OR CLOSE AN ELECTRIC CIRCUIT.



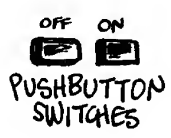
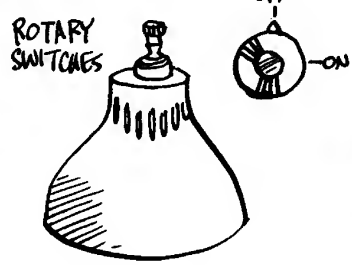
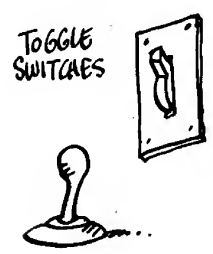


WHEN THE SWITCH IS OPEN, IT BREAKS THE CIRCUIT AND NO CURRENT FLOWS THROUGH THE LIGHT BULB.

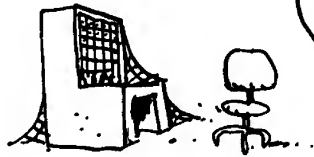
WHEN THE SWITCH IS CLOSED, THE CIRCUIT IS COMPLETE AND THE BULB LIGHTS UP.



A FEW FAMILIAR SWITCHES:



THE TELEPHONE RELAY COULD SWITCH MUCH FASTER THAN THE HUMAN HAND - ABOUT 5 TIMES PER SECOND! IT MADE THE SWITCHBOARD OPERATOR OBSOLETE...



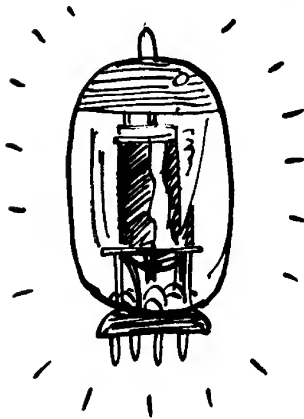
GUESS I'LL APPLY FOR A JOB AT THE RELAY WORKS...



BUT IT COULDN'T HOLD A CANDLE TO ANOTHER TYPE OF SWITCH INVENTED EVEN EARLIER: THE **VACUUM TUBE**.



REMEMBER WHEN TUBES USED TO GLOW IN THE BACK OF THE RADIO? YOU DON'T? SIGH...



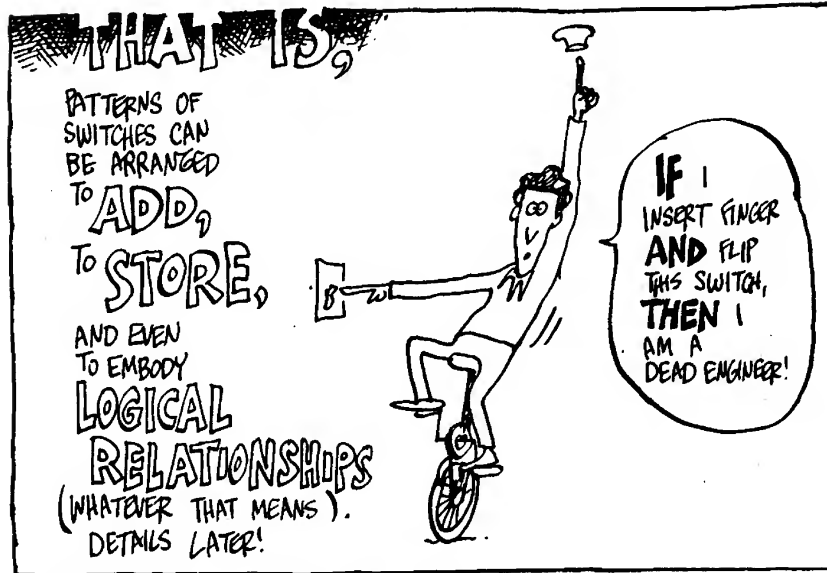
THE TUBE CAN ALSO BE FLIPPED ON AND OFF LIKE A SWITCH, SO FAST YOU CAN'T EVEN SEE IT FLICKER: IT JUST GLOWS... BUT IT CAN SWITCH AS OFTEN AS

1,000,000
TIMES PER SECOND!!!





NOT LONG AFTER THESE SWITCHES WERE INVENTED, PEOPLE REALIZED THAT THEY COULD BE COMBINED INTO COMPUTER COMPONENTS!



BY THE 1930'S, A NUMBER OF PEOPLE HAD SEEN HOW VERY RAPID COMPUTERS COULD BE BUILT FROM HARDWARE STRAIGHT OFF THE SHELF!!





Who built
 THE FIRST ELECTROMECHANICAL COMPUTER? THE VERY FIRST WAS **KONRAD ZUSE** (1910 -). HIS Z-1, BUILT IN 1936, CALCULATED WITH RELAYS AND READ INPUT FROM PUNCHED FILM.

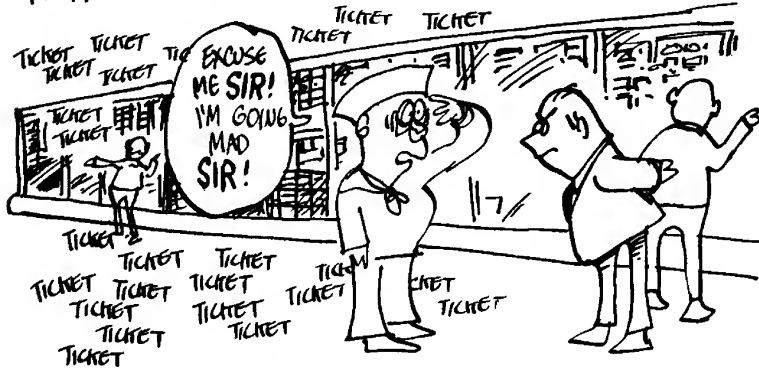




IN THE USA, THE NAVY COLLABORATED WITH HARVARD AND IBM TO CONSTRUCT THE MARK I, AN ELECTROMAGNETIC GIANT LAUNCHED IN 1944.

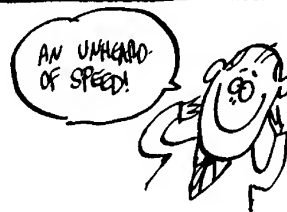


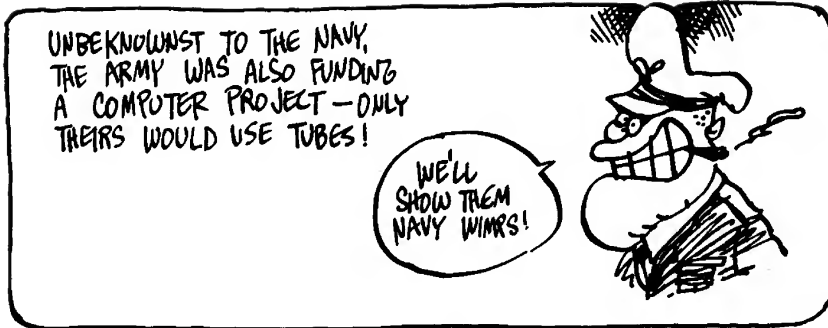
DESIGNED BY HARVARD PROF HOWARD AIKEN, WHO MODELED IT ON BABBAGE'S ANALYTICAL ENGINE, MARK I OCCUPIED SOME 1200 CUBIC FEET AND CONTAINED THOUSANDS OF RELAYS. WHEN IT CRANKED UP, THEY SAY IT SOUNDED LIKE A MILLION KNITTING NEEDLES!!



MARK I COULD MULTIPLY TWO 10-DIGIT NUMBERS (A CONVENIENT MEASURE OF COMPUTER SPEED) IN ABOUT

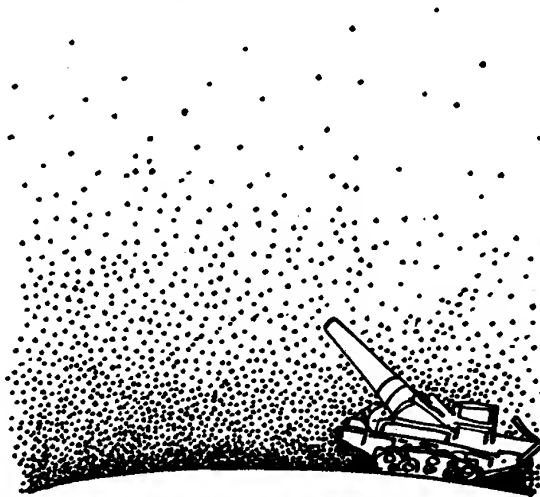
3 SECONDS.





THEIR AIM WAS THE SAME AS TARTAGLIA'S IN THE 1500'S:
TO COMPUTE **BALLISTICS** MORE ACCURATELY.

TARTAGLIA
HAD ERRED
IN SAYING
THAT CANNON-
BALLS FLY IN
PARABOLIC PATHS.
IN REALITY,
AIR RESISTANCE
ALTERS THEIR
TRAJECTORY
APPRECIABLY,
AND IN A
VERY COMPLEX
WAY, BECAUSE
AIR RESISTANCE
DIMINISHES AT
HIGHER
ALTITUDES.



IN WORLD WAR I, THE GERMAN
CANNON "BIG BERTHA" SHOT 94 MILES -
TWICE AS FAR AS EXPECTED FROM
OVERSIMPLIFIED CALCULATIONS!

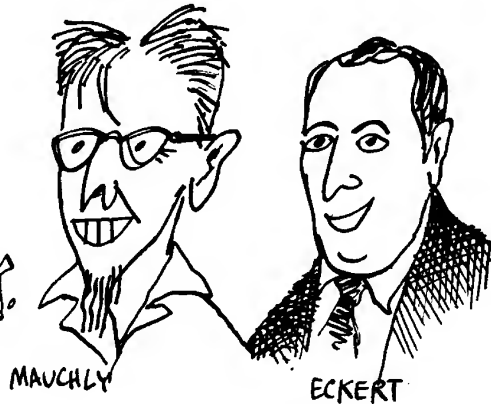
GUNNERS AND BOMBARDIERS THEREFORE NEEDED ACCURATE BALLISTIC TABLES TO AIM BY. THESE COULD HARDLY BE CALCULATED ON THE FLY!

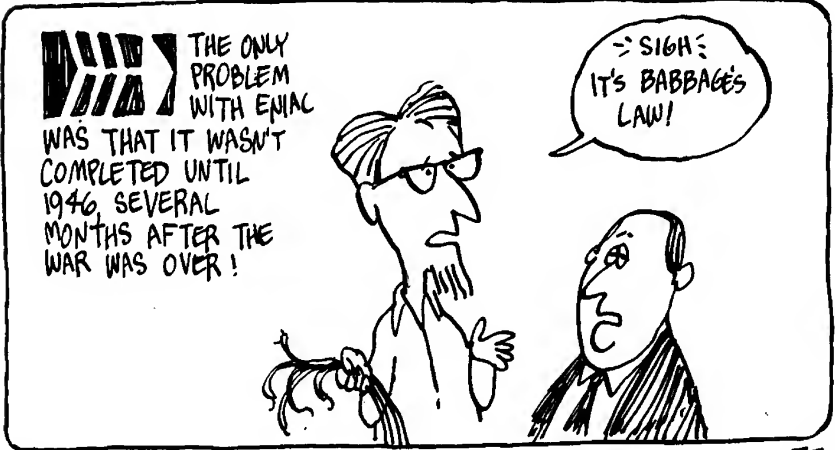
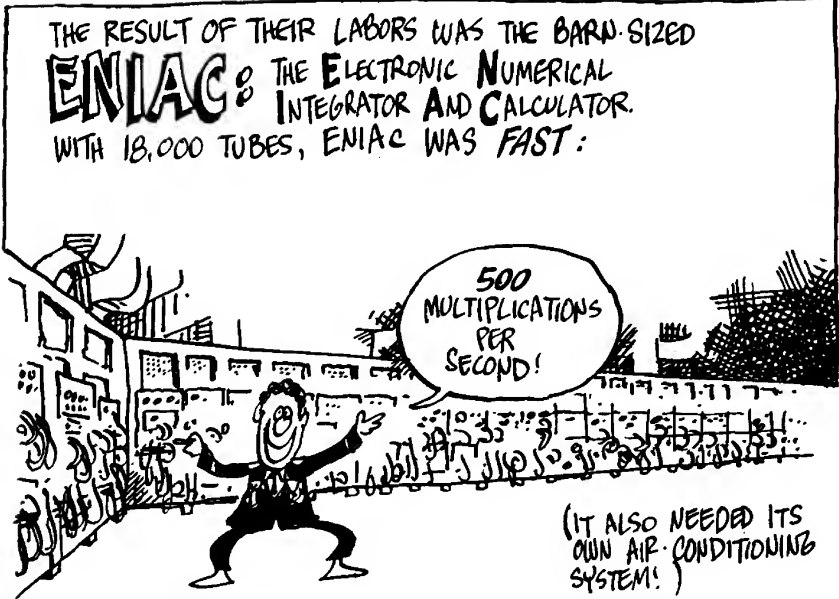


BALLISTIC TABLES USED TO BE CALCULATED BY ROOMFULS OF "GIRLS" WITH ADDING MACHINES — AND EVEN THIS WAS SLOW.



THE CHIEF ENGINEERS IN THE ARMY PROJECT WERE J. PERSPER ECKERT AND JOHN MAUCHLY.







BEST THING ABOUT WAR IS THERE'S ALWAYS ANOTHER!

SO THE ARMY PUT ENIAC TO WORK ON THE NEXT WAR, DOING CALCULATIONS FOR THE NUCLEAR WEAPONS PROGRAM...

ENIAC MAY HAVE BEEN FAST, BUT IN SOME WAYS IT WAS PRETTY DUMB. ITS MEMORY WAS VERY SMALL AND EACH NEW CALCULATION REQUIRED A WHOLE CHANGE OF WIRING.



IN OTHER WORDS, IT'S PROGRAMMED BY SCREWDRIVER!

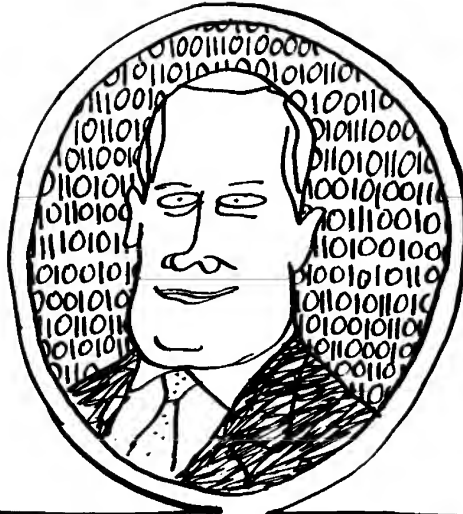
BUT STILL IMPRESSIVE: WITH 18,000 TUBES FLICKERING ON AND OFF 100,000 TIMES PER SECOND, ENIAC HAD TO PERFORM FAR MORE RELIABLY THAN ANY MACHINE EVER CONSTRUCTED.



MISCALCULATION COULD BE FATAL!

NOW ENTERS
**JOHN VON
NEUMANN**

(1903-1957), A
PRINCETON MATH
PROFESSOR WHO
MORE THAN ANYONE
GETS CREDIT FOR
TURNING ELECTRONIC
CALCULATORS
INTO "ELECTRONIC
BRAINS."



VON NEUMANN PONDERED THE COMPUTER'S LOGICAL STRUCTURE
IN THE ABSTRACT: HOW IT CONTROLS ITSELF, HOW MUCH
MEMORY IT NEEDS AND WHAT FOR, ETC... AND HE ASKED
HIMSELF HOW COMPUTERS COULD BE MADE MORE LIKE
HUMAN "WIRING," I.E., THE CENTRAL NERVOUS SYSTEM.



OH, THE
INPUT BONE
CONNECTED TO THE
MEMORY BONE...

CONSIDER HOW A HUMAN BEING "RUNS A PROGRAM":

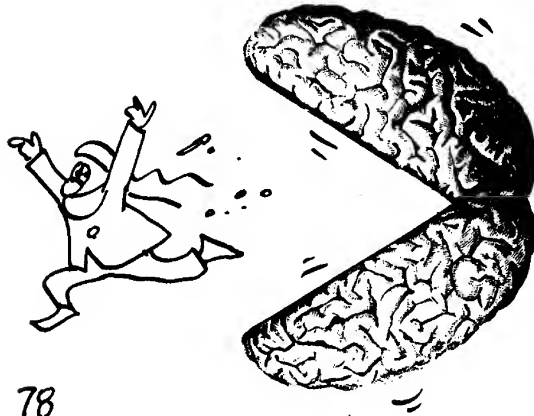
WHEN A SURGEON STARTS TO CUT, IT SHOULDN'T BE NECESSARY TO KEEP REFERRING BACK TO THE TEXTBOOK FOR INSTRUCTIONS.



NO... FIRST THE SURGEON GOES TO MEDICAL SCHOOL, READS THE PROCEDURES, AND COMMITS THEM TO MEMORY.



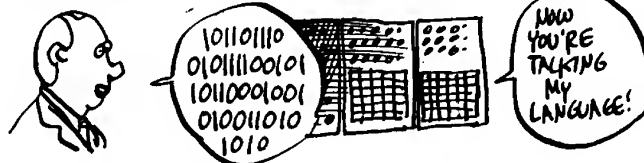
THIS SPEEDS UP SURGERY CONSIDERABLY!



YOUR BRAIN IS FULL OF THESE "STORED PROGRAMS":
YOU KNOW HOW TO TIE YOUR SHOELACES, HOW TO FEED YOURSELF, HOW TO MULTIPLY 94 TIMES 16, HOW TO TALK, HOW TO WALK...

VON NEUMANN PROPOSED TO MAKE COMPUTERS DO LIKEWISE:

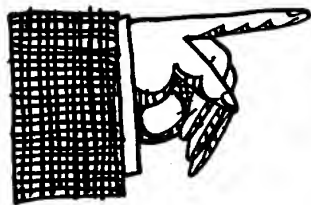
1. FIND A WAY TO **ENCODE** THE INSTRUCTIONS INTO A FORM WHICH COULD BE STORED IN THE COMPUTER'S MEMORY. VON NEUMANN SUGGESTED USING STRINGS OF ONES AND ZEROS.



2. STORE THE INSTRUCTIONS IN MEMORY, ALONG WITH WHATEVER OTHER INFORMATION (NUMBERS, ETC) IS NEEDED TO DO THE PARTICULAR JOB.



3. WHEN RUNNING THE PROGRAM, FETCH THE INSTRUCTIONS STRAIGHT FROM MEMORY, RATHER THAN READING A NEW PUNCHCARD AT EACH STEP.



THIS IS THE CONCEPT OF THE
STORED PROGRAM.

The advantages?



SPEED LIKE THE SURGEON, THE COMPUTER FINDS IT MUCH FASTER TO WHIZ INSTRUCTIONS FROM "BRAIN" TO "FINGERS" THAN TO "RETURN TO THE TEXTBOOK" AFTER EXECUTING EACH STEP.



VERSATILITY:

WITH SEVERAL PROGRAMS STORED AT ONCE, THEY CAN REFER TO ONE ANOTHER RUNNING IN COMBINATION. SURGERY IS ACTUALLY SUCH A COMBINATION.



SELF-MODIFICATION:

IF STORED ELECTRONICALLY, PROGRAMS MAY EASILY BE WRITTEN WHICH CAN MODIFY OR ADJUST THEMSELVES. THIS TURNS OUT TO BE CRITICALLY IMPORTANT!



TO MAKE HIS POINT, VON NEUMANN WROTE SOME CODE FOR A PROGRAM CALLED:

SORT AND MERGE



IT'S A SIMPLE JOB TO DESCRIBE:

GIVEN TWO LISTS OF NAMES (FOR EXAMPLE):

ALABAMA, S.
ANTEATER, J.
ANTEATER, B.
AARDVARK, A.

TARDIGRADE, C.
BEAVER, M.
OWL, H.
ALLIGATOR, A.



AARDVARK, A.
ALABAMA, S.
ALLIGATOR, A.
ANTEATER, B.
ANTEATER, J.
BEAVER, M.
OWL, H.
TARDIGRADE, C.

MAKE ONE LIST IN ALPHABETICAL ORDER.

THIS SEEMINGLY SIMPLE PROCESS BECOMES HORRIBLY TIME-CONSUMING WHEN THE LISTS ARE LONG.

So:
HERE'S ANOTHER IDEAL COMPUTER JOB THAT CONTAINS ESSENTIALLY NO MATH. YOU CAN SEE HOW THIS ONE MIGHT APPEAL TO SOMEONE COMPILING A TELEPHONE DIRECTORY OR A MAILING LIST!!

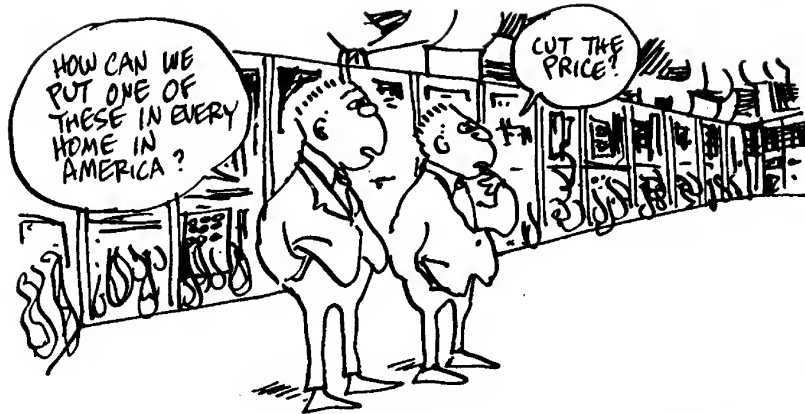


ACTUALLY, THERE'S
SOME ARGUMENT
OVER WHO
INVENTED THE
STORED PROGRAM.
ECKERT AND
MAUCHLY CLAIMED
CREDIT, TOO...
AND THE
ENIAC PROJECT
DISSOLVED IN
A WELTER OF
LAWSUITS OVER
WHO OWNED
WHAT IDEA...



WELCOME
TO THE
COMPUTER
AGE...



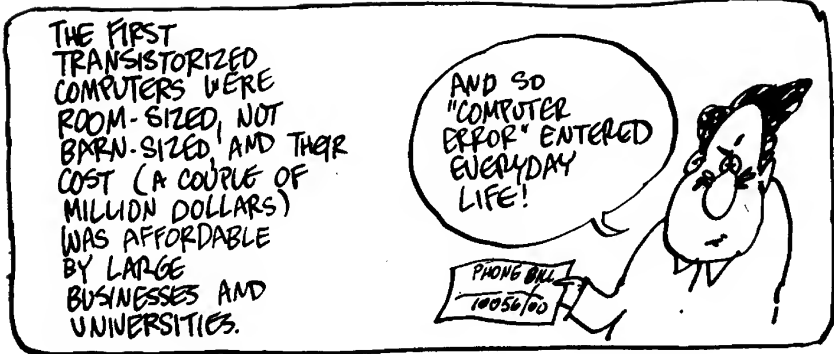


IF COMPUTERS HAD REMAINED AS BULKY AS ENIAC, THEY WOULDN'T BE WHAT THEY ARE TODAY... BUT THEY DIDN'T, AND THEY ARE...

IN 1947, THE YEAR AFTER ENIAC WAS FINISHED, A TEAM AT STANFORD INVENTED THE **TRANSISTOR**, USING ELEMENTS CALLED SEMICONDUCTORS. LIKE TUBES, TRANSISTORS CAN ACT AS SWITCHES, BUT THEY'RE SMALLER, FASTER, COOLER, AND LONGER-LIVED AND THEY DRAW FAR LESS ELECTRIC POWER.

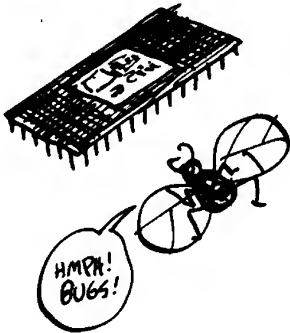


TIME TO RETIRE, FATSO!



THEN THE TRANSISTOR BEGAN TO SHOW AN INCREDIBLE ABILITY TO SHRINK IN SIZE AND PRICE.

FIRST CAME
**INTEGRATED
 CIRCUITS** —
 A WHOLE BOARDFUL
 OF TRANSISTORS
 MANUFACTURED AS
 A SINGLE UNIT... THEN
**LARGE-SCALE AND
 VERY LARGE-SCALE
 INTEGRATION**
 (LSI AND VLSI), WHICH
 PACKED HUNDREDS OF
 THOUSANDS OF
 TRANSISTORS ON A
 TINY CHIP!



AS COMPONENTS SHRANK, THE INDUSTRY EXPLODED!

IN THE '60'S, THE
MINICOMPUTER
APPEARED. IT WAS THE
SIZE OF A DESK!



MAKES IT
LESS MYSTERIOUS
SOMEHOW!



IN THE '70'S CAME THE
MICRO, WHICH CAN BE
AS SMALL AS YOU LIKE.



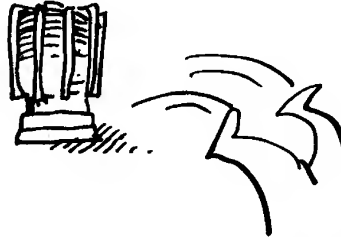
WHAT'S
NEXT?
THE
DISPOSABLE?

BY THIS TIME, BIG COMPUTERS,
ALSO KNOWN AS
MAINFRAMES, HAD
BECOME IMMENSELY POWERFUL.



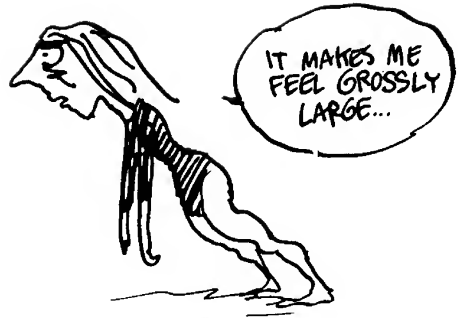
100,000
TRANSISTORS
PER CHIP...
100,000
CHIPS PER
MACHINE...

AND FINALLY THE EXOTIC
SUPERCOMPUTERS,
WHICH CALCULATE AT RATES
UP TO 500 MEGAFLOPS* —
A MILLION TIMES FASTER THAN
ENIAC!

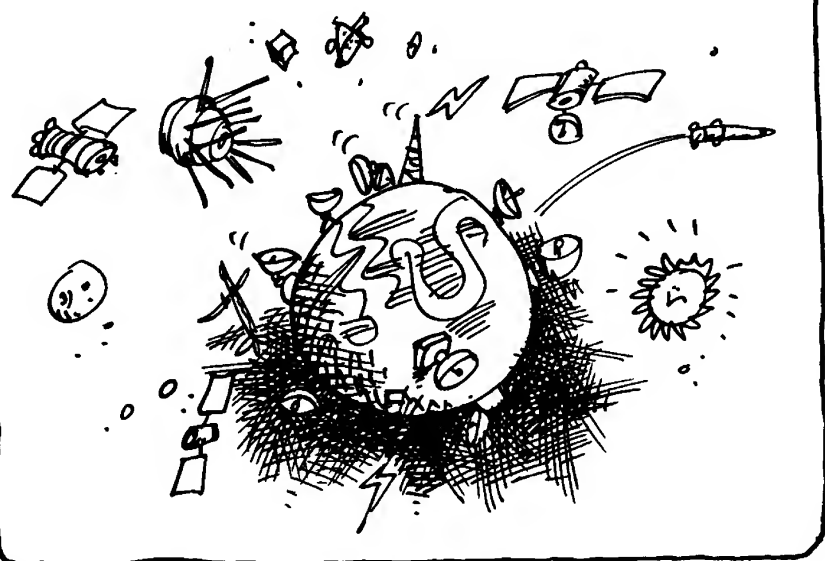


* MILLION FLOATING POINT OPERATIONS
PER SECOND.

THERE'S NO END IN SIGHT... NOW WE HAVE MICROS WITH THE POWER OF MINIS, "SUPERMINIS" THAT RIVAL MAINFRAMES, MINIS ON A CHIP... AND THERE'S TALK OF REDUCING COMPONENTS TO MOLECULAR SIZE USING RECOMBINANT DNA TECHNOLOGY...



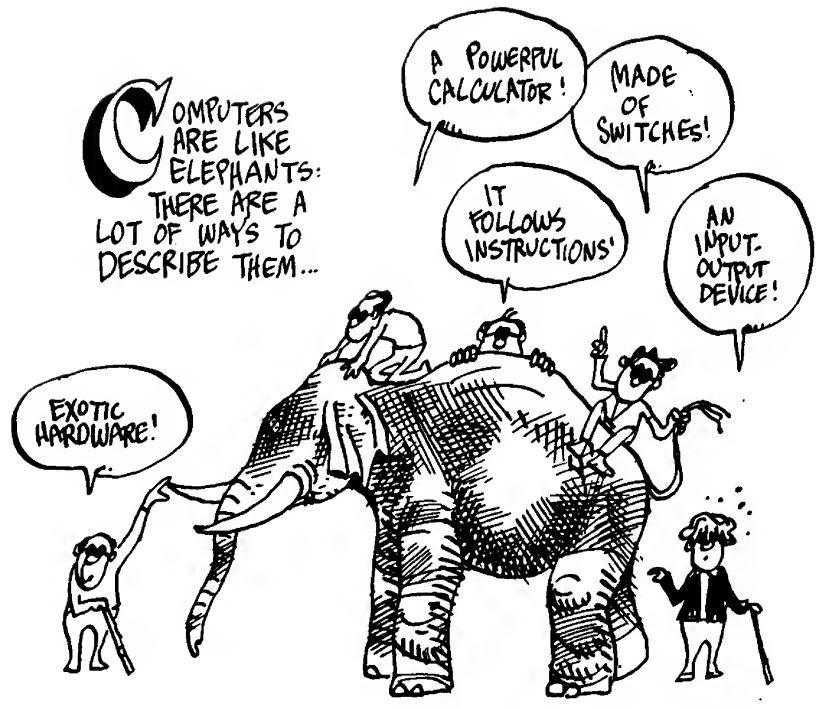
THERE SEEMS TO BE NO SUCH THING AS A COMPUTER WITH TOO MUCH COMPUTING POWER. NO MATTER THE SPEED OR CAPACITY, COMPUTERS ALWAYS FIND JOBS TO DO... AND NO WONDER: THIS IS THE AGE OF EXCESS INFORMATION!



PART II

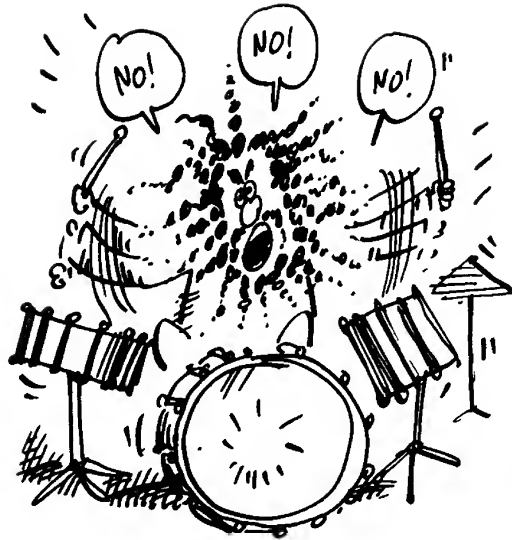
LOGICAL SPAGHETTI





.....
HOW DOES ONE GET TO THE HEART OF THE MATTER?





IF THERE'S ONE
IDEA WE'VE TRIED
TO DRUM IN,
IT'S THAT THE
COMPUTER IS
ESSENTIALLY AN
**INFORMATION
PROCESSOR.**
SO FORGET THE
ELEPHANT...

TO UNDERSTAND INFORMATION PROCESSING, IT HELPS TO
COMPARE IT WITH A MORE FAMILIAR PROCESS: **COOKING.**
SO STEP INTO GRANDMOTHER BABBAGE'S KITCHEN, AS SHE
PREPARES BASIC SPAGHETTI...



HERE'S THE WORLD FAMOUS RECIPE:

1 BRING A KETTLE OF SALTED WATER TO BOIL.



2 ADD 8 OZ. OF RAW SPAGHETTI.



3 BOIL FOR 10 MINUTES.



4 DRAIN THROUGH A SIEVE.



5 SERVE..



THIS SPAGHETTI IS BETTER ANALYZED THAN EATEN!

IT'S NOT HARD TO DISTINGUISH A FEW COMPONENTS IN THIS PROCESS:

FIRST, THE INGREDIENTS, OR **INPUT.**

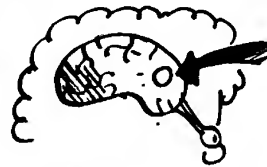


NEXT, THE EQUIPMENT WHICH DOES THE COOKING: HANDS, KETTLE, STOVE, SALTSHAKER, SIEVE, PLATE, SPOON.



THESE FORM THE **PROCESSING UNIT.**

LESS OBVIOUSLY, THERE IS A PART OF THE COOK'S BRAIN WHICH CONTROLS THE PROCESS. IT MONITORS AND DIRECTS THE STEP-BY-STEP UNFOLDING OF THE RECIPE. THIS IS REFERRED TO AS THE



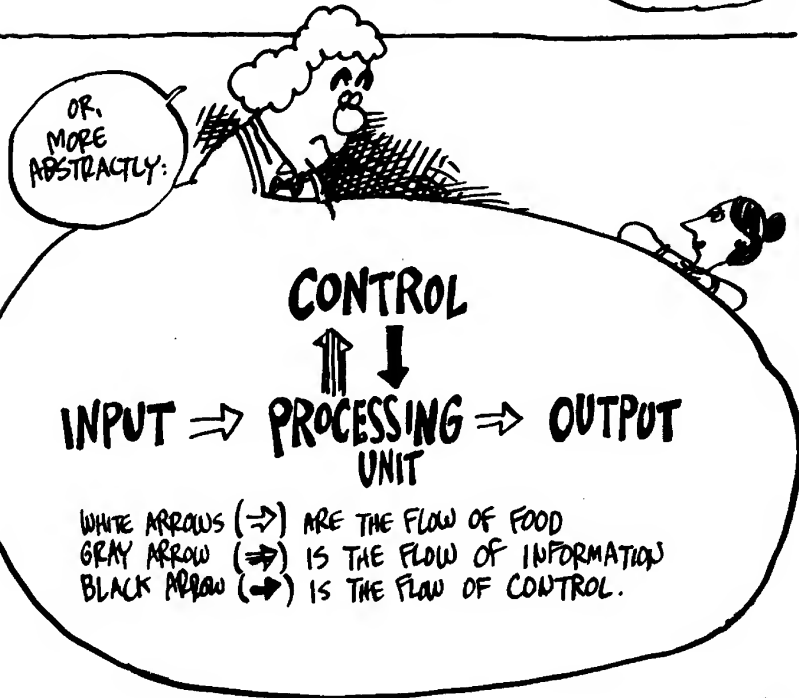
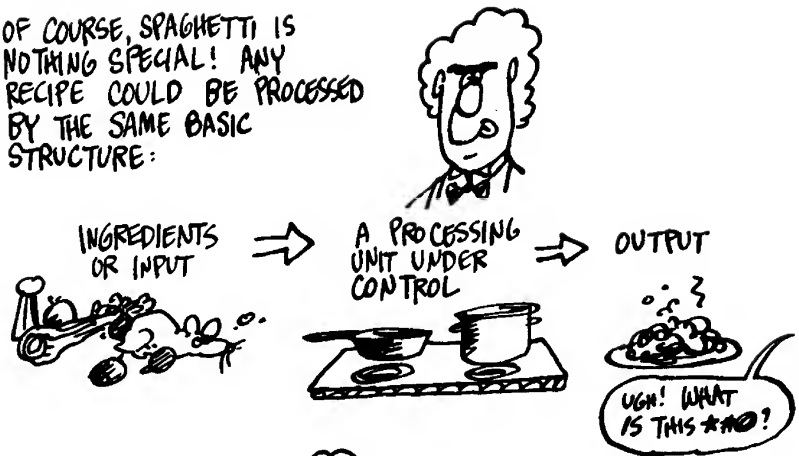
CONTROL UNIT.

AND OF COURSE THE COMPLETED DISH, OR **OUTPUT.**



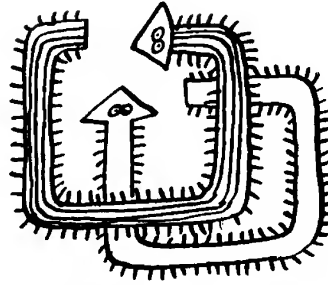
WHICH ALSO RESEMBLES THE COOK'S BRAIN...

OF COURSE, SPAGHETTI IS NOTHING SPECIAL! ANY RECIPE COULD BE PROCESSED BY THE SAME BASIC STRUCTURE:

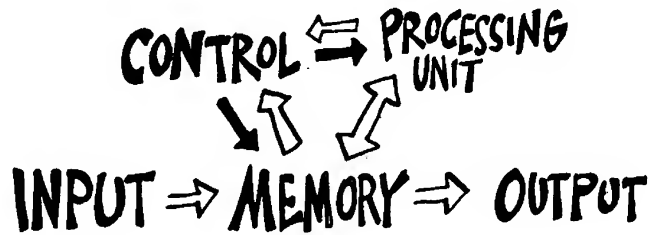


WITH COMPUTERS, THE DIAGRAM IS SLIGHTLY DIFFERENT:

THERE ARE TWO REASONS FOR THIS: ONE IS THE FACT THAT INPUT AND OUTPUT ARE INFORMATION, NOT FOOD — SO THE GRAY ARROW IS THE SAME AS THE WHITE ONES.



THE OTHER IS THE GREAT IMPORTANCE OF ~~MEMORY~~ **MEMORY**, WHICH FORMS THE FIFTH AND FINAL COMPONENT. IN COMPUTERS, ALL INFORMATION PASSES INTO MEMORY FIRST! HERE'S THE DIAGRAM:



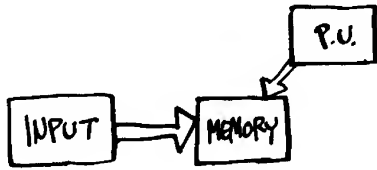
⇒ = INFORMATION FLOW

⇔ = CONTROL FLOW

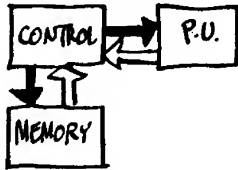


IN THE CASE OF COMPUTERS, THE **INPUT** CONSISTS OF ALL THE "RAW" DATA TO BE PROCESSED — AS WELL AS THE ENTIRE "RECIPE," OR PROGRAM, WHICH SPECIFIES WHAT'S TO BE DONE WITH THEM.

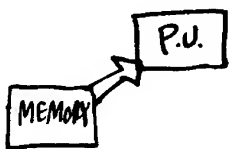
THE **MEMORY** STORES THE INPUT AND RESULTS FROM THE PROCESSING UNIT:



CONTROL READS THE PROGRAM AND TRANSLATES IT INTO A SEQUENCE OF MACHINE OPERATIONS.



THE **PROCESSING UNIT** PERFORMS THE ACTUAL ADDITIONS, MULTIPLICATION, COUNTING, COMPARISON, ETC, ON INFORMATION RECEIVED FROM MEMORY.

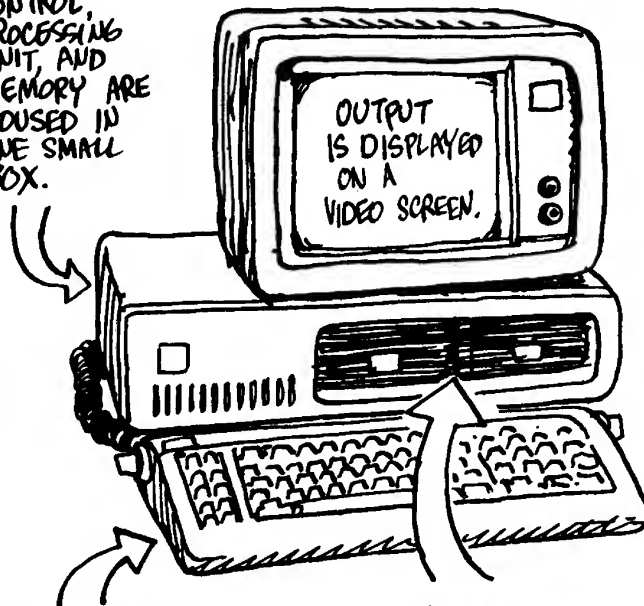


THE **OUTPUT** CONSISTS OF THE PROCESSING UNIT'S RESULTS, STORED IN MEMORY AND TRANSMITTED TO AN OUTPUT DEVICE.



HERE'S THE REAL THING (AN IBM PERSONAL COMPUTER),
JUST TO GIVE ONE EXAMPLE OF HOW THESE COMPONENTS
MAY ACTUALLY LOOK:

CONTROL,
PROCESSING
UNIT, AND
MEMORY ARE
HOUSED IN
ONE SMALL
BOX.



INPUT IS ENTERED
FROM KEYBOARD.

DISK DRIVES
PROVIDE EXTRA
MEMORY STORAGE

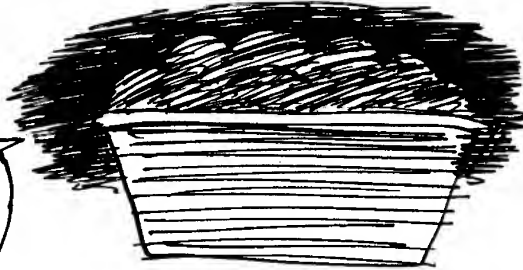
OTHER COMMON INPUT/OUTPUT DEVICES (NOT PICTURED) ARE
A **MODEM**, FOR SENDING AND RECEIVING SIGNALS
OVER THE PHONE, AND A **PRINTER**, FOR PRODUCING
OUTPUT ON PAPER.

LET'S START IN THE MIDDLE, WITH THE

PROCESSING UNIT:

IN THE KITCHEN, A CHEF MAY DISPLAY A RICH REPERTOIRE OF PROCESSING POSSIBILITIES:

BRASE
BROIL
SAUTÉ
ROAST
POACH
STEAM
BOIL
FRY
BAKE...



BUT, AS THE GREAT ESCOFFIER HIMSELF HAS REMARKED, ALL COOKING TECHNIQUES ARE COMBINATIONS OF SIMPLER STEPS: THE APPLICATION OF MORE OR LESS HEAT, WET OR DRY, ETC...

THESE FEW
ARE
ELEMENTARY!



LIKEWISE, ALL THE POWER OF THE COMPUTER DEPENDS ON A COUPLE OF ELEMENTARY OPERATIONS.



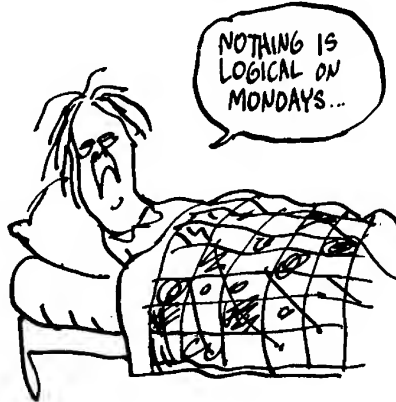
O.K... O.K... NO MORE
BEATING AROUND THE
BUSH WITH CULINARY
METAPHORS...

THE COMPUTER'S ELEMENTARY OPERATIONS ARE

 **LOGICAL** 



WHAT'S A LOGICAL
OPERATION, YOU ASK?
A LOGICAL QUESTION,
CONSIDERING HOW MUCH
EASIER IT IS TO THINK
OF ILLOGICAL OPERATIONS,
LIKE AMPUTATION OF
THE THUMBS OR
GETTING OUT OF BED
ON MONDAYS...



TO EVERYONE'S GOOD FORTUNE, LOGIC ISN'T AS HARD AS IT USED TO BE. IN ARISTOTLE'S TIME, THE SUBJECT WAS DIVIDED INTO INDUCTIVE AND DEDUCTIVE BRANCHES, INDUCTIVE LOGIC BEING THE ART OF INFERRING TRUTHS BY OBSERVING NATURE, WHILE DEDUCTIVE LOGIC DEDUCED TRUTHS FROM OTHER TRUTHS:

1. YOU ARE A MAN.
2. ALL MEN ARE MORTAL.
3. THEREFORE, YOU ARE MORTAL.

DEDUCTIVE 3

INDUCTIVE
:AHEM:
HOW DO YOU KNOW ALL MEN ARE MORTAL??



MEDIEVAL

LOGICIANS COMPOUNDED THE CONFUSION WITH SIX "MODES": A STATEMENT WAS EITHER TRUE, FALSE, NECESSARY, CONTINGENT, POSSIBLE, OR IMPOSSIBLE.



NECESSARY IS TO CONTINGENT AS TRUE IS TO FALSE... POSSIBLY..

THEIR REASONING GREW SO MINDLESS THAT THE MEDIEVAL LOGICIAN DUNS SCOTUS HAS BEEN IMMORTALIZED IN THE WORD "DUNCE"!

THE SUBJECT WAS STRETCHED
TO ABSURD LENGTHS
BY LEWIS
CARROLL:

- (1) GENTILES
HAVE NO
OBJECTION
TO PORK.
- (2) NOBODY WHO
ADMIRES PIGSTIES
EVER READS
HOGG'S POEMS.
- (3) NO
MANDARIN
KNOWS HEBREW.
- (4) EVERYONE, WHO
DOES NOT OBJECT
TO PORK, ADMIRE
TURNSTILES.
- (5) NO JEW IS
IGNORANT
OF HEBREW.

THEREFORE, NO
MANDARIN EVER
READS HOGG'S
POEMS. " *

CLEARLY, IT WAS
TIME TO SIMPLIFY
THE SUBJECT...

* FROM SYMBOLIC LOGIC

THIS STEP WAS
TAKEN BY
**GEORGE
BOOLE** (1815-
1864),
AN ENGLISH MATHEMATICIAN
WHO BUILT AN
"ALGEBRA" OUT OF
LOGIC.



THAT IS, HE MADE
LOGIC FULLY
SYMBOLIC, JUST
LIKE MATH. SENTENCES
WERE DENOTED BY LETTERS
AND CONNECTED BY
ALGEBRAIC SYMBOLS — AN
IDEA GOING BACK TO
LEIBNIZ, WHO HAD DREAMED
OF "JUSTICE BY ALGEBRA."



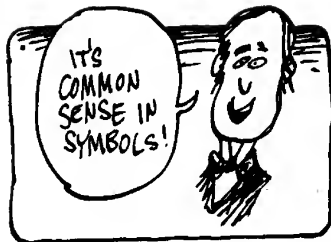
$(1-x) \cdot (1-y) =$
 $1-x-y+xy.$
THEREFORE, 30 YEARS!

WE CAN'T POSSIBLY DESCRIBE BOOLE'S ALGEBRA IN ITS
ENTIRETY. WE'LL LIMIT OURSELVES TO THREE WORDS:

**AND,
OR,
NOT!**



BOOLE LOOKED AT
THE VERY
CONNECTIVE TISSUE
OF LANGUAGE:
THE WORDS "AND",
"OR", AND "NOT".



SUPPOSE P IS ANY STATEMENT... FOR EXAMPLE,

P = "The pig has spots."

ACCORDING TO BOOLE, THIS SENTENCE IS EITHER TRUE (T) OR FALSE (F). NO OTHER OPTION IS ALLOWED! *



NOW LET Q BE ANOTHER STATEMENT - LIKEWISE TRUE OR FALSE:

Q = "The pig is glad."

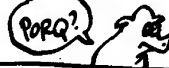


NOW FORM THE COMPOUND SENTENCES:

P AND Q = THE PIG IS SPOTTED AND THE PIG IS GLAD.

P OR Q = THE PIG IS SPOTTED OR THE PIG IS GLAD.

WHEN ARE THESE SENTENCES TRUE?



* IN SOME VERSIONS OF LOGIC, MORE THAN TWO TRUTH VALUES ARE PERMISSIBLE.

THERE ARE FOUR POSSIBLE COMBINATIONS OF TRUTH AND FALSEHOOD FOR P AND Q.



P TRUE, Q TRUE



P FALSE, Q TRUE



P TRUE, Q FALSE



P FALSE, Q FALSE

AND

"THE PIG IS GLAD AND HAS SPOTS."

THIS IS TRUE ONLY IN THE ONE CASE IN WHICH P, Q ARE BOTH TRUE. THIS IS SUMMARIZED IN A TRUTH TABLE:



P	Q	P AND Q
T	T	T
T	F	F
F	T	F
F	F	F

OR

"THE PIG IS GLAD OR HAS SPOTS."

THIS IS TRUE IN THE THREE CASES FOR WHICH EITHER ONE OF THE STATEMENTS P, Q IS TRUE.



P	Q	P OR Q
T	T	T
T	F	T
F	T	T
F	F	F

EXCEPT FOR THE ONE WEIRD EQUATION $1 \oplus 1 = 1$, THESE LOOK LIKE ORDINARY ARITHMETIC... WITH "AND" PLAYING THE ROLE OF "TIMES" AND "OR" IN THE ROLE OF "PLUS."



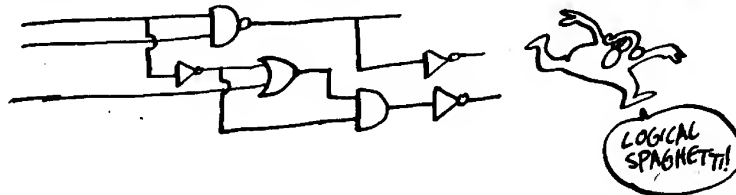
WE'RE NEVER GOING TO USE THE SYMBOLS \cdot AND \oplus ... YOU CAN FORGET ABOUT THEM... BUT USING 1 AND 0 TO REPRESENT TRUE AND FALSE IS VERY USEFUL... SO FROM NOW ON WE'LL WRITE TRUTH TABLES LIKE THIS:

P	Q	P AND Q
1	1	1
1	0	0
0	1	0
0	0	0

P	Q	P OR Q
1	1	1
1	0	1
0	1	1
0	0	0

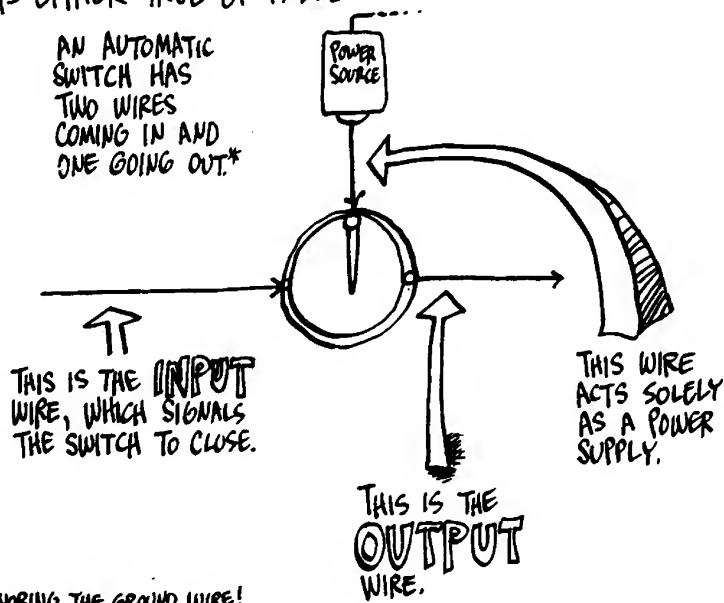
P	NOT-P
1	0
0	1

FROM THESE RELATIONSHIPS, BOOLE BUILT UP AN ENTIRE ALGEBRA, USING ONLY THE NUMBERS 0 AND 1... TODAY THIS **BOOLEAN ALGEBRA** IS USED ALL THE TIME BY COMPUTER ENGINEERS — ONLY THEY EXPRESS IT AS ELECTRICAL CIRCUITS...



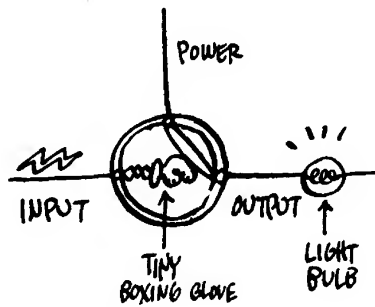
THE KEY IS THE **AUTOMATIC SWITCH**, WHICH IS EITHER OPEN OR CLOSED, AS A LOGICAL PROPOSITION IS EITHER TRUE OR FALSE.

AN AUTOMATIC SWITCH HAS TWO WIRES COMING IN AND ONE GOING OUT.*

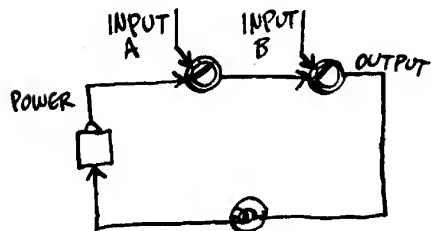


*IGNORING THE GROUND WIRE!

WHEN NO CURRENT FLOWS THROUGH THE INPUT WIRE, THE SWITCH REMAINS OPEN, AS PICTURED ABOVE. WHEN AN INPUT SIGNAL ARRIVES, HOWEVER, THE ELECTRONIC EQUIVALENT OF A MINIATURE BOXING GLOVE "PUNCHES" THE SWITCH CLOSED, RESULTING IN AN OUTPUT SIGNAL.



WHAT IS THE OUTPUT WHEN TWO SWITCHES (A, B) ARE ARRANGED IN SERIES, ONE AFTER THE OTHER? [IN OUR DIAGRAM, PLEASE NOTE THE REARRANGEMENT OF WIRES, MADE FOR CONVENIENCE OF ILLUSTRATION.]



THE CURRENT CAN FLOW ONLY IF BOTH SWITCHES ARE CLOSED— I.E., WHEN INPUT SIGNALS ARRIVE SIMULTANEOUSLY AT A AND B.

WRITING 1 FOR CURRENT AND 0 FOR NO CURRENT, WE CAN THEN WRITE THIS INPUT-OUTPUT TABLE. LOOK FAMILIAR? IT SHOULD! IT'S IDENTICAL TO THE TRUTH TABLE FOR AND!

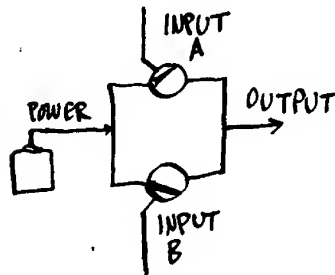
A	B	OUTPUT
1	1	1
1	0	0
0	1	0
0	0	0

THAT'S WHY THIS ARRANGEMENT OF SWITCHES IS CALLED AN

AND-GATE

AND IT HAS ITS VERY OWN SYMBOL →

TWO SWITCHES CONNECTED IN PARALLEL BEHAVE LIKE LOGICAL **OR**: CURRENT CAN PASS FROM POWER TO OUTPUT IF EITHER SWITCH A, B IS CLOSED (OR IF BOTH ARE).



A	B	OUTPUT
1	1	1
1	0	1
0	1	1
0	0	0

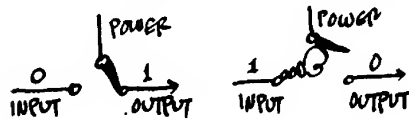
THIS IS THE

OR-GATE

AND ITS SYMBOL IS:



NOT IS NOT ANY MORE DIFFICULT... IT USES A SPECIAL SWITCH THAT REMAINS CLOSED UNTIL AN INPUT SIGNAL OPENS IT — JUST THE REVERSE OF AN ORDINARY SWITCH:

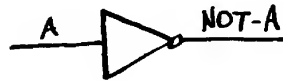


A	OUTPUT
1	0
0	1

THIS KIND OF SWITCH IS CALLED

AN **INVERTER**,

AND IT HAS A SYMBOL, TOO:

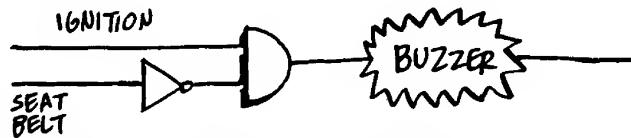


AN EVERYDAY EXAMPLE SHOWS HOW THESE SIMPLE GATES CAN MAKE LOGICAL DECISIONS.

YOU KNOW THOSE BUZZERS THAT GO OFF WHEN YOU START YOUR CAR AND YOUR SEAT BELT ISN'T FASTENED? THE KIND THAT'S SPECIALLY DESIGNED TO PENETRATE HUMAN BONE?



WELL, THAT'S BECAUSE THE SEAT BELT AND IGNITION ARE CONNECTED BY AN AND-GATE, LIKE SO:



THAT IS, **IF** THE IGNITION IS ON **AND** THE SEAT BELT IS **NOT**, THE BUZZER SOUNDS! PRETTY LOGICAL, NO?

CAN YOU THINK OF ANY EXAMPLES OF OR-GATES IN DAILY LIFE?

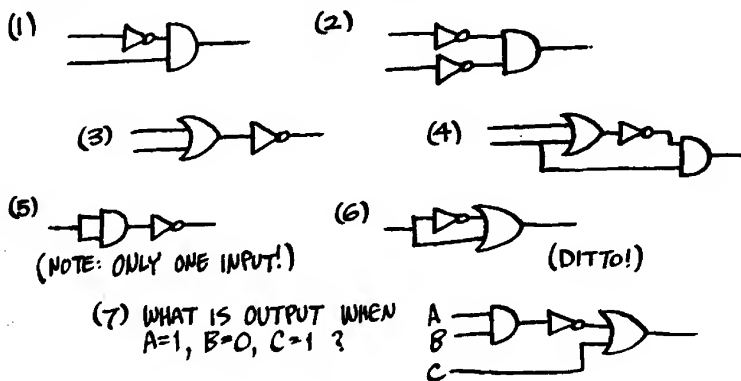


(HOW ABOUT A SMOKE ALARM TRIGGERED BY EITHER OF TWO DIFFERENT DETECTORS?)

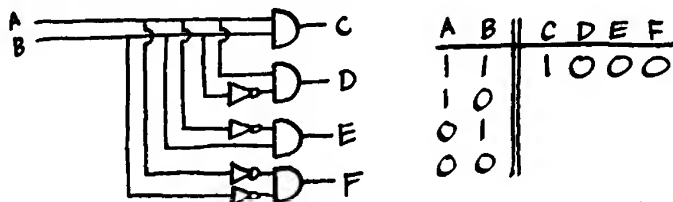


HERE ARE A FEW WARM-UP EXERCISES FOR CHASING THROUGH LOGIC DIAGRAMS:

DO THE INPUT-OUTPUT (I/O) TABLES:



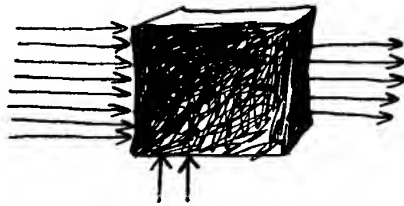
(8) COMPLETE THE I/O TABLE:



DESIGN LOGIC DIAGRAMS WITH THESE I/O TABLES.

(9)		(10)		(11)		(12)	
IN	OUT	IN	OUT	IN	OUT	IN	OUT
1	1	1	0	1	1	1	0
1	0	1	0	0	1	1	0
0	1	0	1	0	0	0	1
0	0	0	0	0	1	0	0

LOGIC GATES HAVE ONLY ONE OR TWO INPUTS AND A SINGLE OUTPUT — BUT COMPUTER COMPONENTS HAVE MANY INPUTS AND OUTPUTS WITH COMPLICATED INPUT/OUTPUT BEHAVIOR:



THE WONDERFUL FACT IS THAT **ANY** INPUT/OUTPUT TABLE CAN BE PRODUCED BY A COMBINATION OF LOGIC GATES!

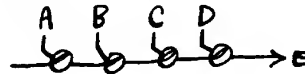
TO DO IT, YOU NEED MULTIPLE-INPUT LOGIC GATES.
HERE'S A 4-INPUT AND-GATE:



A	B	C	D	E
1	1	1	1	1
1	1	1	0	0
1	1	0	1	0
1	0	1	1	0
0	0	0	0	0

} ALL 0'S

THIS MEANS
E=1 IF A=B=C=D=1,
AND E=0 OTHERWISE.
THE GATE CAN BE MADE
WITH FOUR SWITCHES IN
SERIES:



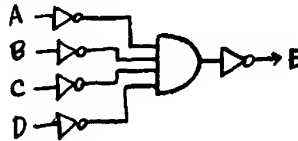
SIMILARLY, THERE'S A MULTIPLE-INPUT OR-GATE:



A	B	C	D	E
1	1	1	1	1
1	1	1	0	1
1	1	0	1	1
1	0	1	1	1
0	0	0	1	1
0	0	0	0	0

} ALL 1'S

IT CAN ACTUALLY BE MADE
FROM AN AND-GATE AND
SOME INVERTERS:



AS AN EXAMPLE OF HOW TO PRODUCE A GIVEN INPUT/OUTPUT TABLE, LET'S SOLVE PROBLEM #12:

IN		OUT
A	B	C
1	1	0
1	0	1
0	1	1
0	0	0

BEGIN BY FINDING ALL ROWS WHERE C=1.

THE TABLE SAYS C=1 IF A=1 AND B=0 OR A=0 AND B=1.
C=0 OTHERWISE.

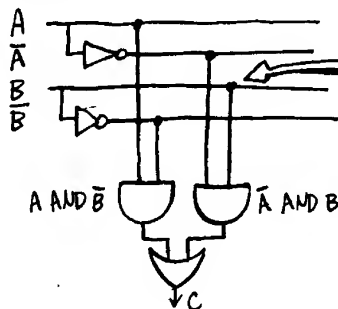
WRITING \bar{A} FOR NOT-A, THIS AMOUNTS TO SAYING

C=1 IF A=1 AND \bar{B} =1 OR \bar{A} =1 AND B=1.
C=0 OTHERWISE.

IN OTHER WORDS,

$$C = (A \text{ AND } \bar{B}) \text{ OR } (\bar{A} \text{ AND } B)$$

TO DRAW THE CIRCUIT, RUN THE INPUT WIRES AND THEIR NEGATIVES IN ONE DIRECTION —



—AND ATTACH THE GATES TO THE APPROPRIATE WIRES.

EXACTLY THE SAME METHOD WORKS FOR MORE INPUTS.
FOR EXAMPLE:

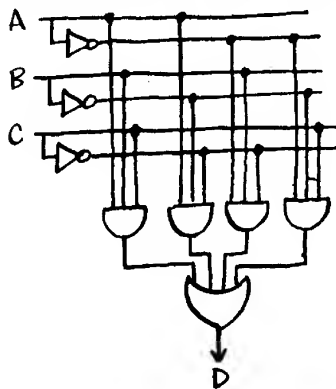
A	B	C	D
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

AGAIN, FIND ALL ROWS WITH OUTPUT = 1.

NOTE ALL POSSIBLE INPUT COMBINATIONS!

IN THIS CASE,
 $D = (A \text{ AND } B \text{ AND } C) \text{ OR } (A \text{ AND } \bar{B} \text{ AND } \bar{C}) \text{ OR } (\bar{A} \text{ AND } B \text{ AND } \bar{C}) \text{ OR } (\bar{A} \text{ AND } \bar{B} \text{ AND } C).$

RUN THE INPUTS AND THEIR NEGATIVES ACROSS THE PAGE, ATTACH AND-GATES, THEN RUN THEM THROUGH AN OR-GATE!



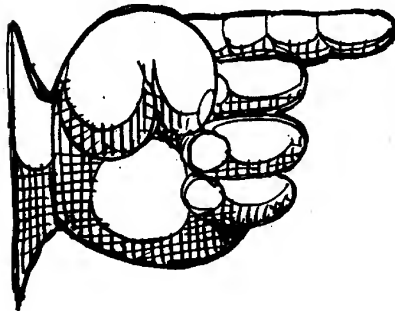
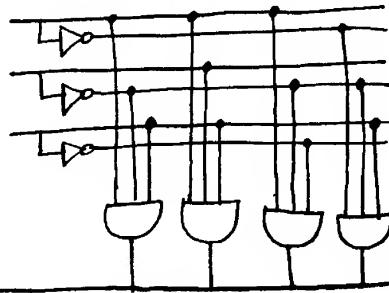
TO REPEAT: BY THE SAME METHOD, YOU CAN PRODUCE INPUT/OUTPUT TABLE!!

WHERE DOES THIS
LEAVE US?



WITH ONE
FOOT IN THE
SWAMP!

BY NOW YOU MAY
BE GETTING THE IDEA
THAT INFORMATION IS
ENCODED INSIDE COMPUTERS
AS STRINGS OF 1'S
AND 0'S, WHICH CAN BE
TRANSFORMED IN ANY
WAY WE LIKE BY THE
RIGHT COMBINATION OF
LOGIC GATES.



BUT WE
HAVEN'T
REALLY SEEN
HOW LOGIC
GATES CAN
DO THE
JOB
COMPUTERS
WERE
DESIGNED
FOR:

NAMELY:
HOW DO COMPUTERS
COMPUTE?



The questions:

□ IS THERE SOME NATURAL WAY TO REPRESENT NUMBERS USING ONLY 0'S AND 1'S? CAN THE OPERATIONS OF ARITHMETIC BE BUILT OUT OF LOGIC?

The answer

(WHICH GOES BACK TO OUR OLD PAL LEIBNIZ):



AS SURE AS I DIDN'T STEAL CALCULUS FROM NEWTON!

THE SYSTEM IS CALLED

BINARY NUMBERS.

THEY'RE BASED ON TWO!

OUR DECIMAL SYSTEM, BASED ON TEN, WAS A RESULT OF OUR HAVING TEN FINGERS — AN ACCIDENT OF NATURE! BINARY NUMBERS ARE WHAT WOULD HAVE EVOLVED IF WE'D BEEN BORN WITH TWO FINGERS, LIKE THE TREE SLOTH.

I'D COUNT BY FOURS, BUT I ONLY HAVE ONE FREE PAW!

TREE SLOTHS ALWAYS COUNT IN BINARY!

10

LOOK AT THE SYMBOL
"10" - "ONE-ZERO." FORGET
THAT IT USUALLY MEANS
TEN! FORGET IT! STOP
CALLING IT THAT! IS
THERE ANYTHING THERE
THAT SAYS "TEN?"
NO!! IT'S JUST A ONE
FOLLOWED BY A ZERO -
IN AND OF ITSELF, IT
HAS NOTHING TO DO
WITH TEN!!!

THE SYMBOL ONLY MAKES "TEN" FLASH THROUGH YOUR
MIND BECAUSE YOU'VE ALWAYS CALLED IT THAT... IT'S
LIKE A RITUAL: PERFORM IT OVER AND OVER AND IT
BECOMES AUTOMATIC!

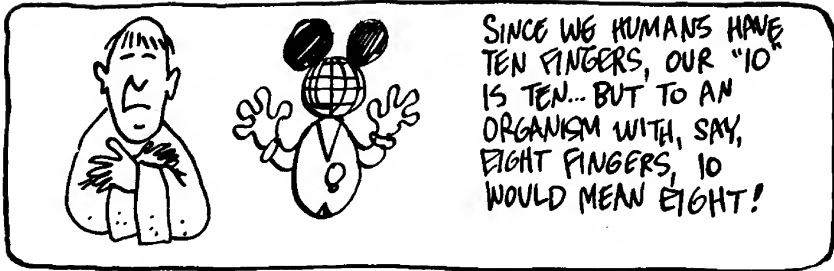


IN ACTUALITY, "10" MEANS:

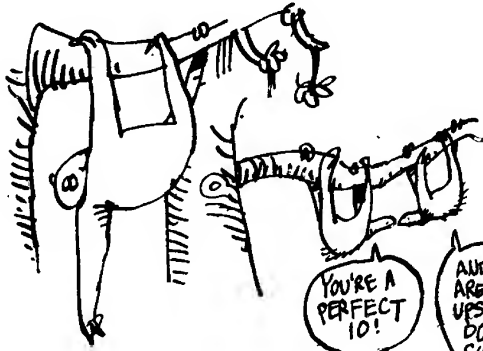


| (ONE) HANDFUL* AND
○ (ZERO) FINGERS LEFT OVER

*REMEMBER - ON P. 24, WE AGREED
TO CALL TEN FINGERS, NOT FIVE, A HUMAN HANDFUL!



IN THE CASE AT HAND, WITH JUST TWO FINGERS IN A HANDFUL... 10 MEANS **TWO!**



SO WE CAN WRITE:

$$10_{\text{BINARY}} = 2_{\text{DECIMAL}}$$

NOTE: DO NOT READ THIS AS "TEN EQUALS TWO." TEN DOES NOT EQUAL TWO!! "ONE-ZERO IN BINARY" EQUALS TWO!!





LIKEWISE, 100 — "ONE-ZERO-ZERO" — MEANS

1 HANDFUL OF HANDFULS.

IN DECIMAL, THAT'S 10×10 , OR A HUNDRED. WELL, IN BINARY IT'S 10×10 ALSO — BUT THAT ONLY AMOUNTS TO **FOUR!**

1000 IS

$$10 \times 10 \times 10 = 2 \times 2 \times 2 = 8$$

⋮

AND GENERALLY,

1 FOLLOWED BY N ZEROS IS:

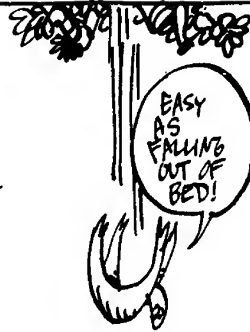
$$\underbrace{2 \times \dots \times 2}_{N \text{ TIMES}} = 2^N$$

("TWO TO THE NTH POWER").

IN THE COMPUTER AGE, EVERYONE WILL BE REQUIRED BY LAW TO MEMORIZE THE POWERS OF TWO, UP TO 2^{10} . BETTER NOT WAIT! AVOID JAIL AND DO IT NOW!

$$\begin{aligned} 1 &= 2^0 = 1 \\ 10 &= 2^1 = 2 \\ 100 &= 2^2 = 4 \\ 1000 &= 2^3 = 8 \\ 10000 &= 2^4 = 16 \\ 100000 &= 2^5 = 32 \\ 1000000 &= 2^6 = 64 \\ 10000000 &= 2^7 = 128 \\ 100000000 &= 2^8 = 256 \\ 1000000000 &= 2^9 = 512 \\ 10000000000 &= 2^{10} = 1024 \end{aligned}$$

TO MAKE THIS A BIT MORE CONCRETE —
 HERE'S HOW TO COUNT UP FROM 1 IN BINARY. IT'S JUST LIKE COUNTING IN DECIMAL, ONLY EASIER. IN DECIMAL, TO COUNT PAST A 9, YOU WRITE 0 AND CARRY 1. IN BINARY, YOU HAVE TO CARRY 1 EVERY OTHER NUMBER!!



BINARY	DECIMAL
0	0
1	1
10	2
11	3
100	4
101	5
110	6
111	7
1000	8
1001	9
1010	10
1011	11
1100	12
1101	13
1110	14
1111	15
10000	16
10001	17
10010	18
10011	19
10100	20
⋮	⋮
ETC!	ETC!

AS YOU MAY HAVE NOTICED, BINARY NUMBERS GET **LONNNNNNNNG** VERY FAST!

THIS MAKES THEM HARD FOR US HUMANS TO USE WITHOUT MAKING MISTAKES — BUT FOR COMPUTERS THEY'RE IDEAL !!



BINARY CALCULATION IS SIMPLE.
THERE ARE ONLY FIVE RULES
TO REMEMBER:

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

$$1 + 1 = 10$$

AND THE HANDY FIFTH RULE:

$$1 + 1 + 1 = 11$$



AS OPPOSED TO
100 SUMS IN
DECIMAL: 9+6,
7+5, 9+3, 8+4,
4+6, ETC ETC
ETC!!!

TO ADD TWO BINARY NUMBERS, PROCEED PLACE BY PLACE
FROM RIGHT TO LEFT, CARRYING A 1 WHEN NECESSARY.
HERE'S A STEP-BY-STEP EXAMPLE:

$\begin{array}{r} 1110 \\ 111 \\ \hline 1 \end{array}$	$\begin{array}{r} 1110 \\ 111 \\ \hline 01 \end{array}$	$\begin{array}{r} 11110 \\ 111 \\ \hline 101 \end{array}$	$\begin{array}{r} 11110 \\ 111 \\ \hline 10101 \end{array}$	<p>THE CARRIES</p>
--	---	---	---	------------------------

.....
A FEW SUMS TO PRACTICE ON:

$\begin{array}{r} 100 \\ + 1 \\ \hline \end{array}$	$\begin{array}{r} 11 \\ + 1 \\ \hline \end{array}$	$\begin{array}{r} 11001 \\ + 1100 \\ \hline \end{array}$	$\begin{array}{r} 11011 \\ + 11011 \\ \hline \end{array}$	$\begin{array}{r} 11111111 \\ + 11111111 \\ \hline \end{array}$
---	--	--	---	---

➔ WHAT IS THE RESULT OF ADDING A BINARY NUMBER TO ITSELF?

ANOTHER WONDERFUL FACT ABOUT BINARY:

SUBTRACTION IS DONE BY ADDING !!



THE METHOD IS CALLED USING "TWO'S COMPLEMENT." FIRST YOU INVERT THE NUMBER TO BE SUBTRACTED, SO THAT ALL ITS 1'S BECOME 0'S AND VICE VERSA. THEN ADD THE TWO NUMBERS AND ADD 1 TO THE SUM. IGNORE THE FINAL CARRY AND THAT'S THE ANSWER!

E.G.

$$\begin{array}{r} 1101 \\ -1100 \\ \hline \end{array}$$

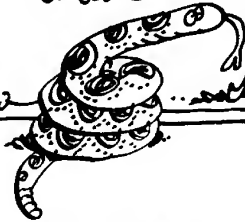
$$\begin{array}{r} 1101 \\ +0011 \leftarrow \text{INVERTED} \\ \hline 10000 \text{ SUM} \\ + \quad 1 \text{ ADD 1} \\ \hline 10001 \\ \leftarrow \text{IGNORE} \\ \hline 0001 \leftarrow \text{ANSWER.} \end{array}$$

BINARY MULTIPLICATION - AND ANY MULTIPLICATION - MAY ALSO BE DONE BY REPEATED ADDITION: TO MULTIPLY $A \times B$, JUST ADD A TO ITSELF B TIMES. LIKEWISE, DIVISION CAN BE DONE BY REPEATED SUBTRACTION.

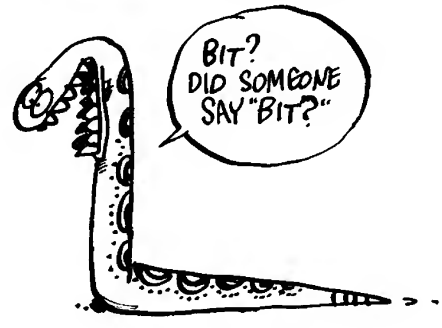
$$110 \times 11 =$$
$$\begin{array}{r} 110 \\ +110 \\ +110 \\ \hline 10010 \end{array} \left. \vphantom{\begin{array}{r} 110 \\ +110 \\ +110 \end{array}} \right\} 11 \text{ TIMES}$$

The computer can do all arithmetic by adding !!

The ADDER



BEFORE SHOWING HOW TO COMBINE LOGIC GATES INTO A BINARY ADDER, WE NEED A BIT OF TERMINOLOGY.



BIT?
DID SOMEONE SAY "BIT?"

BIT IS AN ABBREVIATION OF "BINARY DIGIT." IT REFERS TO A SINGLE 0 OR 1.

IS IT BINARY DIGIT OR BINARY DIGIT?

IT'S VERY COMMON TO GROUP BITS EIGHT AT A TIME, AND ANY STRING OF EIGHT BITS IS CALLED A **BYTE**. THERE ARE 2^8 OR 256, POSSIBLE BYTES, FROM 00000000 TO 11111111.

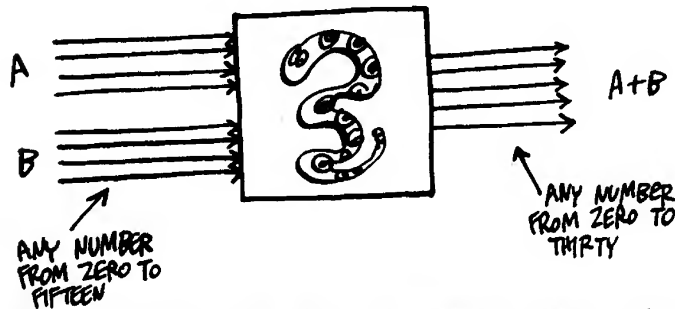
NOW LET'S SEE
WHAT AN ADDER
MIGHT LOOK LIKE.

THIS ADDER LOOKS LIKE
A POISONOUS ROSE...

TO SAVE DRAWING, WE'LL MAKE IT A FOUR-BIT ADDER, CAPABLE
OF ADDING TWO 4-BIT NUMBERS,
OR "NIBBLES." (YES, THEY'RE
REALLY CALLED THAT!)

$$\begin{array}{r} A = 1110 \\ B = 1011 \\ \hline 11001 \end{array}$$

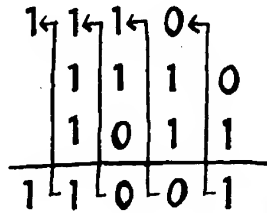
THE INPUT OF OUR ADDER MUST
CONSIST OF EIGHT BITS, FOUR FOR EACH
NIBBLE. THE OUTPUT MUST
BE FIVE BITS, THAT IS, A NIBBLE
PLUS ONE BIT FOR A POSSIBLE CARRY.
LIKE SO:



HOW TO PROCEED? ONE WAY IS TO MAKE A GIANT TRUTH TABLE,
MATCHING EVERY POSSIBLE COMBINATION OF INPUTS WITH THE
CORRECT OUTPUT, AND CONSTRUCTING A HUGE STEW OF ANDs
AND NOTs TO FORCE A SOLUTION. THIS IS
POSSIBLE, BUT THE COMPLEXITY OF THE TASK MIGHT
MAKE YOU THROW UP YOUR HANDS.

OR JUST
THROW UP,
IF YOU
HAVE NO
HANDS!

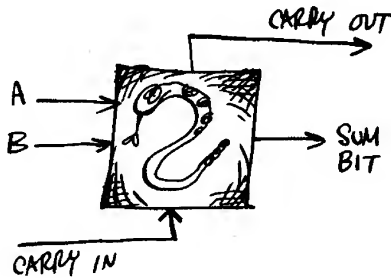
INSTEAD, RECALL HOW
ADDITION WORKS
IN PRACTICE:
COLUMN BY
COLUMN, WITH
A CARRY
BIT CARRYING
OUT OF ONE
COLUMN AND
INTO THE NEXT:



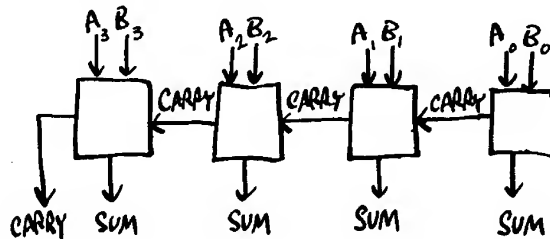
SO IT
SHOULD BE
POSSIBLE TO
MAKE A
4-BIT ADDER
OUT OF FOUR
1-BIT ADDERS!



THE 1-BIT
ADDER MUST HAVE
THREE INPUTS —
ONE FOR EACH OF
THE TWO SUMMAND
BITS AND ONE FOR
THE BIT CARRIED
IN — AND TWO
OUTPUTS — ONE
SUM BIT AND
ONE CARRY-OUT
BIT.



FOUR OF THESE CAN THEN BE HOOKED UP TO PRODUCE A
4-BIT ADDER:

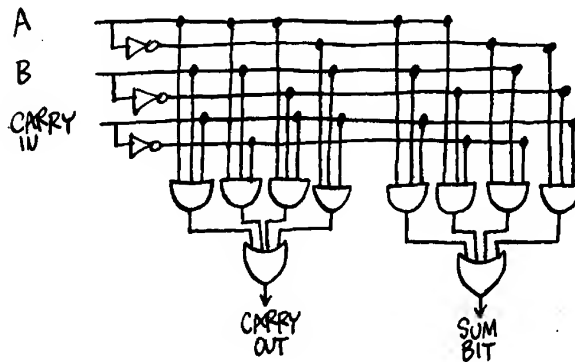


NOTE:
8 INPUTS
AND 5
OUTPUTS, AS
PROMISED!

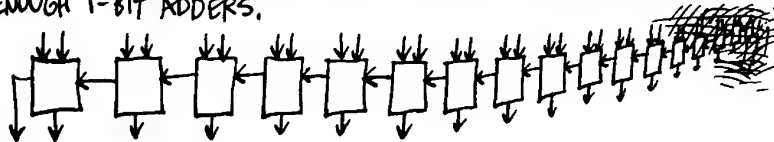
THE INPUT/OUTPUT TABLE FOR THE 1-BIT ADDER:

A	B	CARRY IN	CARRY OUT	SUM BIT
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

NOW THERE'S NOTHING TO IT! REMEMBER, LOGIC GATES CAN BE RIGGED UP TO PRODUCE ANY INPUT/OUTPUT TABLE. IN THIS CASE, JUST TREAT EACH OUTPUT COLUMN SEPARATELY:



YOU CAN ADD TWO NUMBERS OF ANY LENGTH BY HOOKING TOGETHER ENOUGH 1-BIT ADDERS.



CODING & COMBINATION

BINARY TASTES RIGHT!



THE IMPLICATION OF THE LAST TWO SECTIONS IS THAT BINARY IS THE "NATURAL" SYSTEM FOR ENCODING NUMBERS IN A MACHINE MADE OF ON/OFF SWITCHES. EVEN SO, COMPUTERS USE SEVERAL VARIATIONS ON THE BASIC IDEA.

INTEGERS, OR WHOLE NUMBERS — IF THEY AREN'T TOO LARGE — ARE ENCODED IN STRAIGHT BINARY. FOR INSTANCE,

185

WOULD BECOME

1	0	1	1	1	0	0	1
---	---	---	---	---	---	---	---

FLOATING POINT REPRESENTATION IS FOR LARGE OR FRACTIONAL NUMBERS. FOR EXAMPLE, 19,700,030.2 WOULD BE ENCODED AS THE BINARY EQUIVALENT OF

197	5
-----	---

,

MEANING 197×10^5 .

FLOATING POINT REPRESENTATION OFTEN INVOLVES ROUNDING OFF.

BINARY CODED DECIMAL REPRESENTS A NUMBER IN DECIMAL, BUT WITH EACH DIGIT ENCODED IN BINARY. 967, FOR INSTANCE, WOULD BECOME

1001	0110	0111
⏟	⏟	⏟
9	6	7

AND WHAT ABOUT NON-NUMERICAL INFORMATION — THE ALPHABET, PUNCTUATION MARKS, OTHER SYMBOLS, AND EVEN THE BLANK SPACE ??

SINCE THERE IS NO NATURAL WAY TO ENCODE THESE INTO 0'S AND 1'S, COMPUTER SCIENTISTS INVENTED AND ADOPTED A STANDARD CODE BY MUTUAL AGREEMENT:

ASCII,

THE AMERICAN STANDARD CODE FOR INFORMATION INTERCHANGE.

(ACTUALLY, ASCII IS USED BY EVERYONE BUT IBM, WHICH HAS ITS OWN CODE, CALLED EBCDIC.)



128

FIRST THREE BITS

	0	0	0	0	1	1	1	1
	0	0	1	1	0	0	1	1
	0	1	0	1	0	1	0	1
NEXT FOUR BITS	0000	NUL	DLE	SP	0	@	P	~
	0001	SOH	DC1	!	1	A	Q	a
	0010	STX	DC2	"	2	B	R	b
	0011	ETX	DC3	#	3	C	S	c
	0100	EOT	DC4	\$	4	D	T	d
	0101	ENQ	NAK	%	5	E	U	e
	0110	ACK	SYN	&	6	F	V	f
	0111	BEL	ETB	'	7	G	W	g
	1000	BS	CAN	(8	H	X	h
	1001	HT	EM)	9	I	Y	i
	1010	LF	SUB	*	:	J	Z	j
	1011	VT	ESC	+	;	K	[k
	1100	FF	FS	,	<	L	\	l
	1101	CR	GS	-	=	M]	m
	1110	SO	RS	.	>	N	^	n
	1111	SI	US	/	?	O	_	o
								DEL

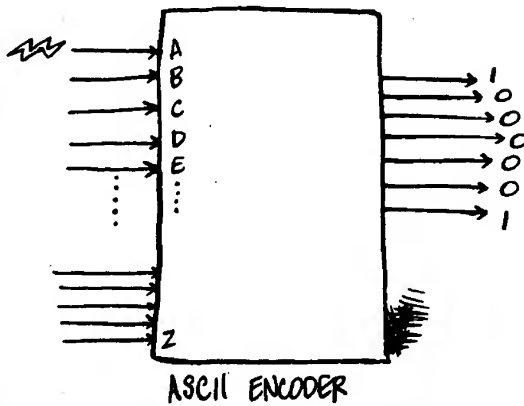
★ THUS, THE LETTER "T" IS ENCODED AS 101 0100... ETC!

★ THE FIRST TWO COLUMNS CONTAIN SYMBOLS FOR SUCH THINGS AS "START OF HEADING" (SOH) AND OTHER TEXTUAL DIRECTIONS.

TO ENCODE AND DECODE DATA, COMPUTERS USE LOGIC DEVICES CALLED, NATURALLY ENOUGH, **ENCODERS** AND **DECODERS**.

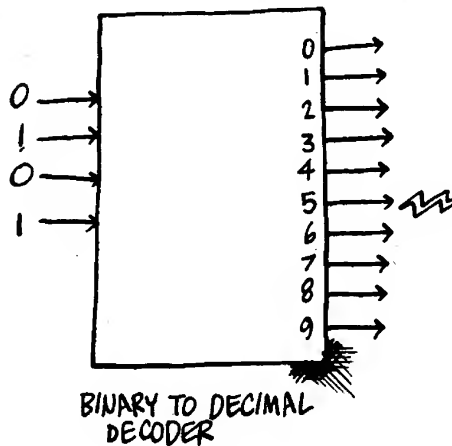
AN **ENCODER**

USUALLY HAS MANY INPUTS AND A FEW OUTPUTS. A SINGLE INPUT SIGNAL PRODUCES A PATTERN OF OUTPUTS. FOR EXAMPLE, A COMPUTER KEYBOARD IS ATTACHED TO AN ENCODER WHICH TRANSLATES A SINGLE KEYSTROKE INTO ITS ASCII CODE.



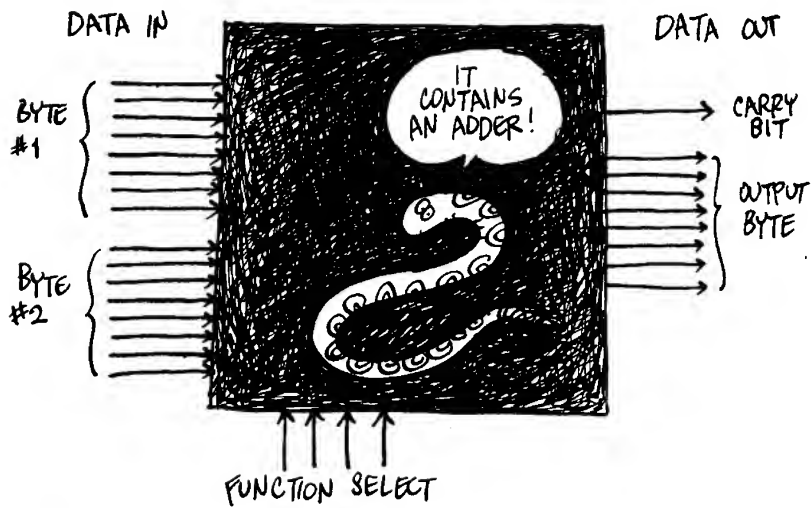
A **DECODER**

WORKS THE OTHER WAY AROUND, TRANSLATING A PATTERN OF BITS INTO A SINGLE OUTPUT SIGNAL. ONE DECODER CONVERTS A BINARY NIBBLE INTO A DECIMAL DIGIT. ANOTHER TRANSFORMS A SPECIFIED LOCATION, OR ADDRESS, IN MEMORY INTO A SIGNAL TO THAT MEMORY CELL. (SEE P. 155.)



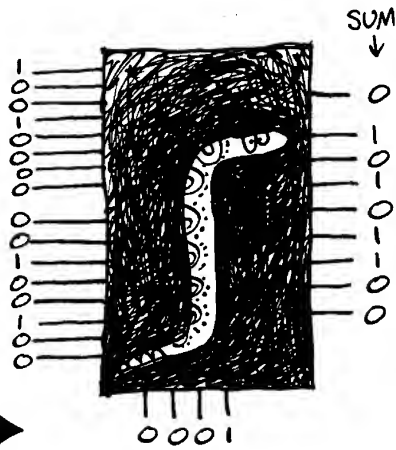
ONCE ALPHANUMERIC INFORMATION IS ENCODED IN BINARY STRINGS, IT IS READY TO BE PROCESSED BY THE COMPUTER'S MOST ELABORATE COMBINATION OF LOGIC GATES, THE

ARITHMETIC LOGIC UNIT (OR ALU, FOR SHORT)

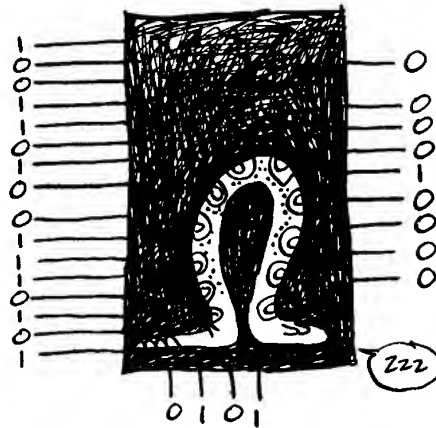


THIS IS THE MACHINE'S CENTRAL PROCESSOR, WHICH CAN ADD, SUBTRACT, MULTIPLY, COMPARE, SHIFT, AND PERFORM A WEALTH OF OTHER LOGICAL FUNCTIONS. THE DRAWING ABOVE REPRESENTS AN 8-BIT ALU, BUT THEY CAN RANGE FROM FOUR TO SIXTY BIT CAPABILITY, DEPENDING ON THE COMPUTER.

THE FUNCTION SELECT INPUTS DETERMINE WHICH ARITHMETIC OR LOGICAL FUNCTION THE ALU IS TO PERFORM, EACH FUNCTION HAVING ITS OWN BINARY CODE. FOR EXAMPLE, 0001 APPLIED TO FUNCTION SELECT MIGHT MEAN **ADD**, IN WHICH CASE



ANOTHER FUNCTION (0101, SAY) MIGHT **COMPARE** TWO BYTES, BIT BY BIT, AND OUTPUT A 1 WHEREVER THEY AGREE. (MEANWHILE, THE ADDER TAKES A NAP.)



YOU CAN GET AN IDEA OF A FANCY ALU'S CAPABILITIES FROM THE LIST ON PAGE 181.


THE ALU WOULD BE A COMPLETE CENTRAL PROCESSING UNIT, EXCEPT FOR ONE THING: IT'S UNABLE TO STORE RESULTS. RETURNING TO THE COOKING ANALOGY, WE MIGHT SAY THE ALU LACKS "COUNTER SPACE." WHERE WOULD GRANDMA BABBAGE BE WITHOUT SOMEPLACE TO SET DOWN HER SPAGHETTI?



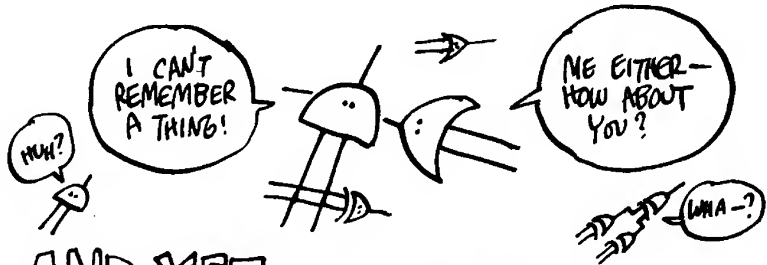
ALTHOUGH THE ALU CAN PERFORM MIRACLES OF INPUT/OUTPUT, IT CAN'T REMEMBER ANYTHING — AND THAT'S WHERE FLIP-FLOPS COME IN...



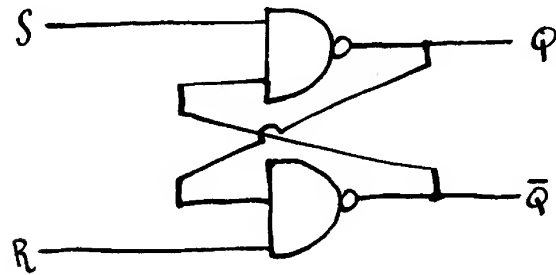
FLIP-FLOPS

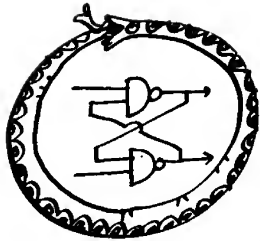


VERSATILE AS THEY MAY BE, THE LOGICAL COMBINATIONS WE'VE BEEN SKETCHING STILL HAVE NO MEMORY. THEIR OUTPUT CONTINUES ONLY AS LONG AS THE INPUT IS APPLIED.



AND YET — THERE IS A WAY TO HOOK THESE LOGICAL BUT SEMILE GATES TOGETHER INTO A GADGET THAT HOLDS AN OUTPUT INDEFINITELY: THE **FLIP-FLOP**. STARE AT THIS A MINUTE !!



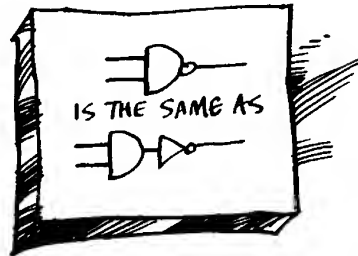


BESIDES THE STRANGE WAY A FLIP-FLOP EATS ITS OWN TAIL, PLEASE NOTE THE UNFAMILIAR GATE USED IN THE CONSTRUCTION. IT'S CALLED A

NAND GATE

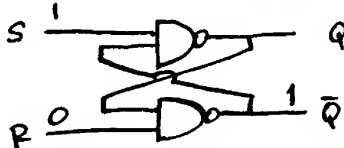
WHICH IS MERELY AN ABBREVIATION OF "NOT-AND."

A	B	NAND
1	1	0
0	0	1
0	1	1
1	0	1

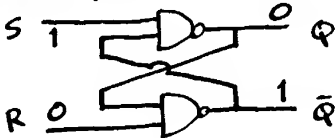


NOW FOR THE FLIP-FLOP IN ACTION:

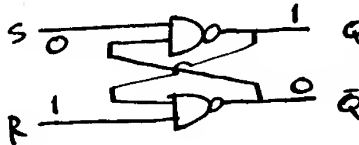
SUPPOSE THE INPUT IS $S=1, R=0$



THEN \bar{Q} MUST BE 1, BECAUSE NAND OUTPUTS 1 IF EITHER INPUT IS 0. COUPLING THIS BACK TO THE UPPER GATE GIVES $Q=0$.

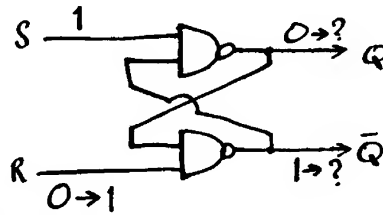


AND IF $S=0, R=1$? WELL, THAT'S JUST THE PREVIOUS DIAGRAM TURNED UPSIDE DOWN:

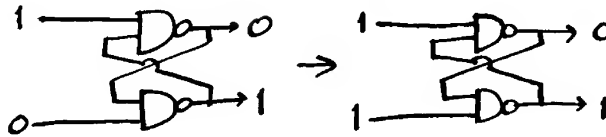


NOW WHAT HAPPENS
WHEN THE INPUT
CHANGES?

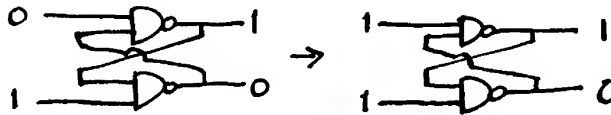
SUPPOSING WE BEGIN WITH
THE INPUT ($S=1, R=0$),
WHAT DOES CHANGING
IT TO ($S=1, R=1$)
DO TO THE FLIP-FLOP'S
OUTPUT?



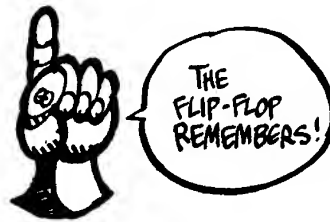
THE ANSWER IS: **NOTHING!** THE LOWER
NAND-GATE'S INPUT BECOMES (0, 1), SO ITS
OUTPUT \bar{Q} IS STILL 1, SO Q REMAINS 0.



BUT PRECISELY THE SAME LINE OF REASONING SHOWS NO CHANGE
IN OUTPUT WHEN INPUT CHANGES TO ($S=1, R=1$) FROM
($S=0, R=1$):



A LITTLE WEIRD, ISN'T IT?
THE SAME INPUT ($S=R=1$) CAN
PRODUCE TWO DIFFERENT
OUTPUTS, DEPENDING ON THE
PREVIOUS INPUT!



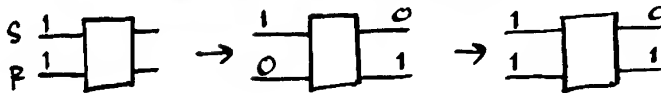
THE WAY A FLIP-FLOP IS USED IS THIS: IT BEGINS BY SITTING THERE WITH A CONSTANT INPUT OF (S=1, R=1) AND AN OUTPUT OF GOD-KNOWS-WHAT:



YOU **SET** THE FLIP-FLOP [I.E., MAKE Q=1] BY FLASHING A 0 MOMENTARILY DOWN THE S-WIRE, AND THEN RETURNING IT TO 1:



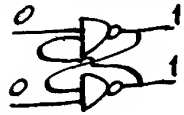
OR YOU CAN **RESET** IT [MAKE Q=0] BY FLASHING A 0 DOWN THE R-WIRE, THEN RETURNING IT TO 1:



IN EITHER CASE, AS LONG AS (1,1) KEEPS COMING IN, THE FLIP-FLOP WILL MAINTAIN ITS OUTPUT UNTIL IT'S CHANGED WITH ANOTHER INCOMING 0.



THE ONLY INPUT COMBINATION WE HAVEN'T CHECKED IS $(R=S=0)$.
IT'S EASY TO VERIFY THAT IT PRODUCES OUTPUT OF $Q=\bar{Q}=1$:

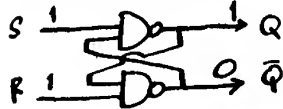


WHAT HAPPENS
WHEN THE
INPUT RETURNS
TO $(1, 1)$?

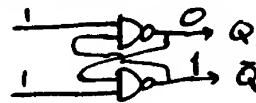


THE ANSWER IS NOT SO CLEAR: IT DEPENDS ON WHICH OUTPUT
HAPPENS TO FLOP FIRST!! (ONE OF THEM MUST.)

IF \bar{Q} IS FIRST TO
CHANGE, WE GET:



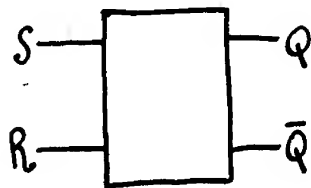
IF Q FLOPS FIRST,
HOWEVER:



SINCE THERE IS **NO WAY** OF KNOWING WHICH OF THESE
WILL ACTUALLY HAPPEN, AND WE DON'T WANT OUR
FLIP-FLOPS IN RANDOM STATES, THE INPUT $(S=0, R=0)$ IS

DISALLOWED.

WE CAN SUMMARIZE THE BASIC "RS" FLIP-FLOP LIKE SO:



S	R	Q	\bar{Q}
1	1	NO CHANGE	
1	0	0	1
0	1	1	0
0	0	DISALLOWED!	

FLIP-FLOP INPUTS ARE ALWAYS ARRANGED TO MAKE CERTAIN
THE DISALLOWED STATE CANNOT ARRIVE.



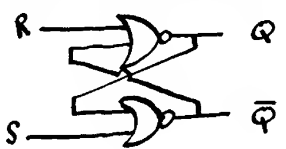
A LITTLE EXERCISE:

A NOR-GATE IS A SHORTHAND WAY OF WRITING "NOT OR:" I.E.

I AM THE TRUTH!

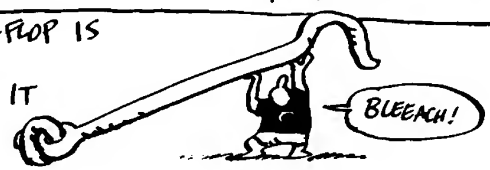
A	B	NOR
1	1	0
1	0	0
0	1	0
0	0	1

A BASIC R-S FLIP-FLOP MAY ALSO BE MADE OUT OF NOR-GATES:



1. WHAT IS THE OUTPUT WHEN R=0, S=1? WHEN S=0, R=1?
2. WHAT HAPPENS WHEN EACH OF THESE INPUT CONDITIONS CHANGES TO R=0, S=0?
3. WHAT IS THE OUTPUT WHEN R=1, S=1? WHAT HAPPENS WHEN THIS CHANGES TO R=0, S=0?
4. WHAT INPUT COMBINATION MUST BE DISALLOWED?
5. IF R=0, S=0, HOW DO YOU SET THIS FLIP-FLOP (I.E. MAKE Q=1)? HOW DO YOU RESET IT?

BY THE WAY, A FLIP-FLOP IS ALSO CALLED A **LATCH**, BECAUSE IT "LOCKS IN" DATA.

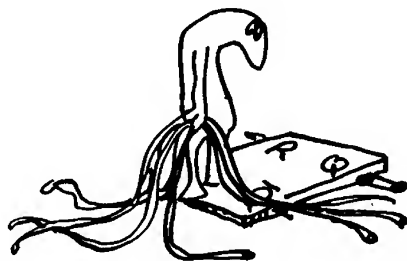


REGISTERS, COUNTERS, & GLITCHES

IF THE FLIP-FLOP IS A DEVICE FOR STORING ONE BIT, A REGISTER STORES SEVERAL BITS SIMULTANEOUSLY. IT'S LIKE A ROW OF BOXES, EACH HOLDING ONE BIT.

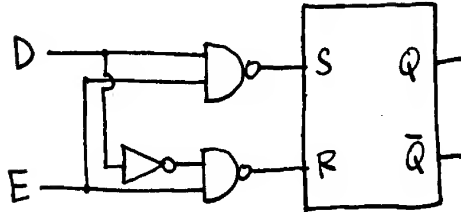


A ROW OF FLIP-FLOPS SHOULD DO THE JOB:...

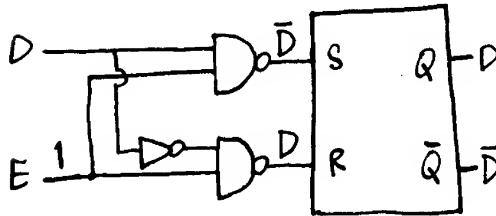


...SORT OF! BUT
IF YOU TRY AND
MAKE THIS WORK
BY HOOKING UP
SOME INPUTS TO
RS FLIP-FLOPS,
YOU MAY FIND
YOURSELF GROWING
CONFUSED!

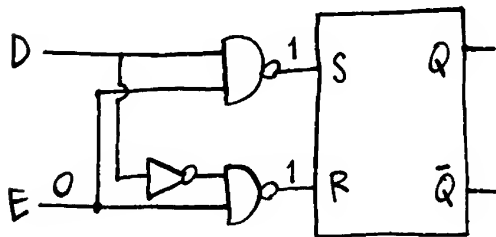
THE SOLUTION IS TO ADD A "GATING NETWORK" TO THE BASIC R-S FLIP-FLOP.



HERE "D" STANDS FOR DATA, AND "E" STANDS FOR ENABLE. NOTE THAT THE GATING NETWORK MAKES IT IMPOSSIBLE FOR R AND S TO BE ZERO SIMULTANEOUSLY.

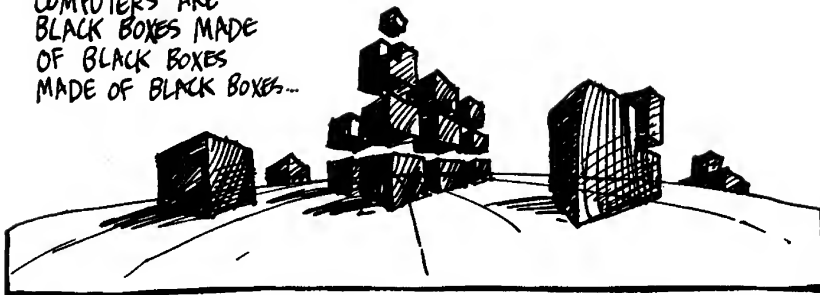


WHEN $E=1$, THEN $R=D$ AND $S=\bar{D}$ (NOT-D). HENCE, THE VALUE OF D IS STORED AT Q. IN OTHER WORDS, $E=1$ **ENABLES** THE BIT D TO BE LOADED INTO THE FLIP-FLOP.

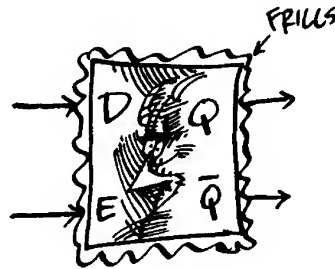


WHEN $E=0$, S AND R BOTH BECOME 1, AND THE FLIP-FLOP DOES NOT CHANGE. THAT IS, $E=0$ **BLOCKS** THE ARRIVAL OF MORE DATA.

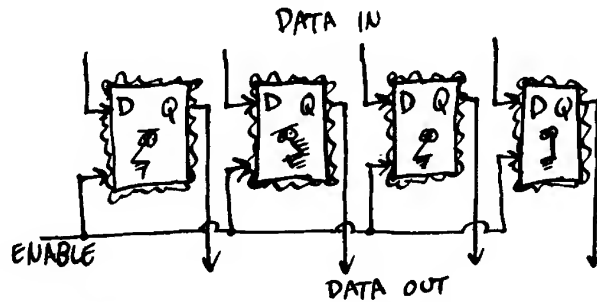
COMPUTERS ARE
BLACK BOXES MADE
OF BLACK BOXES
MADE OF BLACK BOXES...



SO - IN THE SPIRIT OF
IGNORING THE INNER WORKINGS
ONCE THEY'RE UNDERSTOOD
[OR EVEN WITHOUT EVER
UNDERSTANDING THEM],
WE INCORPORATE THE
GATING NETWORK
INTO THE BOX, AND DRAW
THE GATED LATCH LIKE
SO

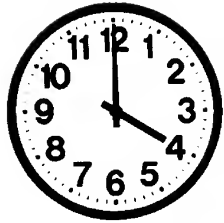


THEN HERE'S A **PARALLEL REGISTER**: NOT THE ONLY
KIND OF REGISTER, BUT A GENUINE MEMBER OF THE BREED!



ON E=1,
FOUR BITS
ARE SIMUL-
TANEOUSLY
LOADED
INTO THE
LATCHES!

NOW WHAT CONTROLS THE 'ENABLE' INPUT?



A BASIC FACT OF COMPUTER LIFE:

AS SOON AS YOU BEGIN STORING DATA, QUESTIONS OF TIMING ARISE: HOW LONG DO YOU STORE IT? WHEN DO YOU MOVE IT? HOW DO YOU SYNCHRONIZE SIGNALS? THESE ISSUES ARE SO CRITICAL THAT LOGIC WITH MEMORY

IS CALLED **SEQUENTIAL**, TO DISTINGUISH IT FROM THE PURELY **COMBINATIONAL** LOGIC OF MEMORY-LESS NETWORKS. TO KEEP THE SEQUENTIAL LOGIC IN STEP,

ALL COMPUTERS HAVE CLOCKS!

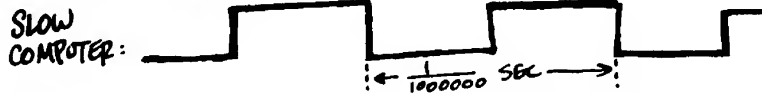
THE CLOCK'S PULSE IS THE COMPUTER'S HEARTBEAT—ONLY INSTEAD OF A WARM, RAGGED HUMAN HEARTBEAT, LIKE THIS—



THE COMPUTER'S PULSE IS SQUARE AND COLD:



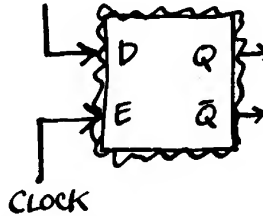
ONE CLOCK PULSE IS THE BURST OF CURRENT WHEN CLOCK OUTPUT = 1. ONE CYCLE IS THE INTERVAL FROM THE BEGINNING OF A PULSE TO THE BEGINNING OF THE NEXT. DEPENDING ON THE COMPUTER, THE CLOCK FREQUENCY MAY BE HUNDREDS OF THOUSANDS TO **BILLIONS** OF CYCLES PER SECOND!



THE IDEA OF USING A CLOCK IS THAT THE COMPUTER'S LOGICAL STATE SHOULD CHANGE ~~ONLY~~ ON THE CLOCK PULSE. IDEALLY, WHEN THE CLOCK HITS 1, ALL SIGNALS MOVE, THEN STOP ON CLOCK=0. THEN GO... THEN STOP... THEN GO...



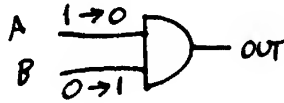
A TYPICAL EXAMPLE IS TO ATTACH THE CLOCK TO THE "ENABLE" INPUT OF A GATED LATCH, IN WHICH CASE THE LATCH BECOMES KNOWN AS A "D FLIP-FLOP."



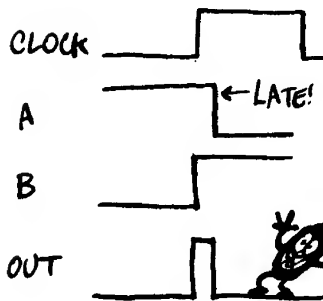
THEN A NEW BIT OF DATA IS LOADED AT EVERY CLOCK PULSE!

HEE HEE

UNFORTUNATELY, THINGS ARE RARELY IDEAL! IT TAKES A NON-ZERO TIME FOR A SIGNAL TO PASS ALONG A WIRE, SO THINGS ARE NEVER PERFECTLY SYNCHRONIZED. FOR EXAMPLE, SUPPOSE AT AN AND GATE, ONE INPUT IS CHANGING FROM 1 TO 0, AND THE OTHER FROM 0 TO 1:



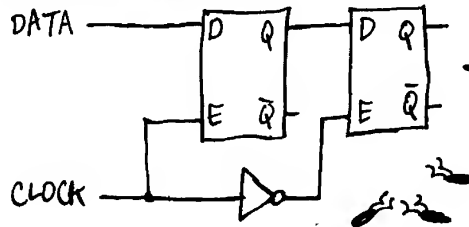
IF A CHANGES AFTER B, THE OUTPUT WILL HAVE AN UNWANTED PULSE.



THAT PULSE IS A **GLITCH**, AND BRIEF AS IT IS, IT CAN CAUSE A FLIP-FLOP TO FLOP!

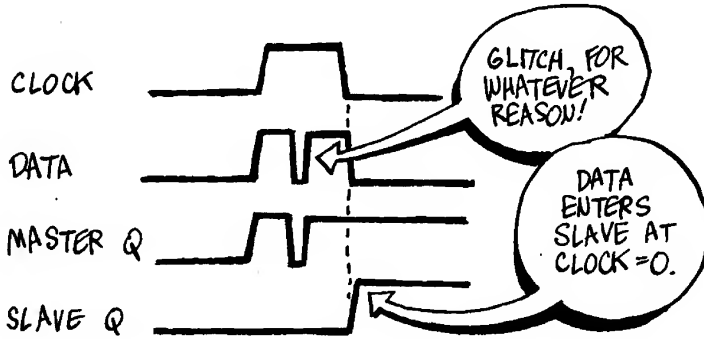
WE'RE UNAVOIDABLE!

THE GLITCH IS DEFEATED BY THE
MASTER-SLAVE FLIP-FLOP:

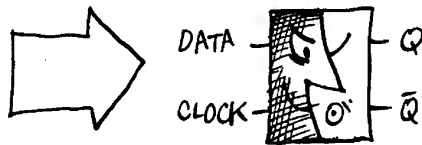


WE WERE BORN EQUAL, MASTER, AND ONLY MY POSITION MAKES ME A SLAVE!

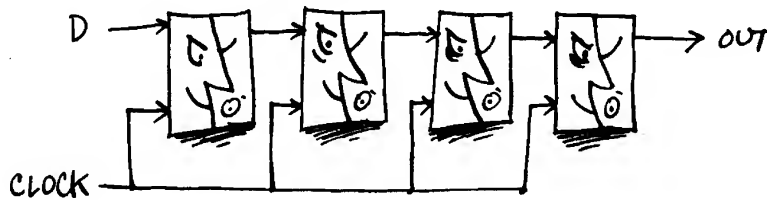
THE INVERTED CLOCK SIGNAL TO THE SLAVE FLIP-FLOP DELAYS THE DATA INPUT FROM ARRIVING AT THE SLAVE UNTIL THE **END** OF A CLOCK PULSE, AFTER ALL GLITCHES HAVE DIED OUT. FOR EXAMPLE, SUPPOSE WE WANT TO LOAD THE BIT 1 INTO THE FLIP-FLOP.



AS USUAL, WE DRAW THE WHOLE THING AS A SINGLE BOX!

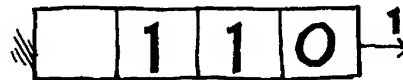


STRINGING A NUMBER OF MASTER-SLAVE FLIP-FLOPS TOGETHER MAKES A **SHIFT REGISTER**



DATA ENTER A SHIFT REGISTER ONE BIT AT A TIME, SHIFTING TO THE RIGHT WITH EACH NEW CLOCK PULSE.

FOR EXAMPLE, THE NIBBLE 1101 WOULD ENTER THE SHIFT REGISTER LIKE THIS:



EACH CLOCK PULSE BRINGS A NEW BIT INTO THE REGISTER. (WHY DOESN'T THE BIT TRAVEL ALL THE WAY THROUGH ON ONE PULSE? BECAUSE OF THE MASTER-SLAVE FLIP-FLOPS!)

LIKewise, THE NIBBLE SHIFTS OUT ONE BIT AT A TIME.



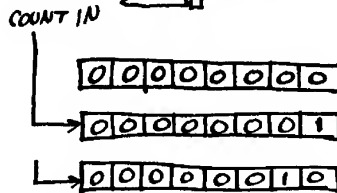
SHIFT REGISTERS ARE USEFUL WHEN INFORMATION IS TO BE TRANSMITTED **SERIALLY**, OR ONE BIT AT A TIME.

FINALLY, A SPECIAL KIND OF REGISTER: THE **COUNTER**.

IS THAT LIKE THE COUNTER MONTE CRISTO?

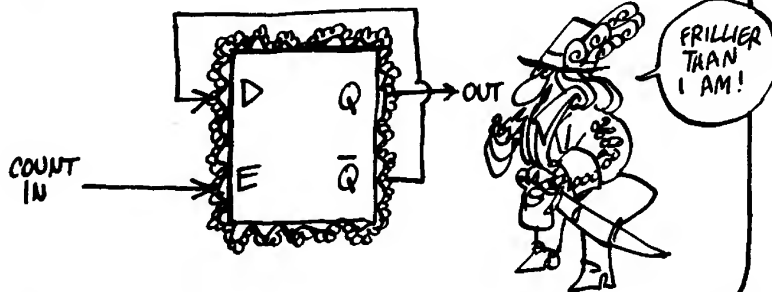


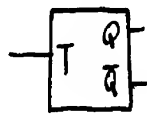
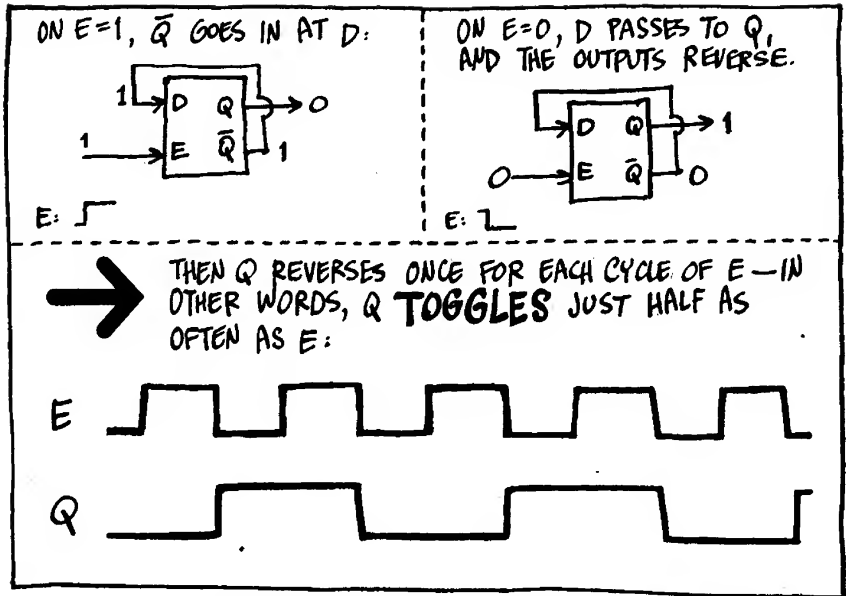
A COUNTER IS JUST WHAT IT SOUNDS LIKE: SOMETHING THAT COUNTS. IN OTHER WORDS, IT'S A REGISTER THAT **INCREMENTS** ITSELF—ADDS 1 TO ITS CONTENTS—WHENEVER A "COUNT" SIGNAL ARRIVES:



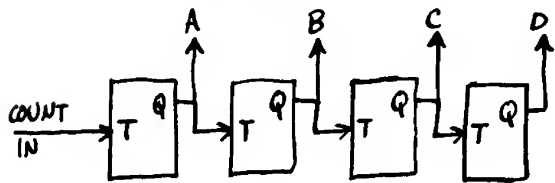
ETC!

DESCRIBED IN THAT WAY, A COUNTER SOUNDS EASY TO MAKE: JUST COMBINE AN ADDER WITH A REGISTER! THIS WOULD IN FACT WORK, BUT THERE'S AN EVEN SLICKER WAY, BASED ON ANOTHER FANCY FLIP-FLOP. CONSIDER THIS MASTER-SLAVE FLIP-FLOP, COUPLED BACK ON ITSELF:





AS USUAL, WE ABBREVIATE THE WHOLE CIRCUIT BY THIS SIMPLER BOX. THE "T" IS FOR TOGGLE, TO INDICATE THAT THE FLIP-FLOP TOGGLES WHENEVER $T=1$. THEN HERE'S OUR COUNTER: EACH FLIP-FLOP TOGGING AT HALF THE RATE OF THE ONE TO ITS LEFT:



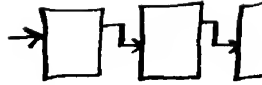
COUNT IN	D	C	B	A
0	0	0	0	0
1	0	0	0	1
2	0	0	1	0
3	0	0	1	1
4	0	1	0	0
5	0	1	0	1
6	0	1	1	0
7	0	1	1	1
8	1	0	0	0
9	1	0	0	1
⋮	⋮	⋮	⋮	⋮



A FEW ITEMS OF NOTE:

1

THIS COUNTER IS CALLED AN "ASYNCHRONOUS RIPPLE COUNTER," BECAUSE THE COUNT RIPPLES THROUGH FROM ONE FLIP-FLOP TO THE NEXT. THIS CAUSES A SLIGHT DELAY BEFORE THE COUNT IS REGISTERED.



2

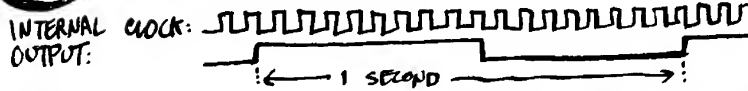
WHEN THE 16TH COUNT PULSE ARRIVES, THE COUNTER RETURNS TO 0. TO GO HIGHER THAN 15, MORE FLIP-FLOPS ARE NEEDED.



THIS 14-BIT COUNTER CAN GO FROM 0 TO $2^{14}-1 = 16,383$

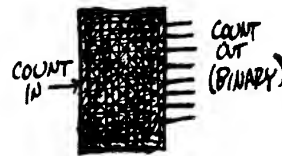
3

THE NTH FLIP-FLOP IN A RIPPLE COUNTER DIVIDES THE INCOMING PULSE BY 2^N . THIS IS THE PRINCIPLE ON WHICH DIGITAL WATCHES ARE BASED: A HIGH-FREQUENCY INTERNAL CLOCK PULSE IS DIVIDED TO A RATE OF PRECISELY ONE CYCLE PER SECOND.



4

THERE ARE ALSO SYNCHRONOUS COUNTERS, WHICH REGISTER ALL BITS SIMULTANEOUSLY, AND COUNTERS WHICH RETURN TO 0 ON ANY PREASSIGNED NUMBER. IN ANY CASE, FROM NOW ON, A COUNTER IS JUST ANOTHER BLACK BOX !!

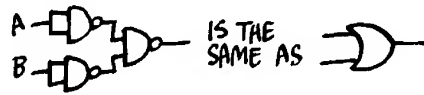
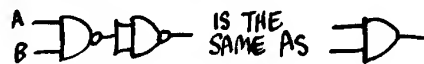
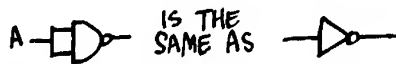


EXERCISES

THE AMAZING NAND:



1. SHOW THAT



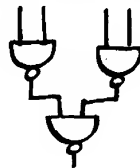
CONCLUDE THAT \Rightarrow ALL LOGIC CAN BE DERIVED FROM THE SINGLE RELATION NAND!!!

2. CAN THE SAME BE DONE WITH NOR?

3. SHOW THAT



IS THE SAME AS



REDRAW THE ADDER ON P. 126 USING ONLY NAND-GATES.

4. GIVEN A 4-BIT SHIFT REGISTER,



SHOW ITS CONTENTS AFTER EACH OF FOUR CLOCK PULSES AS THE NIBBLE 0011 IS ENTERED.

5. HOW WOULD YOU ATTACH A BUZZER TO A COUNTER TO SOUND WHEN THE COUNT HITS NINE (=1001 IN BINARY)?
HINT: LOOK AT THE SEAT BELT BUZZER ON P. 109.

6. CONVINCED YOURSELF THAT ATTACHING INVERTERS TO THE OUTPUTS MAKES A COUNTER COUNT BACKWARDS.

NOW IN CASE YOU'RE FEELING STRANGLER BY SPAGHETTI—

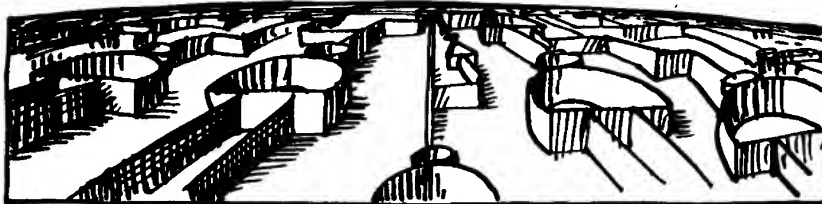
THE TANGLED DIAGRAMS
ON THE PRECEDING
PAGES WERE NEVER
INTENDED TO TRACE
THE COMPLETE WIRING
DIAGRAM OF ANY
COMPUTER. RATHER,
THEY ARE MEANT TO
DEMONSTRATE HOW
THE COMPUTER'S
ESSENTIAL FUNCTIONS—
MATH, COMPARISON,
DECODING, DATA
SELECTION AND STORAGE—
ALL DEPEND ON
SIMPLE LOGIC.



NOW THAT YOU PRESUMABLY
BELIEVE IN THE POWER OF
LOGIC, NO MORE WIRING
DIAGRAMS ARE NEEDED!

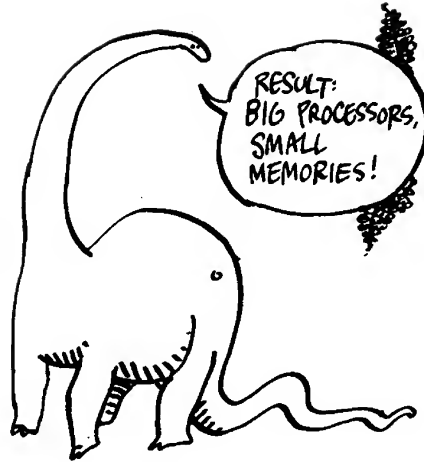


ONWARD,
TO HIGHER
LEVELS!



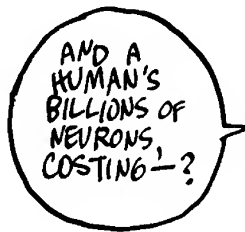


IN THE INFANCY OF ELECTRONIC COMPUTING, MEMORY WAS ALWAYS MORE EXPENSIVE THAN SHEER COMPUTING POWER. PLENTY OF PROCESSING COULD BE DONE WITH RELATIVELY FEW COMPONENTS, BUT EVERY INCREASE IN MEMORY SIMPLY MEANT **MORE** = MORE ACTUAL, PHYSICAL PLACES TO STORE THINGS!



RESULT:
BIG PROCESSORS,
SMALL
MEMORIES!

SINCE THEN, RESEARCH INTO MEMORY TECHNOLOGY HAS BROUGHT DOWN THE COST CONSIDERABLY. FOR A FEW HUNDRED DOLLARS YOU CAN BUY A MICRO WITH OVER 64,000 BYTES OF MEMORY, COMPARED WITH ENIAC'S MEMORY OF ABOUT 100 NUMBERS* — AT A COST OF MILLIONS!!



AND A HUMAN'S BILLIONS OF NEURONS, COSTING —?



*ENIAC DID NOT COMPUTE IN BINARY.

THE SAME RESEARCH EFFORT,
HOWEVER, HAS PRODUCED A
BEWILDERING ARRAY OF MEMORY
TYPES AND TECHNOLOGIES !!

CARD MEMORIES,
TAPE MEMORIES,
DRUM, DISK, BUBBLE,
OPTICAL, CORE, CHARGE-
COUPLED DEVICE, AND
SEMICONDUCTOR MEMORIES;
VOLATILE AND NON-VOLATILE,
DYNAMIC AND STATIC,
DESTRUCTIVE AND
NON-DESTRUCTIVE; READ-
WRITE, READ-ONLY,
PROGRAMMABLE READ-ONLY,
ERASABLE PROGRAMMABLE
READ-ONLY...:PANT:
:PUFF:

HAVE I
FORGOTTEN
ANYTHING?

I DON'T
REMEMBER..

WELL, ONE HAS TO BEGIN SOMEWHERE !! 

AN IMPORTANT
DISTINCTION EXISTS
BETWEEN

ELECTRONIC

AND

**ELECTRO-
MECHANICAL**

MEMORY DEVICES.

ELECTRONIC MEMORIES, WITH NO MOVING PARTS, ARE AS FAST AS THE REST OF THE COMPUTER.

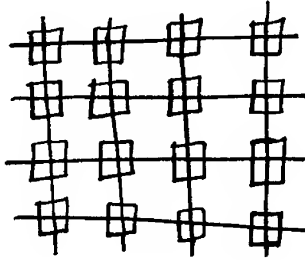
ELECTROMECHANICAL MEMORIES HAVE MOVING PARTS, LIKE DISKS OR REELS OF TAPE. THIS MAKES THEM SLOW—HOW SLOW DEPENDING ON THE TYPE OF MEMORY.



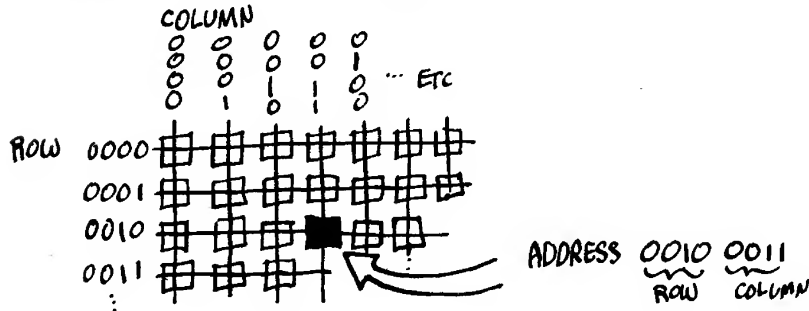
ELECTRONIC MEMORIES' SPEED MAKES THEM IDEAL FOR THE COMPUTER'S MAIN, OR INTERNAL MEMORY, WHILE ELECTRO-MECHANICAL MEMORIES ARE USED FOR AUXILIARY STORAGE OUTSIDE THE MACHINE.

ELECTROMAGNETIC MEMORIES COMPENSATE FOR THEIR SLOWNESS WITH A GIGANTIC CAPACITY. ONE HARD DISK CAN STORE UP TO TEN MILLION BYTES, COMPARED WITH A TYPICAL MICRO'S MAIN MEMORY OF 65,536 ($=2^{16}$) BYTES.

INTERNAL MEMORY CAN BE THOUGHT OF AS A SIMPLE GRID, WITH A CELL AT EACH INTERSECTION. DEPENDING ON THE COMPUTER, EACH CELL CAN HOLD ONE BYTE, TWO BYTES, OR MORE.

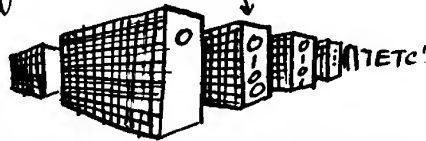


EVERY CELL HAS A UNIQUE ADDRESS, SPECIFYING WHERE IT SITS IN THE GRID.



IN PRACTICE, THERE MAY BE MANY SUCH GRIDS, IN WHICH CASE THE ADDRESS SPECIFIES THE GRID NUMBER, AS WELL AS THE ROW AND COLUMN WITHIN IT.

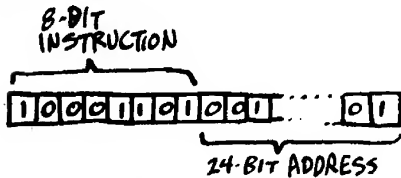
ADDRESS
0101 1001 1110
 GRID ROW COLUMN



NOTE:

DO NOT CONFUSE A CELL'S ADDRESS WITH ITS CONTENTS !!

WHAT IS THE MAXIMUM NUMBER OF CELLS THE COMPUTER CAN ADDRESS? THIS DEPENDS ON THE LENGTH AND STRUCTURE OF THE COMPUTER'S "WORDS." FOR EXAMPLE, A 32-BIT MACHINE MAY INTERPRET THE FIRST 8 BITS AS AN INSTRUCTION...



... AND THE REMAINING 24 BITS AS AN ADDRESS.

IN THAT CASE, ADDRESSES CAN BE ANYTHING BETWEEN

0000...0
AND

1111...1 = $2^{24} - 1$

GIVING 2^{24} POSSIBLE MEMORY CELLS.



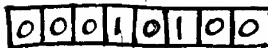
AN 8-BIT MICRO, ON THE OTHER HAND, MIGHT PROCESS THREE BYTES IN SUCCESSION:



AN INSTRUCTION,

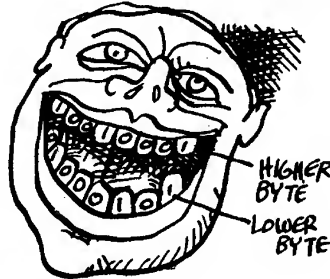


THE FIRST HALF OF AN ADDRESS,

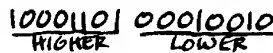


AND THE SECOND HALF OF AN ADDRESS.

▶ HERE THE ADDRESS IS 16 BITS LONG, GIVING $2^{16} = 65,536$ POSSIBLE ADDRESSES.



16-BIT WORDS ARE OFTEN SPLIT LIKE THIS INTO HIGHER-LEVEL AND LOWER-LEVEL BYTES.



TO MAKE ADDRESSES SHORTER AND MORE READABLE, THEY'RE OFTEN EXPRESSED IN

HEXADECIMAL

OR BASE-16, NUMERALS.

$$10_{\text{HEX}} = 16_{\text{DECIMAL}}$$

$$100_{\text{HEX}} = 16^2 = 256$$

$$1000_{\text{HEX}} = 16^3 = 4096$$

⋮

ETC!



JUST AS BASE-10 NUMBERS REQUIRE THE DIGITS 0-9, SO HEXADECIMAL NEEDS DIGITS FROM 0 TO FIFTEEN. THE EXTRAS ARE REPRESENTED BY THE LETTERS A-F:

DECIMAL	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
HEX	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F

FOR EXAMPLE:

$$4A0D_{\text{HEX}} =$$

$$4 \times 16^3$$

$$+ 10 \times 16^2$$

$$+ 0 \times 16$$

$$+ 13 \times 1$$

$$18,957_{\text{DECIMAL}}$$

TO CONVERT BINARY TO HEX:
GROUP THE BINARY NUMBER INTO NIBBLES, STARTING FROM THE RIGHT. CONVERT EACH NIBBLE TO A HEX DIGIT!

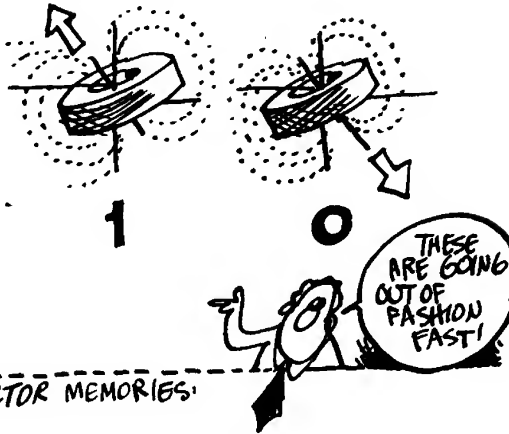
$$\begin{array}{cccc} 101 & 1110 & 0101 & 1011 \\ \hline 5 & C & 5 & B \end{array}$$

TO CONVERT HEX TO BINARY, JUST REVERSE THE PROCESS.

FROM THE HARDWARE POINT OF VIEW, THERE ARE THREE MAIN TYPES OF INTERNAL MEMORY.

CORE

MEMORIES USE LITTLE MAGNETIC DOUGHNUTS - "CORES." EACH CORE CAN BE ELECTRICALLY MAGNETIZED IN ONE OF TWO DIRECTIONS, REPRESENTING 0 AND 1.



AND TWO SEMICONDUCTOR MEMORIES:

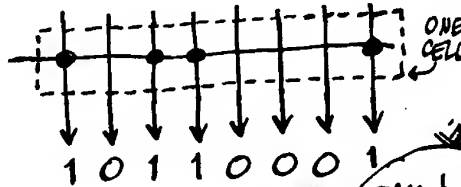
RAM

USES FLIP-FLOPS TO STORE BITS - SO EACH MEMORY CELL IS ESSENTIALLY A (PARALLEL) REGISTER!



ROM

INDICATES A 1 OR 0 AT EACH GRID POINT BY THE PRESENCE OR ABSENCE OF AN ELECTRIC CONNECTION THERE.



RAM

STANDS FOR
"RANDOM ACCESS MEMORY,"
MEANING THAT ANY CELL
CAN BE ACCESSED DIRECTLY.
ROM AND CORE MEMORIES
ALSO PROVIDE RANDOM
ACCESS, BUT FOR SOME
REASON, RAM HOGGED THE
NAME!

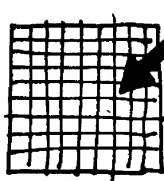


A
CASE OF
SPECIES
CONFUSION...

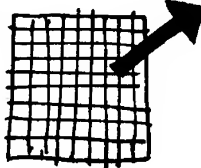
ROM

STANDS
FOR "READ-ONLY MEMORY."

ROM-AN
STYLE
LETTERING



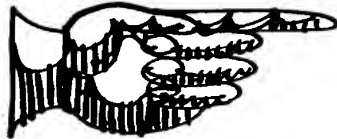
RAM



ROM

THE PRACTICAL
DIFFERENCE
BETWEEN THEM IS
THAT YOU CAN ONLY
READ WHAT'S IN
ROM, WHILE WITH
RAM YOU CAN READ
THINGS OUT OR
WRITE THEM IN
WITH EQUAL EASE.

IN GENERAL!



WHEN YOU LOAD A PROGRAM
INTO THE COMPUTER, IT IS
STORED IN **RAM**.

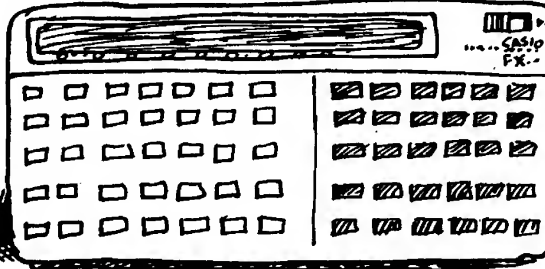
UNFORTUNATELY,
RAM IS
VOLATILE,



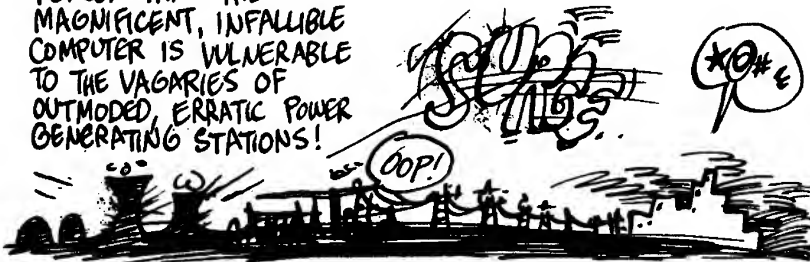
IT FORGETS EVERYTHING WHEN THE POWER IS TURNED OFF.

FOR EXAMPLE, I OWN A BATTERY-POWERED POCKET COMPUTER WITH 1680 BYTES OF RAM. IT CAN STORE UP TO TEN PROGRAMS EVEN WHEN I TURN IT OFF BECAUSE IT KEEPS SOME ELECTRICITY RUNNING THROUGH MEMORY.

BUT WHEN THE BATTERY DIES...
BYE-BYE,
PROGRAMS!



RAM VOLATILITY IS ONE REASON THAT THE MAGNIFICENT, INFALLIBLE COMPUTER IS VULNERABLE TO THE VAGARIES OF OUTMODDED, ERRATIC POWER GENERATING STATIONS!



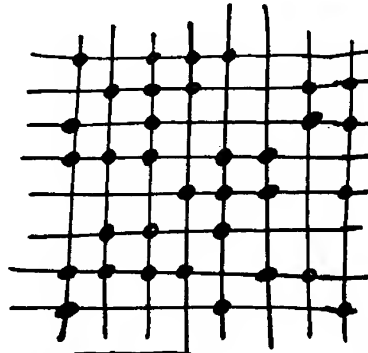
ROM - "READ-ONLY MEMORY" -

ONCE ITS CONTENTS ARE ENTERED,
CAN NEVER BE REWRITTEN.*
ORDINARILY, ROM IS PROGRAMMED
AT THE FACTORY, BUT THERE ARE
NOW ALSO PROMS - PROGRAMMABLE
ROMS - WHICH CAN BE CUSTOM-
PROGRAMMED TO THE USER'S
SPECIFICATIONS.

* EXCEPT FOR
EPROM -
ERASABLE
PROGRAMMABLE
ROM - BUT ONE
WON'T GET
INTO THAT!



UNLIKE RAM, ROM
IS NON-VOLATILE:
IT KEEPS ITS
CONTENTS EVEN
WITHOUT POWER.
AFTER ALL, IT'S
NOTHING BUT A
HUGE GRID OF WIRES
WITH PHYSICAL
CONNECTIONS AT SOME
INTERSECTIONS.
THE CONNECTIONS
REMAIN, REGARDLESS
OF ELECTRIC CURRENT.



AND WHEN I SAY
"HUGE," I MEAN
"TINY"!!

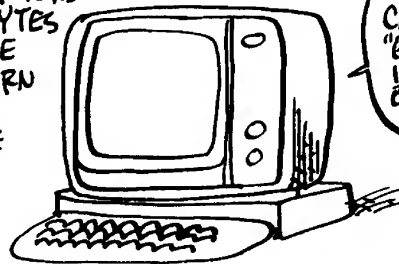


SOME TYPICAL USES OF ROM:

MOST VIDEO GAME CARTRIDGES ARE PROGRAMMED IN ROM. JUST PLUG IT IN AND IT'S READY TO GO! BUT OF COURSE, IT CAN'T BE REPROGRAMMED EITHER...



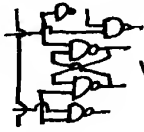
MANY PERSONAL COMPUTERS HAVE THOUSANDS OF BYTES OF ROM TO STORE THE PROGRAM WHICH IN TURN ALLOWS THE MACHINE TO "UNDERSTAND" THE LANGUAGE CALLED BASIC.



IT'S CALLED "BUILT-IN BASIC!"

AND, AS WE'LL SEE, ROM PLAYS AN IMPORTANT ROLE IN THE COMPUTER'S CONTROL SECTION.

BEHIND THE
EXPLOSIVE
GROWTH OF
RAM AND
ROM IS...
THE INCREDIBLE
SHRINKING
TECHNOLOGY!



ETCHED ON SILICON CHIPS,
THE DENSITY OF
COMPONENTS PER
CHIP HAS BEEN
DOUBLING
EVERY YEAR!



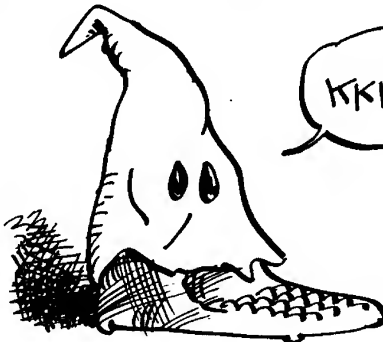
THE STANDARD MEASURE OF CHIP STORAGE IS THE
K, SHORT FOR "KILO" ("CHILLO" IS GREEK FOR 1000),
IN COMPUTERESE IT MEANS 2^{10} , THE POWER OF
TWO CLOSEST TO 1000:

ALMOST
GREEK
FOR ALMOST
1000!



K = 1024

THE FIRST RAM
CHIP WITH 1K
BITS OF STORAGE
WAS A
SENSATION —
BUT NOW
64K IS
COMMON, AND THE
256K CHIP
HAS ARRIVED!
WHAT'S
NEXT?





DESPITE THE GROWTH OF RAM CAPACITY,
SOMETIMES IT IS NOT THE ANSWER TO
EVERY PRAYER !!

SHOW US
THE WAY TO
STORE
MORE
THAN INTERNAL
RAM CAN
HOLD!

LET US
PROTECT
OUR DATA
FROM
POWER
LOSSES!

GRANT
US A PROGRAM
LIBRARY OF
FREQUENTLY
USED
ROUTINES!

GIVE
US
LASER-
POWERED
GADGETS!



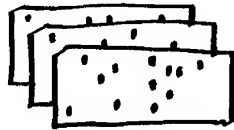
THE ANSWER?

mass storage.

AS THE NAME IMPLIES,
MASS STORAGE IS MEMORY
THAT CAN STORE A LOT!!
ALMOST ALL MASS STORAGE
DEVICES ARE NON-VOLATILE
AND HAVE A MECHANICAL
COMPONENT THAT MAKES THEM
MUCH SLOWER THAN ELECTRONIC
RANDOM ACCESS MEMORIES:

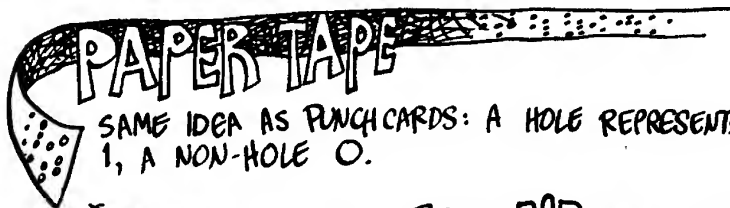


FOR EXAMPLE.



PUNCH CARDS.

THE CARDS OF JACQUARD, BABBAGE,
AND HOLLERITH ARE STILL IN USE!



PAPER TAPE

SAME IDEA AS PUNCH CARDS: A HOLE REPRESENTS
1, A NON-HOLE 0.



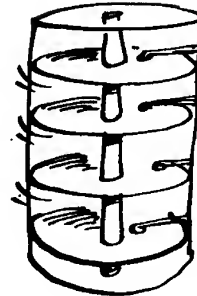
MAGNETIC TAPE

STORES BITS AS SMALL MAGNETIC REGIONS, WHICH MAY
BE MAGNETIZED IN ONE OF TWO DIRECTIONS,
REPRESENTING 1 OR 0.

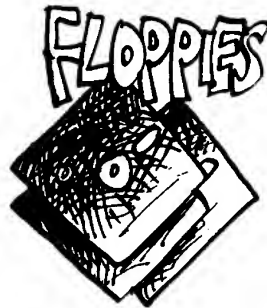
FASTER, LESS BULKY, AND THE CURRENT STORAGE OF CHOICE IS THE



DISKS ALSO STORE BITS AS TINY MAGNETIZED REGIONS — UP TO 10 MILLION BYTES PER DISK!



A BIG COMPUTER SYSTEM USUALLY HAS MULTIPLE DISK DRIVES, WITH PHONOGRAPH-ARMLIKE READ/WRITE HEADS DARTING BACK AND FORTH ACROSS THE WHIRLING PLATTERS.

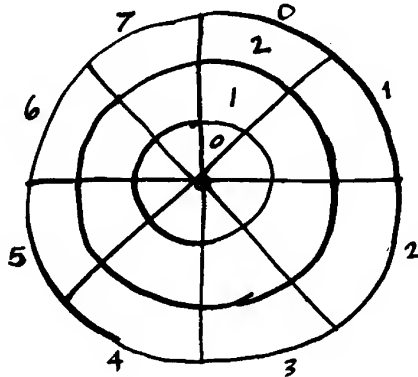


ARE SMALL, LOW-COST MAGNETIC DISKS MADE OF PLASTIC. THEY ALWAYS STAY IN THEIR JACKETS, BECAUSE A SPECK OF DUST CAN CREATE A MONSTER GLITCH!



OTHER, MORE EXOTIC MASS STORAGE TECHNOLOGIES INCLUDE BUBBLE MEMORIES, CHARGE-COUPLED DEVICES, AND OPTICAL DISKS READ BY LASERS.

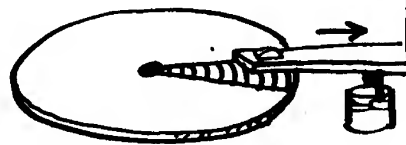
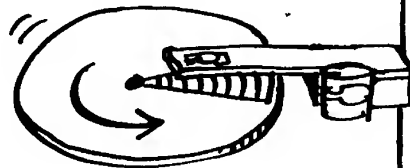
LIKE INTERNAL MEMORY, MASS STORAGE MUST BE ORGANIZED, OR "FORMATTED." TAKE THE FLOPPY DISK FOR EXAMPLE:



FLOPPIES ARE FORMATTED INTO RINGS AND SECTORS — THREE RINGS AND EIGHT SECTORS, IN THIS VERY OVER-SIMPLIFIED DISK. (IT'S MORE LIKE 26 SECTORS AND 77 RINGS IN A GENUINE DISK.)

TO ACCESS A PARTICULAR BLOCK OF DATA, YOU SPECIFY THE RING NUMBER AND SECTOR NUMBER. THEN THE DISK DRIVE

- 1) SPINS THE DISK UNTIL THAT SECTOR LIES UNDER THE READ/WRITE HEAD
- 2) MOVES THE HEAD IN OR OUT TO THE PROPER RING.



THIS PROCESS TAKES MILLISECONDS — AN ETERNITY TO A COMPUTER!

SOME TYPICAL USES OF MASS STORAGE :

A GERBIL RANCHER,
USING A MICROCOMPUTER
TO IMPROVE PRODUCTIVITY,
BUYS THE APPROPRIATE
PROGRAMS (FROM GERBYTE,
INC.) STORED ON
FLOPPIES.



THE PHONE COMPANY
STORES IN BUBBLE
MEMORY THE MESSAGE:
"THE NUMBER YOU HAVE
REACHED IS NOT IN
SERVICE..."



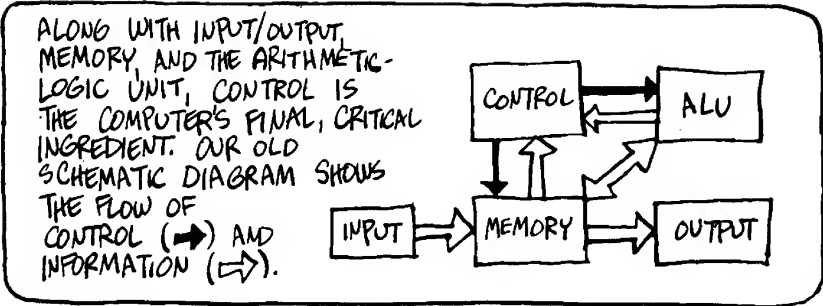
HM... SOUNDS
LIKE A
VERBAL
GERBIL...

WELL, YOU GET THE PICTURE...
NOW IT'S TIME TO MOVE ON...

GETTING EVERYTHING UNDER CONTROL

IN WHICH ALL
THE BLACK BOXES
ARE FINALLY SEEN
TO FIT TOGETHER...

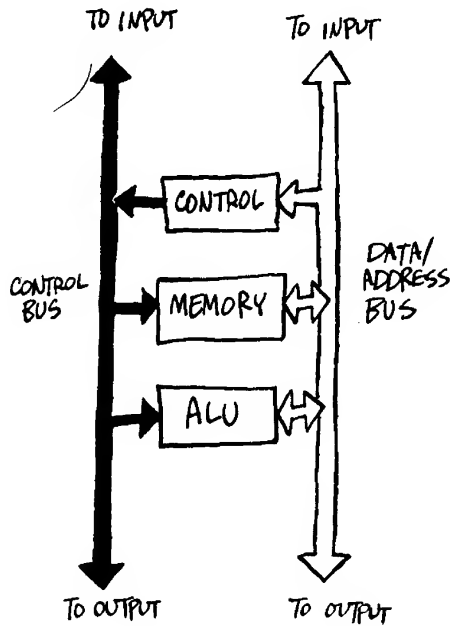




IT HELPS TO REDRAW THIS DIAGRAM IN A WAY THAT BETTER REFLECTS A GENUINE COMPUTER DESIGN KNOWN AS "BUS ARCHITECTURE."

THE VERTICAL ARROWS, REPRESENTING ELECTRICAL PATHWAYS A BYTE OR MORE WIDE, ARE THE BUSES.

ACCORDING TO SIGNALS PASSED ALONG THE CONTROL BUS, ADDRESSES AND DATA GET ON AND OFF THE DATA/ADDRESS BUS, WITH THE PROVISIO THAT ONLY ONE "PASSENGER" CAN RIDE THE BUS AT A TIME.

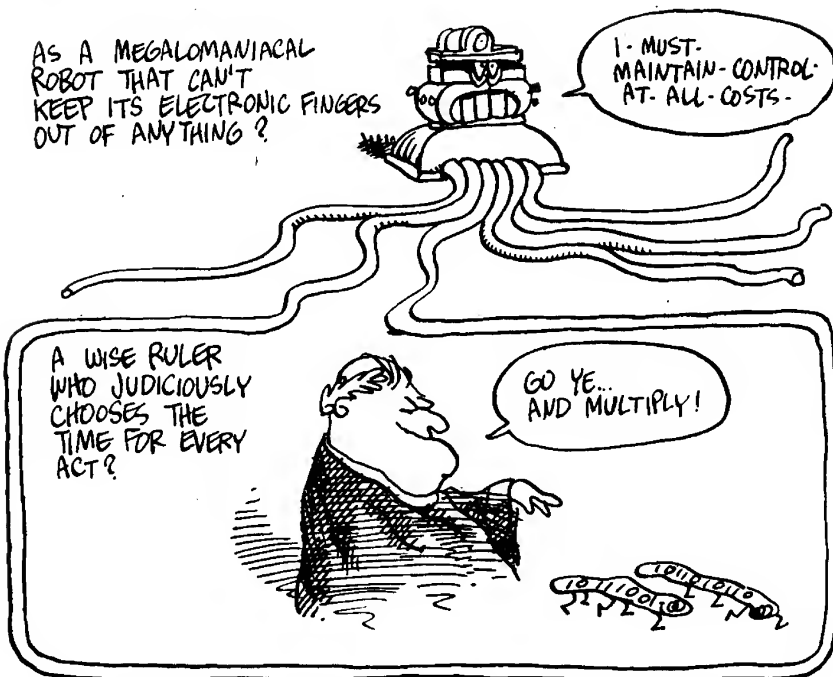


NOTE THAT ALL THE ARROWS ON THE CONTROL BUS POINT AWAY FROM THE CONTROL SECTION.

HOW ARE WE TO IMAGINE THIS CONTROL, FROM WHICH ALL DARK ARROWS POINT AWAY??

AS A MEGALOMANIACAL ROBOT THAT CAN'T KEEP ITS ELECTRONIC FINGERS OUT OF EVERYTHING?

I - MUST - MAINTAIN - CONTROL - AT - ALL - COSTS -



A WISE RULER WHO JUDICIOUSLY CHOOSES THE TIME FOR EVERY ACT?

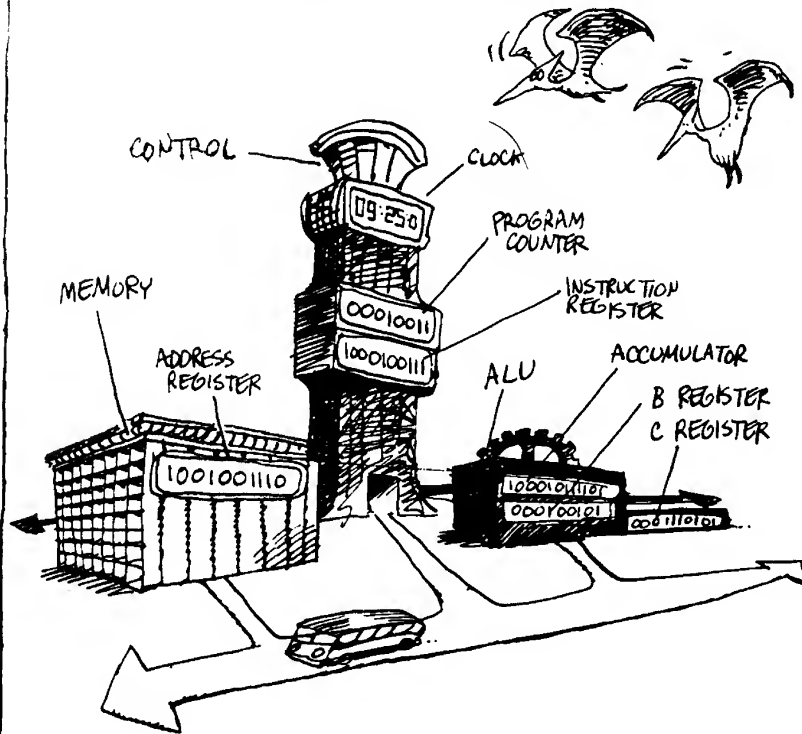
GO YE... AND MULTIPLY!



A RELENTLESS TYRANT WHO WIELDS A WHIP HAND OVER REBELLIOUS GLITCHES?







WELL, AT LEAST THE BUSES RUN ON TIME

LIKE ANYONE ELSE, CONTROL REVEALS ITS CHARACTER BY ITS BEHAVIOR... SO LET'S FOLLOW WHAT HAPPENS IN THIS OVERSIMPLIFIED COMPUTER, WHICH FLESHES OUT THE DIAGRAM OF TWO PAGES BACK WITH SOME ESSENTIAL COUNTERS AND REGISTERS.



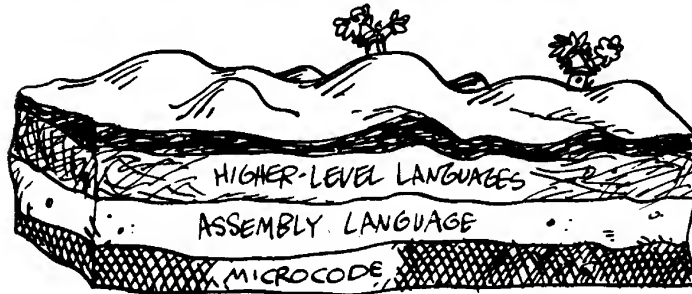
THIS IS A MINIMAL COLLECTION OF EQUIPMENT. A TYPICAL COMPUTER HAS MORE REGISTERS AND COUNTERS, BUT ALL COMPUTERS HAVE THE ONES SHOWN HERE.

HERE'S WHAT THEY'RE FOR:

<p>PROGRAM COUNTER: TICKS OFF THE INSTRUCTIONS ONE BY ONE.</p>  <p>Two... Two...?</p>	<p>INSTRUCTION REGISTER: HOLDS AN ENCODED VERSION OF THE INSTRUCTION BEING PERFORMED.</p>  <p>"BOIL SPAGHETTI TEN MINUTES."</p>
<p>ADDRESS REGISTER: HOLDS THE ADDRESS OF WHATEVER IS TO ENTER OR LEAVE MEMORY.</p>  <p>GET ME BYTE *0101!</p>	<p>ACCUMULATOR: THE ALU'S MAIN REGISTER, KEEPING A RUNNING TOTAL OF ALU OPERATIONS.</p>  <p>COULDN'T SOLVE 1+1 WITHOUT IT!</p>
<p>B REGISTER: AN AUXILIARY REGISTER TO HOLD NUMBERS ON THEIR WAY TO ALU.</p>  <p>LIKE A MOTEL THAT RENTS ROOMS BY THE MICROSECONDS!</p>	<p>C REGISTER: HOLDS DATA ON THE WAY TO OUTPUT.</p>  <p>IS THERE CONTROL IN THE OUTSIDE WORLD?</p>

➤ IN FACT, CONTROL SPENDS MOST OF ITS TIME JUST MOVING THE CONTENTS OF THESE REGISTERS AROUND!

TO SEE HOW CONTROL WORKS, LET'S FOLLOW WHAT HAPPENS WHEN THE COMPUTER **ADDS TWO NUMBERS**—OUR VERY FIRST PROGRAM!

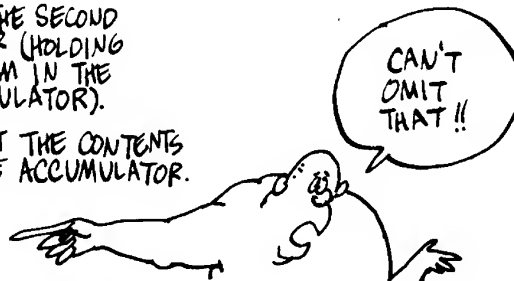


LIKE EVERYTHING ABOUT COMPUTERS, PROGRAMS CAN BE DESCRIBED AT VARIOUS LEVELS. WE BEGIN WITH

ASSEMBLY LANGUAGE,

WHICH SPECIFIES THE COMPUTER'S ACTUAL MOVES, BUT OMITTS THE FINE DETAILS. AT THIS LEVEL, HERE'S HOW TO ADD TWO NUMBERS:

0. LOAD THE FIRST NUMBER INTO THE ACCUMULATOR.
1. ADD THE SECOND NUMBER (HOLDING THE SUM IN THE ACCUMULATOR).
2. OUTPUT THE CONTENTS OF THE ACCUMULATOR.
3. HALT.



TO EXPRESS THIS IN PROPER ASSEMBLY LANGUAGE, WE MUST SPECIFY THE PRECISE LOCATION IN MEMORY OF THE TWO NUMBERS TO BE ADDED, AND CONDENSE THE WORDY STATEMENTS INTO MNEMONIC* ABBREVIATIONS. SUPPOSE, FOR EXAMPLE, THAT THE NUMBERS ARE STORED AT ADDRESSES 1E AND 1F (HEXADECIMAL). OUR PROGRAM BECOMES:

A TRUE ASSEMBLY LANGUAGE PROGRAM!

0. LDA 1E	("LOAD ACCUMULATOR WITH CONTENTS OF 1E")
1. ADD 1F	("ADD CONTENTS OF 1F")
2. OUT	("OUTPUT CONTENTS OF ACCUMULATOR.")
3. HALT	

*MNEMONIC = MEMORY-AIDING



IN GENERAL, ASSEMBLY-LANGUAGE STATEMENTS HAVE TWO PARTS:

THE OPERATOR, WHICH DESCRIBES THE STEP TO BE PERFORMED

THE OPERAND, WHICH GIVES THE ADDRESS ON WHICH THE OPERATOR ACTS


AS IN: PERFORM AN APPENDECTOMY ON THE RESIDENT OF 151 FIRST STREET!

LDA 1E

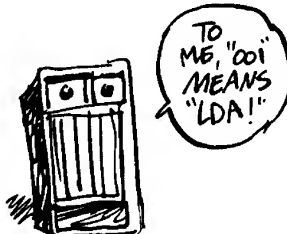
NOTE HOWEVER! SOME OPERATORS DON'T NEED AN EXPLICIT OPERAND. "OUT", FOR INSTANCE, IS UNDERSTOOD TO APPLY TO THE ACCUMULATOR.

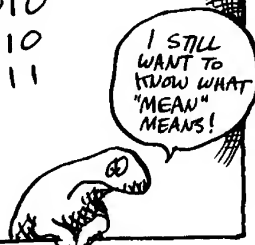


NOW THAT WE HAVE AN ASSEMBLY-LANGUAGE PROGRAM, HOW DO WE FEED IT TO THE MACHINE — WHICH ONLY UNDERSTANDS 0'S AND 1'S ?

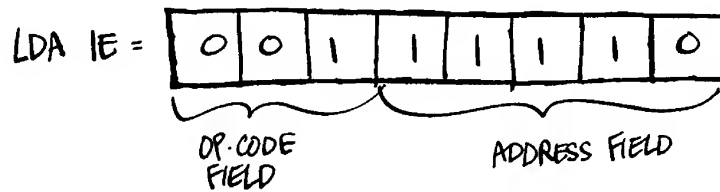
 THE ANSWER IS CLEAR: WITHIN THE MACHINE, EACH OPERATOR IS ENCODED AS A STRING OF BITS CALLED ITS "OP-CODE." SOME SIMPLE SAMPLES:

OPERATOR	OP-CODE
LDA	001
ADD	010
OUT	110
HALT	111





THEN A MACHINE INSTRUCTION CONSISTS OF AN OP-CODE SEGMENT, OR "FIELD," FOLLOWED BY AN ADDRESS FIELD GIVING THE OPERAND IN BINARY:



SO HERE'S OUR PROGRAM TRANSLATED INTO MACHINE LANGUAGE:

0. LDA 1E	001 11110
1. ADD 1F	010 11111
2. OUT	110 XXXXX
3. HALT	111 XXXXX

ANY 5 BITS ARE O.K. FOR THESE ADDRESS FIELDS, AS THEY'LL BE IGNORED!

NOW

(ASSUMING AN INPUT DEVICE)

THE PROGRAM STEPS ARE READ INTO CONSECUTIVE MEMORY ADDRESSES, BEGINNING WITH 0. THE CONTENTS OF MEMORY ARE THEN

ADDRESS	CONTENTS
0	001 11110
1	010 11111
2	110 00000
3	111 00000

NOTE THAT THE PROGRAM STEP NUMBER IS THE ADDRESS WHERE IT'S STORED!

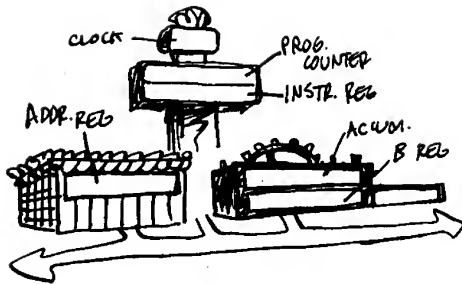


AND WE ALSO NEED TO ENTER THE DATA: THE TWO NUMBERS TO BE ADDED. ANY TWO NUMBERS WILL DO, SAY 5 AND 121. THEY GO IN ADDRESSES 1E AND 1F:

1E	00000101
1F	01111001



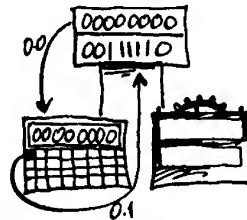
HOW CAN THE COMPUTER DISTINGUISH DATA FROM INSTRUCTIONS? BY ASSUMING EVERYTHING IS AN INSTRUCTION, UNLESS INSTRUCTED TO DO OTHERWISE!!



ONCE THE PROGRAM IS STORED, CONTROL CAN BEGIN EXECUTION, IN A SERIES OF EVEN MORE PRIMITIVE STEPS CALLED MICROINSTRUCTIONS, ONE MICROINSTRUCTION OCCURRING WITH EACH CLOCK PULSE. ARE YOU READY FOR THE GORY DETAILS?

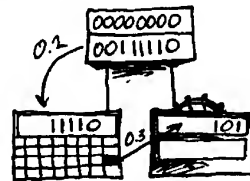
CONTROL BEGINS BY **FETCHING** THE FIRST INSTRUCTION. IT—

- 0.0. MOVES CONTENTS OF PROGRAM COUNTER (00000000 TO BEGIN WITH) TO ADDRESS REGISTER
- 0.1 MOVES CONTENTS OF THAT MEMORY ADDRESS TO INSTRUCTION REGISTER



THE INSTRUCTION REGISTER NOW HOLDS THE FIRST INSTRUCTION. CONTROL "READS" IT AND—

- 0.2. MOVES THE INSTRUCTION REGISTER'S ADDRESS FIELD TO ADDRESS REGISTER
- 0.3. MOVES CONTENTS OF THAT MEMORY ADDRESS TO ACCUMULATOR



THE ACCUMULATOR IS NOW LOADED WITH THE FIRST PIECE OF DATA. ONE MICROINSTRUCTION REMAINS:

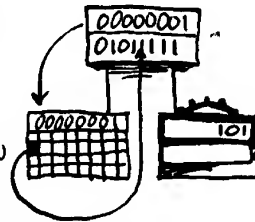
- 0.4 INCREMENT PROGRAM COUNTER



A BIT CONFUSED? LET'S GO THROUGH IT AGAIN WITH THE NEXT STEP, ADD.

AGAIN CONTROL BEGINS WITH A "FETCH PHASE":

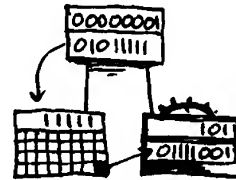
1.0 MOVE CONTENTS OF PROGRAM COUNTER (NOW 00000001) TO ADDRESS REGISTER



1.1 MOVE CONTENTS OF THAT ADDRESS TO INSTRUCTION REGISTER

THE INSTRUCTION IN THE INSTRUCTION REGISTER, 01011111, CAUSES CONTROL TO:

1.2 MOVE ADDRESS FIELD FROM INSTRUCTION REGISTER TO ADDRESS REGISTER



1.3 MOVE CONTENTS OF THAT MEMORY ADDRESS TO B REGISTER

1.4 SIGNAL THE ALU TO ADD AND PUT THE SUM IN ACCUMULATOR



AGAIN, THERE'S ONE MORE STEP:

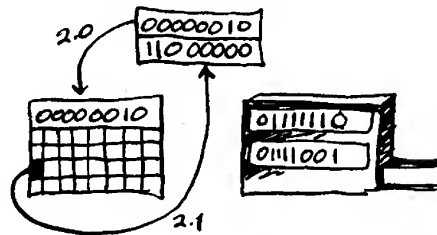
1.5 INCREMENT PROGRAM COUNTER



AND FINALLY?

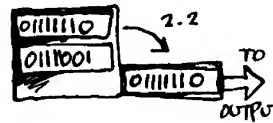
WELL, LUCKILY THE LAST TWO INSTRUCTIONS ARE EASIER:

2.0 AND 2.1 ARE THE SAME FETCH INSTRUCTIONS AS BEFORE, PUTTING INSTRUCTION 2 ("OUT") IN THE INSTRUCTION REGISTER:



THIS OP. CODE (110) CAUSES CONTROL TO —

- 2.2. MOVE CONTENTS OF ACCUMULATOR TO C REGISTER
- 2.3. INCREMENT PROGRAM COUNTER

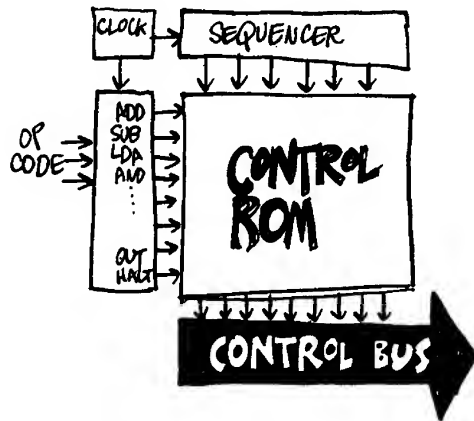


FINALLY, CONTROL FETCHES THE INSTRUCTION 111 ("HALT"), WHICH CAUSES CONTROL TO —

- 3.2 DO NOTHING



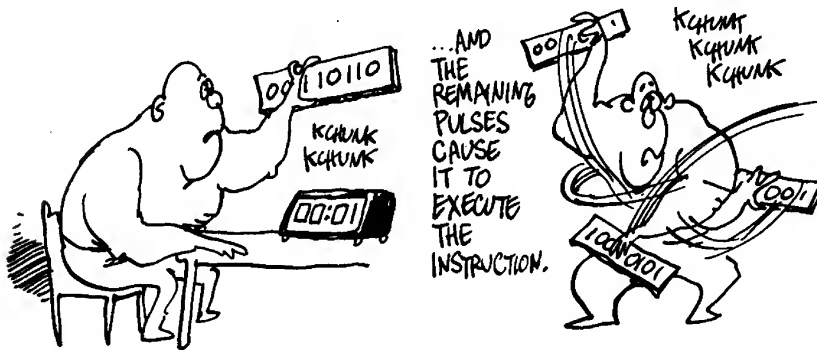
WITHOUT TOO MANY DETAILS, YOU CAN THINK OF CONTROL ROUGHLY LIKE THIS:



ITS INPUT CONSISTS OF CLOCK PULSES AND OP-CODES. ITS OUTPUT CONSISTS OF A SEQUENCE OF SIGNALS TO THE REGISTERS, COUNTERS, ALU, AND MEMORY.

THE "MICROPROGRAM," WHICH CONNECTS THE INPUTS TO THE PROPER OUTPUT COMBINATIONS, IS STORED IN A READ-ONLY MEMORY DEDICATED STRICTLY TO THIS PURPOSE.

THE FIRST COUPLE OF CLOCK PULSES CAUSE CONTROL TO FETCH AN INSTRUCTION...



IN REAL LIFE THE SITUATION IS MORE COMPLICATED IN DETAIL BUT THE SAME IN PRINCIPLE. THERE ARE MORE REGISTERS, AND OP-CODES ARE LONGER THAN THREE BITS, ALLOWING CONTROL TO RESPOND TO A MUCH LARGER SET OF INSTRUCTIONS. HERE'S THE INSTRUCTION SET OF A GENUINE PROCESSOR, THE MOTOROLA 6800.

ARITHMETIC

ADD
 ADD WITH CARRY
 SUBTRACT
 SUBTRACT WITH CARRY
 INCREMENT
 DECREMENT
 COMPARE
 NEGATE

LOGICAL

AND
 OR
 EXCLUSIVE OR
 NOT
 SHIFT RIGHT
 SHIFT LEFT
 SHIFT RIGHT ARITHMETIC
 ROTATE RIGHT
 ROTATE LEFT
 TEST

DATA TRANSFER

LOAD
 STORE
 MOVE
 CLEAR
 CLEAR CARRY
 CLEAR OVERFLOW
 SET CARRY
 SET OVERFLOW

BRANCH

BRANCH
 BRANCH IF ZERO
 BRANCH IF NOT ZERO
 BRANCH IF EQUAL
 BRANCH IF NOT EQUAL
 BRANCH IF CARRY
 BRANCH IF NO CARRY
 BRANCH IF POSITIVE
 BRANCH IF NEGATIVE
 BRANCH IF OVERFLOW
 BRANCH IF NO OVERFLOW
 BRANCH IF GREATER THAN
 BRANCH IF GREATER THAN OR EQUAL
 BRANCH IF LESS THAN
 BRANCH IF LESS THAN OR EQUAL
 BRANCH IF HIGHER
 BRANCH IF NOT HIGHER
 BRANCH IF LOWER
 BRANCH IF NOT LOWER

SUBROUTINE CALL

CALL SUBROUTINE

SUBROUTINE RETURN

RETURN FROM SUBROUTINE
 RETURN FROM INTERRUPT

MISCELLANEOUS

NO OPERATION
 PUSH
 POP
 WAIT
 ADJUST DECIMAL
 ENABLE INTERRUPT
 DISABLE INTERRUPT
 BREAK

ONE GROUP OF THESE INSTRUCTIONS DESERVES SPECIAL MENTION: THE **BRANCH, OR JUMP, INSTRUCTIONS.**

AS WE'LL SEE, THESE GIVE THE COMPUTER A LOT OF ITS "INTELLIGENCE." THEIR EFFECT IS TO **TRANSFER CONTROL** TO ANOTHER PART OF THE PROGRAM. THE SIMPLEST JUMP INSTRUCTION IS JUST PLAIN "JUMP," AS IN:

JMP 123



⇒ "JMP 123" CAUSES CONTROL TO ENTER 123 IN THE PROGRAM COUNTER... AND PROCEED WITH THE PROGRAM FROM THERE.

EVEN "SMARTER" ARE CONDITIONAL JUMPS. THEY TRANSFER CONTROL **IF** SOME CONDITION IS SATISFIED: FOR INSTANCE, "JUMP IF ZERO" MEANS JUMP IF THE ACCUMULATOR HOLDS 0.

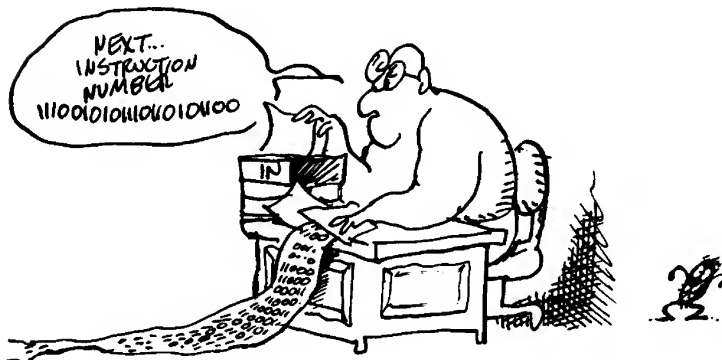
JZ 321



SO YOU SEE,
CONTROL IS
NO TYRANT
AT ALL. IT
ONLY DOES
WHAT IT'S TOLD—
COMPLETELY
AUTOMATICALLY!!



IF YOU REALLY WANT TO IMAGINE THE CONTROL SECTION'S
PERSONALITY, THINK OF A PERFECTLY EFFICIENT
BUREAUCRAT, ACTING IN STRICT OBEDIENCE TO THE
COMPUTER'S REAL BOSS: THE **PROGRAM!**



PART III

SOFTWARE

"GO TO...
GO TO..."
-SHAKESPEARE

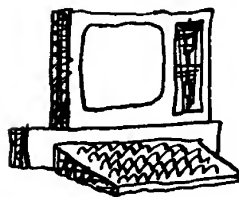
IF PROGRAMS REALLY RULE
THE COMPUTER, THEY DESERVE
A PROPER SCIENTIFIC NAME...
SOMETHING IN GREEK OR
LATIN, PREFERABLY...

TECHNICALCULUS?
REEMULA RATIONOCERUS?
CERHALONEURALGIA?



* * * * *

BUT THAT'S NOT HOW IT IS IN COMPUTER SCIENCE...
INSTEAD, PROGRAMS IN GENERAL ARE CALLED **SOFTWARE**,
TO DISTINGUISH THEM FROM THE CIRCUIT BOARDS, CATHODE
RAY MONITORS, DISK DRIVES, KEYBOARDS, AND OTHER
ITEMS OF COMPUTER **HARDWARE**.



HARDWARE



SOFTWARE



TUPPERWARE

LIKEWISE THERE'S A DIFFERENCE BETWEEN THE IMAGE OF
HARDWARE AND SOFTWARE WORKERS —



HARDWARE TYPES ARE
ENGINEERS... INTO GADGETS...
MOSTLY MEN... BOUND
BY THE LAWS OF
PHYSICS...

PROGRAMMERS HAVE NO TOOL BUT THEIR BRAINS... THEY'RE
MORE OFTEN WOMEN... SUPPOSED TO BE SOLITARY
DREAMERS WHOSE IDEAS HAVE NOTHING TO DO WITH
THE LAWS OF PHYSICS!!



GRONK!

PROGRAMS THESE DAYS ARE
SO COMPLEX THAT NO
ONE PERSON CAN
UNDERSTAND THEM — SO
THESE LOVERS HAVE TO
WORK IN TEAMS — A
SPECTACLE I LEAVE TO
THE READER'S IMAGINATION...



WHILE ADA
LOVELACE WAS
THE ORIGINAL
PROGRAMMER,
THE FIRST PERSON
TO PROVE THE
FULL POWER OF
SOFTWARE WAS

**ALAN
TURING**
(1912-1954)



TURING, WHO ENJOYED LONG-DISTANCE RUNNING BACK
WHEN THAT WAS CONSIDERED WEIRD, PROBABLY WENT
INTO COMPUTERS TO SHRINK THE SIZE OF HIS
JOGGING CLOCK.

IT'S
MAKING
MY
STRIDE
LOPSIDED!



IN 1936
HE DREAMED
UP THE
TURING
MACHINE...

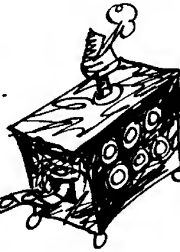
TURING MACHINES
AREN'T REAL MACHINES...
THEY'RE ABSTRACT
MACHINES, EXISTING ONLY
IN THEORY...



A
SOFTWARE
ENGINEER'S
DREAM —
NO
HARDWARE!

ROUGHLY SPEAKING, A
TURING MACHINE IS AN
INPUT-OUTPUT DEVICE: A
BLACK BOX THAT READS A
SEQUENCE OF 0'S AND 1'S.

THE OUTPUT
DEPENDS ONLY ON THE
PRESENT INPUT (0 OR 1)
AND THE PREVIOUS OUTPUT.



THE NATURE OF
THE OUTPUT IS
UNIMPORTANT.

THE MAIN THING IS
THAT THE CHANGES
FROM ONE OUTPUT
STATE TO THE NEXT
ARE GIVEN BY
DEFINITE RULES,
CALLED THE
TRANSITION
RULES.

THE REASON TURING
MACHINES ARE IMPORTANT
IS THAT THEY ARE A
WAY OF THINKING
PHYSICALLY ABOUT LOGIC.
ANY WELL-DEFINED, STEP-BY-STEP
LOGICAL PROCEDURE
CAN BE EMBODIED
IN SOME TURING MACHINE.



TUTT
THERE'S A
TURING MACHINE
THAT CAN
ADD!

*FOR DETAILS, SEE J. WEIZENBUM'S COMPUTER POWER AND HUMAN REASON,
CHAPTER 2.

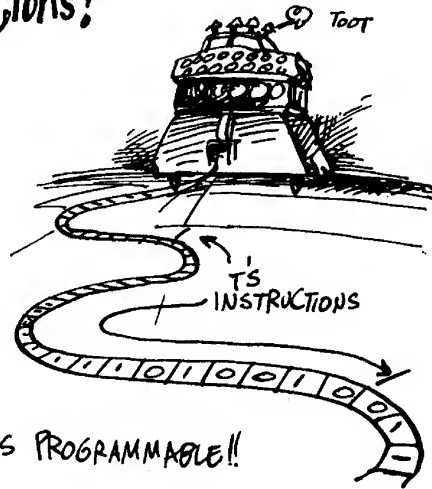
WHAT TURING PROVED:
IT'S THEORETICALLY
POSSIBLE TO
CONSTRUCT A
SINGLE TURING
MACHINE, THE
UNIVERSAL
TURING MACHINE,
WHICH CAN IMITATE
ALL OTHER
TURING MACHINES!!!



THE TRICK IS THAT THE UNIVERSAL TURING MACHINE CAN...

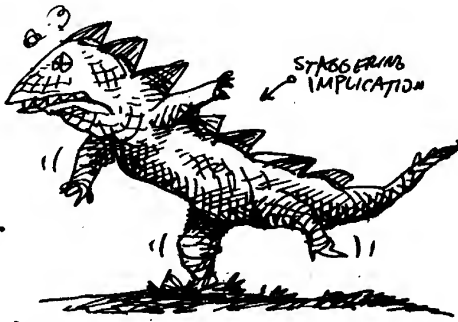
➔ **read instructions!**

THAT IS, TO MAKE THE
UNIVERSAL TURING MACHINE
(U) ACT LIKE MACHINE T,
YOU ENCODE T'S
TRANSITION RULES ONTO
U'S TAPE. AT EACH
STEP, U OBSERVES ITS
OWN INPUT, THEN
REFERS TO T'S
TRANSITION RULES TO
SEE WHAT TO DO.



➔➔ IN OTHER WORDS, U IS PROGRAMMABLE!!

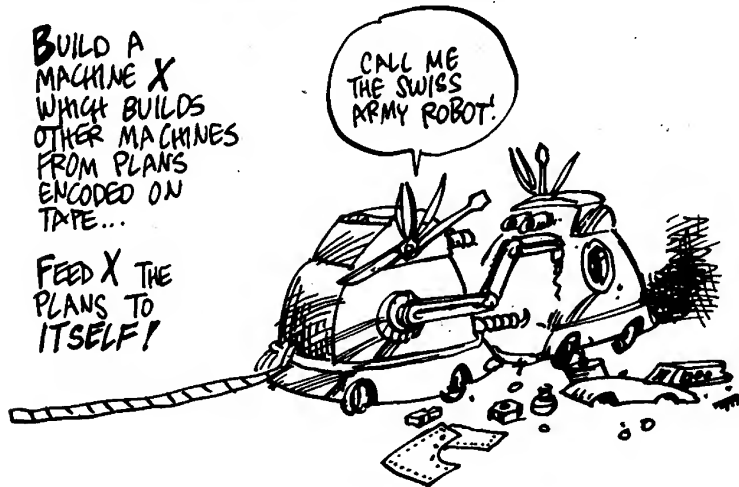
THE IMPLICATIONS ARE STAGGERING:
 A SINGLE, PROGRAMMABLE MACHINE CAN PERFORM ANY WELL-DEFINED, STEP-BY-STEP LOGICAL PROCEDURE. (REMEMBER, TURING SAW THIS TEN YEARS BEFORE A REAL COMPUTER WAS BUILT.)



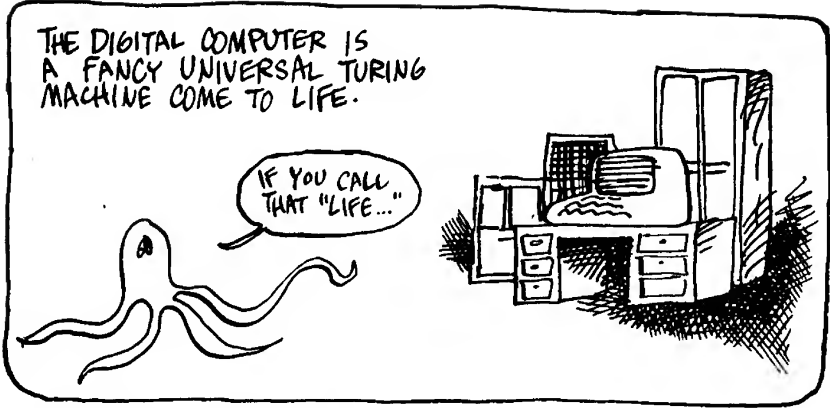
JOHN VON NEUMANN CARRIED TURING'S IDEAS A STEP FURTHER. VON NEUMANN REALIZED THAT ONE COULD:

BUILD A MACHINE X WHICH BUILDS OTHER MACHINES FROM PLANS ENCODED ON TAPE...

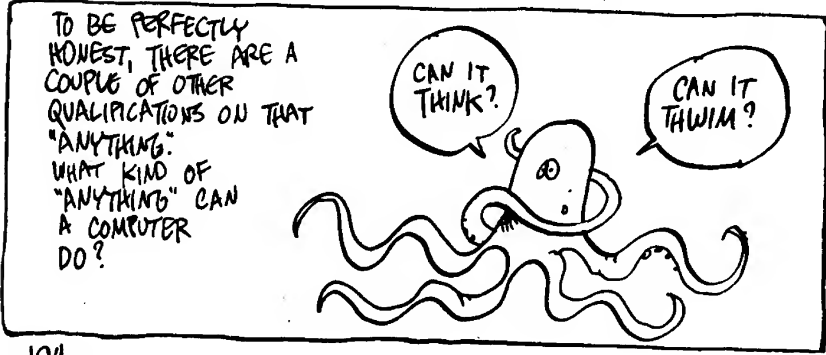
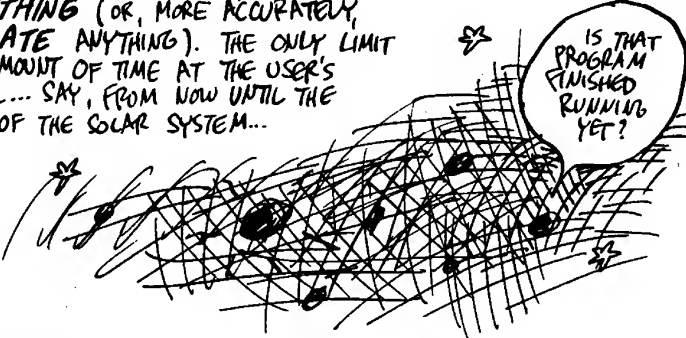
FEED X THE PLANS TO ITSELF!



➔ SELF-REPRODUCING MACHINES ARE POSSIBLE!!



THEREFORE, AS TURING PROVED, IT CAN DO ANYTHING (OR, MORE ACCURATELY, SIMULATE ANYTHING). THE ONLY LIMIT IS THE AMOUNT OF TIME AT THE USER'S DISPOSAL... SAY, FROM NOW UNTIL THE DEATH OF THE SOLAR SYSTEM...

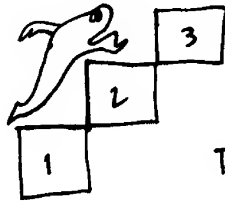


IN A WORD,
COMPUTERS DO

ALGORITHMS



AN ALGORITHM
IS SIMPLY
ANY WELL DEFINED,
STEP-BY-STEP
PROCEDURE: A
RECIPE, IF YOU
WILL!



STEP-BY-STEP,
MEANING EACH STEP
IS COMPLETED BEFORE
THE NEXT IS BEGUN.

WELL DEFINED,
MEANING EACH STEP
IS COMPLETELY
DETERMINED BY
CURRENT INPUT AND
THE RESULTS OF
PREVIOUS STEPS.
NO AMBIGUITY ALLOWED!



EXAMPLES OF ALGORITHMS:



↳ IT'S AN ALGORITHM BECAUSE I ALWAYS KNOW WHAT TO DO:

1. CHECK TO SEE IF WARHEADS ARE FALLING
2. IF YES, LIE DOWN + ENJOY!
3. IF NO, GO TO WORK.



LIKEWISE, ALGEBRAIC FORMULAS REPRESENT ALGORITHMS

$y = x^2 + 2x + 10$ MEANS -

- (1) INPUT A NUMBER x
- (2) MULTIPLY x TIMES ITSELF
- (3) MULTIPLY x TIMES 2
- (4) ADD THE RESULTS OF (2) AND (3)
- (5) ADD 10 TO THE RESULT OF (4)

IF YOU UNDERSTAND, LIE DOWN AND ENJOY YOURSELF!



EXAMPLES OF NON-ALGORITHMS:



ANOTHER?

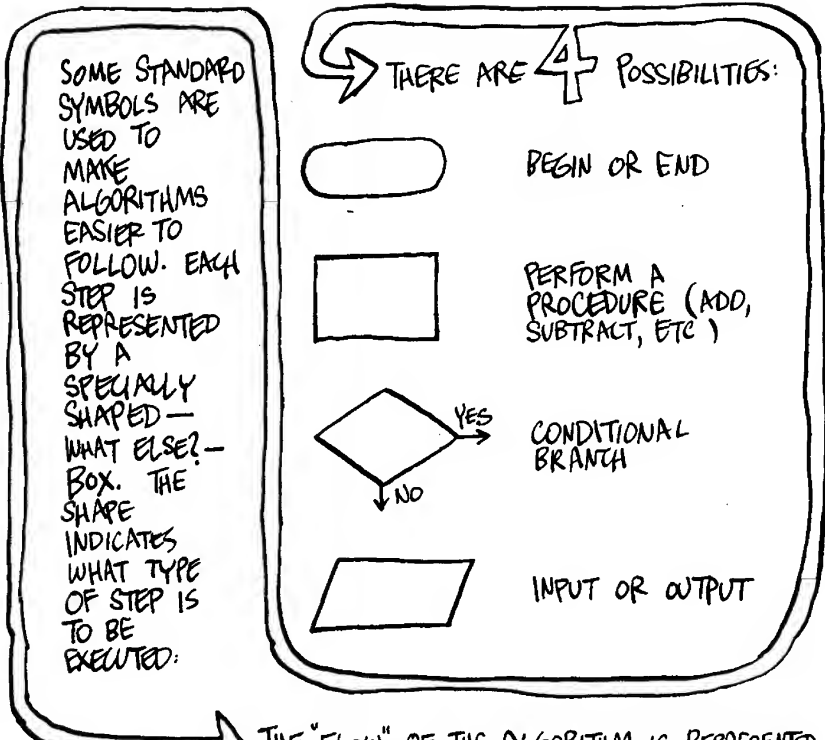
HOW ABOUT

$$y = x^2 ++ 2x - 10 ?$$

THIS IS NO ALGORITHM BECAUSE IT'S NOT EXPRESSED IN PROPER "ALGEBRAIC GRAMMAR." WE ASSIGN NO MEANING TO THE SYMBOLS "++".



IF YOU TRY TO MAKE A COMPUTER DO A NON-ALGORITHM, IT WILL JUST SIT THERE FLASHING ERROR MESSAGES!

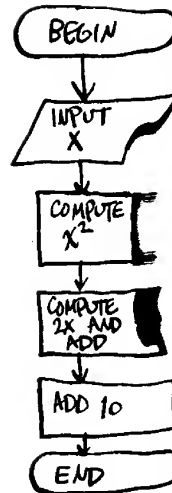
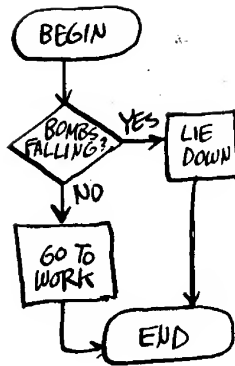


THE "FLOW" OF THE ALGORITHM IS REPRESENTED BY ARROWS →, AND WHEN ALL THE SYMBOLS ARE COMBINED, IT'S A

FLOW CHART



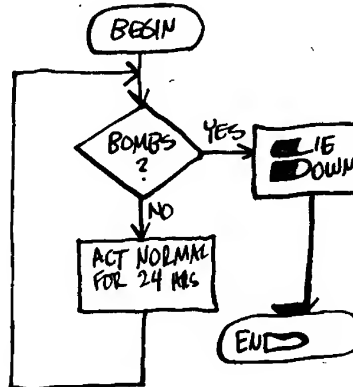
HERE ARE THE
FLOW CHARTS
OF THE
ALGORITHMS
FROM A
COUPLE OF PAGES
BACK:



IN BOTH ALGORITHMS, THE FLOW PROCEEDS
IN ONE DIRECTION, FROM START TO FINISH.

IT'S ALSO POSSIBLE FOR THE FLOW OF
ALGORITHMS TO JUMP FORWARD
OR BACKWARD. FOR EXAMPLE,
LET'S REWRITE THAT FIRST
ALGORITHM:

1. IF BOMBS ARE FALLING,
GO TO STEP 2. OTHERWISE,
GO TO STEP 4.
2. LIE DOWN AND ENJOY!
3. GO TO STEP 6.
4. LEAD A NORMAL LIFE
FOR 24 HOURS
5. GO TO STEP 1
6. END



YOU MAY FIND THE FLOW CHART
EASIER TO GRASP THAN THE
WRITTEN "PROGRAM." NOTE THAT
IT MAY CONTINUE INDEFINITELY!!

FLOW CHARTS ARE USEFUL
IN HELPING TO DESIGN
ALGORITHMS - SIMPLE ONES,
ANYWAY - AND DESIGNING
ALGORITHMS IS WHAT
COMPUTER PROGRAMMING
IS ALL ABOUT !!




THE FIRST STEP IN WRITING ANY PROGRAM IS TO
ANALYZE THE JOB TO BE DONE, AND SEE
HOW TO DO IT ALGORITHMICALLY!

FAILURE TO
THINK
ALGORITHMICALLY
HAS CAUSED
MANY
SOFTWARE
NIGHTMARES !!
MOST
SOFTWARE
DESIGNERS
HAVE
HORROR STORIES
ABOUT
CUSTOMERS
WHO DIDN'T
KNOW
EXACTLY
WHAT THEY
WANTED !!

A cartoon illustration of a hand pointing at a speech bubble. The speech bubble contains the text: "YES... THAT INFORMATION SHOULD GO INTO MY FILES... OR MAYBE NOT... THE VICE PRESIDENT'S ARE JUST AS GOOD... OR MAYBE THE TREASURER'S...". To the right, another speech bubble says "NO! NO!".

LET'S TRY A COUPLE MORE EXAMPLES...
A LITTLE MORE LIKE WHAT A COMPUTER
MIGHT ACTUALLY BE ASKED TO DO...

"ROOMMATE RECEIPTS"



TWO ROOMMATES, LISA AND SOPHIE,
SHARE THEIR MEALS. THEY
BOTH SHOP FOR FOOD AND SAVE
THEIR RECEIPTS. AT THE END OF
THE MONTH, THEY WANT TO
KNOW WHO OWES WHOM HOW
MUCH.

**"MULTIPLE
PLUG-INS"**

THIS ONE ASKS THE
COMPUTER TO
EVALUATE THE
EXPRESSION

$$x^2 + 2x + 10$$

NOT JUST AT ONE
VALUE OF X, BUT
FOR MANY VALUES,
NAMELY

$x=0, 0.1, 0.2, 0.3,$
...AND SO ON...
UP TO 2.0.



FOR "ROOMMATE RECEIPTS"
WE REASON LIKE SO:

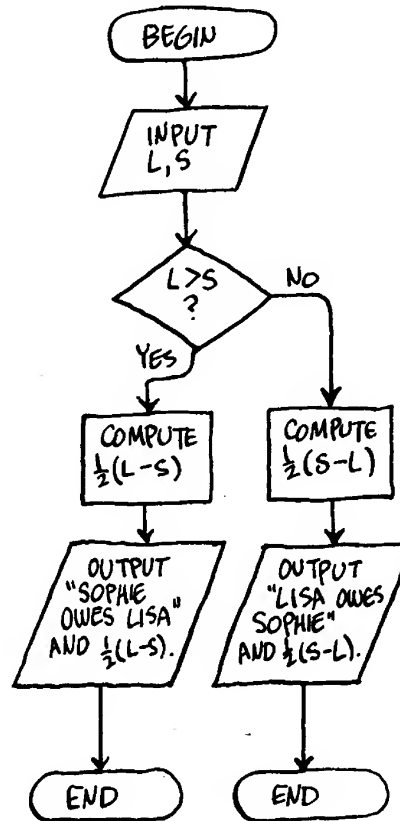
LET S = SOPHIE'S EXPENSES
 L = LISA'S EXPENSES

THEN THE TOTAL EXPENSE
IS $S+L$, AND EACH
ROOMMATE'S SHARE IS
 $\frac{1}{2}(S+L)$.

IF LISA OUTSPENT SOPHIE,
SO $L > S^*$, THEN
SOPHIE OWES LISA
 $\frac{1}{2}(S+L) - S$, OR
 $\frac{1}{2}(L-S)$.

OTHERWISE (WHEN $S \geq L^*$),
LISA OWES SOPHIE
 $\frac{1}{2}(S-L)$.

THE ALGORITHM'S OUTPUT
IS TO TELL US WHO
OWES WHOM AND
HOW MUCH.



* $>$ MEANS "IS GREATER THAN"; \geq MEANS "IS GREATER THAN OR EQUAL TO";
 $<$ MEANS "IS LESS THAN"; \leq MEANS "IS LESS THAN OR EQUAL TO".

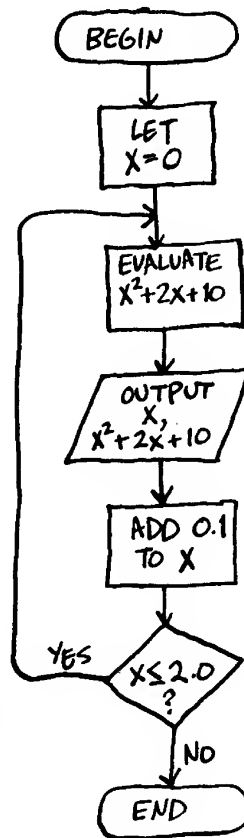
IN "MULTIPLE PLUG-INS," WE WANT TO EVALUATE A SINGLE EXPRESSION, $x^2 + 2x + 10$, REPEATEDLY AT DIFFERENT VALUES OF x (NAMELY 0.0, 0.1, 0.2, ..., 1.9, 2.0)

THE CORE OF THE ALGORITHM WILL BE THIS LOOP:

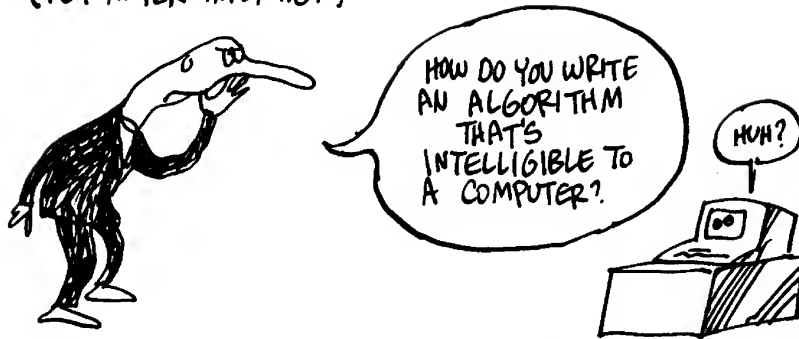
1. PLUG THE CURRENT VALUE OF x INTO $x^2 + 2x + 10$
2. PRINT THE RESULT
3. NEXT x
4. RETURN TO STEP 1.

WE ALSO HAVE TO SPECIFY WHAT x TO START WITH, WHEN TO STOP, AND HOW TO COMPUTE "NEXT x ."

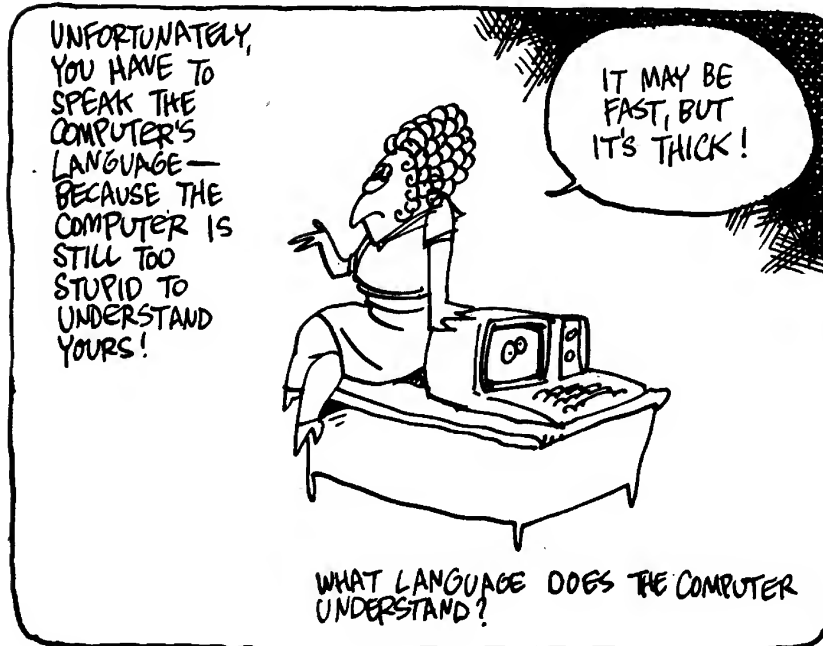
NOTE HOW THE FLOW CHART SHOWS HOW THE PROGRAM LOOPS BACK, PLUGGING IN SUCCESSIVE VALUES OF x UNTIL x EXCEEDS 2.



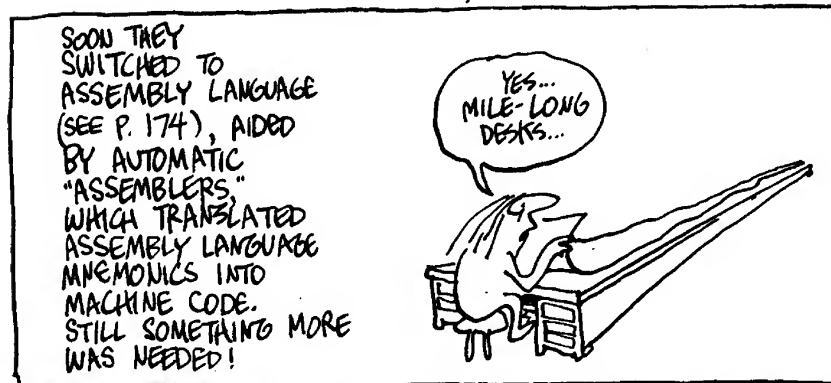
NOW THE #738 QUESTION:
(\$64 AFTER INFLATION):



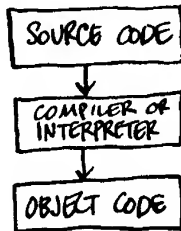
IN OTHER WORDS, HOW DO YOU PROGRAM A COMPUTER?



AT THE VERY BEGINNING, PROGRAMMERS WROTE DIRECTLY IN "MACHINE LANGUAGE" — BINARY CODE. THIS WAS OBVIOUSLY A HEADACHE!

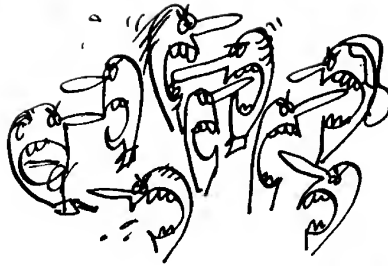


AND FINALLY,



THE **HIGHER-LEVEL** PROGRAMMING LANGUAGES WERE INVENTED. THESE CONTAIN FAMILIAR ENGLISH-LIKE COMMANDS, SUCH AS "PRINT," "READ," AND "DO," WHICH ARE TRANSLATED INTO MACHINE LANGUAGE BY COMPLEX PROGRAMS CALLED COMPILERS OR INTERPRETERS. HIGHER-LEVEL PROGRAMS ARE SOMETIMES CALLED "SOURCE CODE," AND THE MACHINE-LANGUAGE TRANSLATION IS CALLED "OBJECT CODE."

THE FIRST HIGHER-LEVEL LANGUAGE WAS **FORTRAN** ("FORMULA TRANSLATOR"), WHICH MADE ITS DEBUT IN THE EARLY 1950'S. SINCE THEN, LITERALLY HUNDREDS OF LANGUAGES HAVE BEEN WRITTEN, EACH WITH ITS OWN ARMY OF RABID DEVOTEES!



WE'RE GOING TO TAKE A QUICK LOOK AT **BASIC** — BEGINNER'S ALL-PURPOSE SYMBOLIC INSTRUCTION CODE. BASIC IS EASY TO LEARN AND WIDELY USED, DESPITE CRITICISM (ESPECIALLY BY PASCAL ADMIRERS) THAT IT PROMOTES "BAD PROGRAMMING HABITS."



WITH APOLOGIES TO PASCAL, THEN HERE'S A LITTLE BASIC...

BASIC BASIC



THERE ARE TWO
WAYS TO WRITE A
BASIC PROGRAM:
WITH PENCIL AND
PAPER, OR DIRECTLY
AT THE COMPUTER.

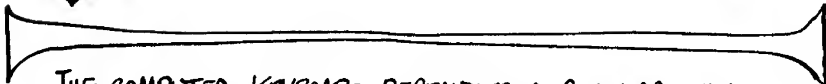
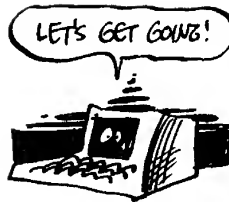
IT'S GOOD PRACTICE TO PLAN PROGRAMS ON PAPER FIRST, TO
WORK OUT THE ESSENTIAL IDEAS AND STRUCTURE, BUT
EVENTUALLY YOU MUST SIT DOWN AT THAT KEYBOARD!



SOME MACHINES
ARE READY FOR
BASIC AS
SOON AS YOU
TURN THEM ON.
OTHERS ONLY
BRING IT UP
ON COMMAND.
IF IN DOUBT,
ASK!



WHEN THE COMPUTER IS READY, IT GIVES YOU A "PROMPT" OF SOME KIND: THE WORD "READY" OR JUST THE SIGN ">".



THE COMPUTER KEYBOARD RESEMBLES A STANDARD TYPEWRITER'S "QWERTY" KEYBOARD... EXCEPT THAT AS YOU TYPE, CHARACTERS APPEAR ON THE CRT (CATHODE RAY TUBE) SCREEN, INSTEAD OF ON PAPER. TO GO TO THE NEXT LINE, HIT THE **RETURN** (**↵**) KEY. HERE'S A SIMPLE BASIC PROGRAM:

```
10 REM BASIC MULTIPLICATION
20 READ A, B
30 DATA 5.6, 1.1
40 LET C = A * B
50 PRINT "THE PRODUCT IS"; C
60 END
```

BASIC MATH:
 A+B } AS USUAL
 A-B }
 A*B... A TIMES B
 A/B... A DIVIDED BY B
 A^B... A TO THE BTH POWER

THE PROGRAM IS NOW STORED IN MEMORY. TO RUN IT, TYPE "RUN", FOLLOWED BY THE RETURN KEY. THE SCREEN DISPLAYS:

```
RUN
THE PRODUCT IS 6.16
```





> EVERY LINE BEGINS WITH A **LINE NUMBER** (10, 20, ...). EVERY LINE OF A BASIC PROGRAM MUST HAVE A NUMBER! IT'S WISE TO COUNT BY TENS, SO YOU CAN INSERT LINES LATER.

> THE FIRST LINE (10) IS A **REMARK**. REMARKS EXPLAIN THE PROGRAM BUT AREN'T EXECUTED BY THE COMPUTER. THE PREFIX "REM" IDENTIFIES REMARKS. WE MIGHT INSERT ONE HERE:

```
20 READ A, B
25 REM THESE ARE THE #S TO BE MULT'D
30 DATA 5.6, 1.1
```

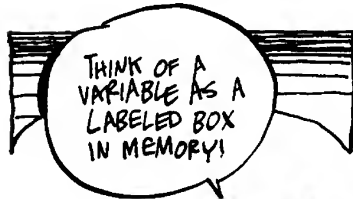
> PROGRAM **STATEMENTS** CONSIST OF INSTRUCTIONS ("LET", ETC), NUMBERS (5.6, 1.1), VARIABLES (A, B, C), **TEXT** ("THE PRODUCT IS"), AND PUNCTUATION.

```
50 PRINT "THE PRODUCT IS"; C
```

↑ ↑ ↑
QUOTES SPACES SEMICOLON

> EACH OF THESE HAS A PRECISE MEANING!

(NUMERICAL) VARIABLES



A NUMERICAL VARIABLE IN BASIC IS LIKE A VARIABLE IN ALGEBRA. IT ASSUMES A NUMERICAL VALUE, WHICH MAY VARY (BUT IT HAS ONLY ONE VALUE AT A TIME!). ONLY THESE SYMBOLS CAN BE USED AS VARIABLES:

A, B, C, D,	Z
A0, B0, ...	AND	...Z0
A1, B1, ...	EVERYTHING	...Z1
⋮	IN	⋮
A9, B9, ...	BETWEEN!	⋮
		...Z9

THERE ARE SEVERAL WAYS TO ASSIGN A VALUE TO A VARIABLE: ONE IS THE **READ = DATA** STATEMENT.

```
20 READ A, B
30 DATA 5.6, 1.1
```

COMMAS ARE ESSENTIAL!

THIS INSTRUCTS THE COMPUTER TO ASSIGN THE NUMERICAL VALUES IN THE **DATA** STATEMENT - IN ORDER - TO THE VARIABLES IN THE **READ** STATEMENT.

```
20 READ A, B, C
30 DATA 5.6, 1.1
```

THIS IS A **BUG!**

DISGUSTING!



ANOTHER WAY TO ASSIGN VALUES TO VARIABLES IS WITH

LET.


```
10 LET Q=6.5
20 LET R=2*Q
30 LET S=Q^2+R+10
```

MAKES R=13

MAKES S=
 $(6.5)^2 + 13 + 10$
 $= 65.25$

THE LET STATEMENT ASSIGNS THE VALUE ON THE RIGHT OF THE EQUALITY SIGN, "=", TO THE VARIABLE ON THE LEFT. THE RIGHT-HAND SIDE MAY BE A NUMBER, OR SOME MATHEMATICAL EXPRESSION INVOLVING OTHER VARIABLES — AS LONG AS THEY ALREADY HAVE VALUES!!

```
10 LET Q=6.5
20 LET Q=0.5*R
30 LET S=Q^2+R+10
```



HERE STATEMENT 20 DOES NOT ASSIGN ANY VALUE TO R, BECAUSE R IS NOT ON THE LEFT SIDE OF "=". IN FACT, IF R HASN'T BEEN ASSIGNED SOME VALUE EARLIER IN THE PROGRAM, THEN STATEMENT 20 GIVES Q AN INDETERMINATE VALUE! BUT—

```
10 LET M=0
20 LET M=M+1
30 LET M=M+1
```

MAKES M=1

MAKES M=2

THESE STRANGE-LOOKING STATEMENTS ARE PERFECTLY O.K.! "LET M=M+1" MEANS "ASSIGN TO THE VARIABLE M A VALUE EQUAL TO ITS CURRENT VALUE PLUS 1."

PRINT

THIS IS AN OUTPUT COMMAND, MEANING "DISPLAY ON THE SCREEN," NOT "PRINT ON PAPER."

WHAT CAN BE
PRINTED?



YOU CAN PRINT ANY TEXT:

```
10 PRINT "ANY NUKES TODAY?"  
RUN  
ANY NUKES TODAY?
```

QUOTATION MARKS
ESSENTIAL!

QUOTATION MARKS
REMOVED

PRINT A VARIABLE AND YOU GET
ITS VALUE:

```
10 LET X=77001  
20 PRINT X  
RUN  
77001
```

BUT—

```
10 LET X=77001  
20 PRINT "X"  
RUN  
X
```

QUOTATION MARKS MAKE
THE COMPUTER TREAT X
AS A TEXT.

PRINT A MATHEMATICAL EXPRESSION AND YOU GET ITS VALUE:

```
10 LET Z=1.5  
20 PRINT Z↑2 + 2*Z + 10  
RUN  
15.25
```

BECAUSE
 $(1.5)^2 + 2 \times 1.5 + 10$
 $= 2.25 + 3.0 + 10 = 15.25$

SEMICOLON (;)

A SEMICOLON AFTER A PRINT STATEMENT CAUSES THE NEXT PRINT STATEMENT TO DISPLAY ITS OUTPUT ON THE SAME LINE AND DIRECTLY AFTER THE FIRST ONE'S :

```
10 LET A = 1
20 PRINT "INFINITY IS MORE THAN";
30 PRINT A
RUN
INFINITY IS MORE THAN 1
```

IT'S O.K. TO ABBREVIATE THIS :

```
10 LET A = 1
20 PRINT "INFINITY IS MORE THAN"; A
RUN
INFINITY IS MORE THAN 1
```

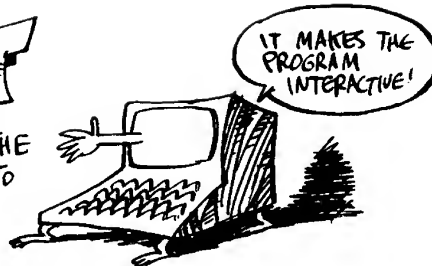
FOR EXAMPLE, WE COULD REWRITE THE PROGRAM ON P. 208.

```
10 REM BASIC MULTIPLICATION
20 READ A, B
30 DATA 5.6, 1.1
40 LET C = A * B
50 PRINT "THE PRODUCT OF"; A; "AND"; B; "IS"; C; "."
60 END
RUN
THE PRODUCT OF 5.6 AND 1.1 IS 6.16.
```

➤ THERE ARE ALSO SOME NIFTY TRICKS USING THE **COMMA** AND PRINT, BUT WE WON'T GET INTO IT...

INPUT

THIS STATEMENT ALLOWS THE USER TO ASSIGN VALUES TO VARIABLES WHILE THE PROGRAM IS RUNNING.



THE FORM OF THE STATEMENT:

```
INPUT A
```

WHEN THE PROGRAM RUNS AND REACHES AN INPUT STATEMENT, THE SCREEN DISPLAYS:

```
?
```

THIS INDICATES THAT THE PROGRAM HAS HALTED, WANTING INPUT. YOU TYPE SOME NUMBER (FOLLOWED BY "RETURN," AS ALWAYS!):

```
5.6
```

AND THE PROGRAM CONTINUES RUNNING. "INPUT" AND "PRINT" CAN BE USED IN COMBINATION TO LET YOU KNOW WHAT SORT OF INPUT IS EXPECTED:

```
10 BASIC DIVISION
20 PRINT "TYPE THE NUMERATOR."
30 INPUT N
40 PRINT "TYPE THE NON-ZERO DENOMINATOR."
50 INPUT D
60 PRINT N; "/"; D; "="; N/D
70 END
```

```
RUN
```

```
TYPE THE NUMERATOR.
```

```
? 5
```

```
TYPE THE NON-ZERO DENOMINATOR.
```

```
? 8
```

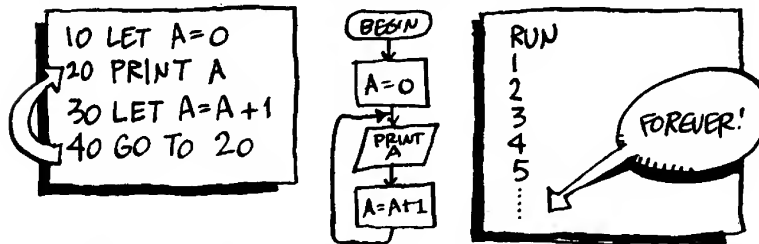
```
5/8 = 0.625
```

TYPED BY THE USER.

GO TO

THIS IS THE UNCONDITIONAL BRANCHING INSTRUCTION.

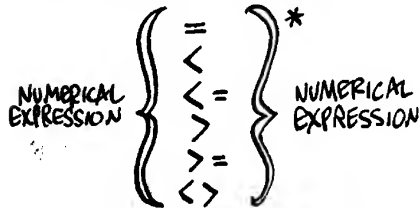
"GO TO (LINE NUMBER)" TRANSFERS CONTROL TO A LINE OTHER THAN THE NEXT. THE PROGRAM THEN CONTINUES FROM THERE, AS IN THIS ENDLESS LOOP:



IF-THEN

IS THE "SMART" CONDITIONAL JUMP.

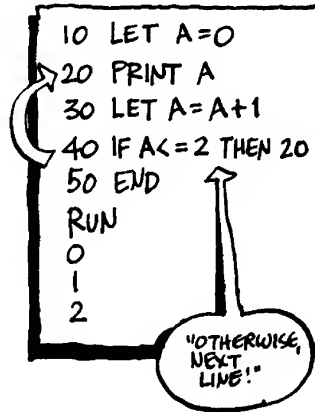
IT HAS THE GENERAL FORM IF (CONDITION) THEN (LINE NUMBER). THE CONDITION HAS THE FORM:



AS IN IF A <= B THEN 30

THIS ALWAYS INCLUDES THE UNSTATED INSTRUCTION, "OTHERWISE, GO TO THE NEXT LINE."

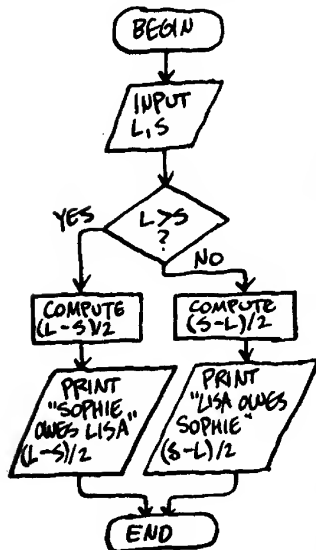
* < LESS THAN, <= LESS THAN OR EQUAL TO, > GREATER THAN, >= GREATER THAN OR EQUAL TO, <> DOES NOT EQUAL.



THIS IS ENOUGH TO WRITE BASIC PROGRAMS FOR THE TWO ALGORITHMS FROM P. 201:

ROOMMATE RECEIPTS

THE FLOW CHART:



NOW WE NEED A PROGRAM TO ROUND OFF THE HALF PENNY!



THE PROGRAM:

```

10 PRINT "LISA SPENT"
20 INPUT L
30 PRINT "SOPHIE SPENT"
40 INPUT S
50 IF L > S THEN 80
60 PRINT "LISA OWES SOPHIE"; (S-L)/2
70 GO TO 90
80 PRINT "SOPHIE OWES LISA"; (L-S)/2
90 END
  
```

SEE HOW "IF-THEN" AND "GO TO" ARE USED? IF $L > S$, THEN LINES 60 AND 70 ARE NOT EXECUTED. OTHERWISE, THEY ARE EXECUTED, AND LINE 80 ENSURES THAT LINE 80 IS SKIPPED.

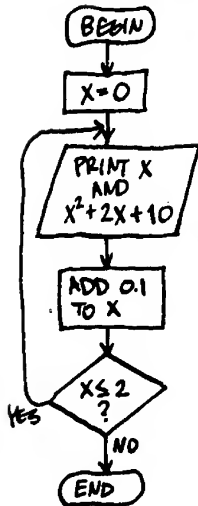
IF THE PROGRAM IS RUN:

```

RUN
LISA SPENT
? 93.75
SOPHIE SPENT
? 77.38
SOPHIE OWES LISA 8.185
  
```

MULTIPLE PLUGGINS

THE FLOW CHART:



THE PROGRAM:

```

10 REM LINE 20 PRINTS A HEADING
20 PRINT "X          Y"
30 LET X=0
40 LET Y=X^2+2*X+10
50 PRINT X;"          ";Y
60 LET X=X+0.1
70 IF X<=2 THEN 40
80 END
  
```

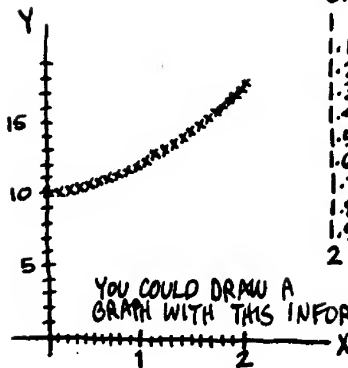
Annotations: "5 SPACES" points to the space between X and Y in line 20. "3 SPACES" points to the space between X and Y in line 50.

RUN

X
0
0.1
0.2
0.3
0.4
0.5
0.6
0.7
0.8
0.9
1
1.1
1.2
1.3
1.4
1.5
1.6
1.7
1.8
1.9
2

Y
10
10.21
10.44
10.69
10.96
11.25
11.56
11.89
12.24
12.61
13
13.41
13.84
14.29
14.76
15.25
15.76
16.29
16.84
17.41
18

GAPS A RESULT OF NOT USING ALL OF BASIC'S FORMATTING ABILITIES!



YOU COULD DRAW A GRAPH WITH THIS INFORMATION!

THE "MULTIPLE PLUG-INS" LOOP IS SO TYPICAL THAT ALL PROGRAMMING LANGUAGES HAVE SPECIAL COMMANDS JUST FOR SUCH REPETITIONS. IN BASIC, IT'S →

FOR
NEXT

THIS REPLACES THESE THREE LINES:

```
30 LET X=0
...
60 LET X=X+0.1
70 IF X<=2 THEN 30
...
```

WITH THESE TWO:

```
30 FOR X=0 TO 2 STEP 0.1
...
60 NEXT X
...
```

THE STATEMENT INITIALLY SETS THE VARIABLE EQUAL TO THE LOWER LIMIT, EXECUTES THE LINES UP TO "NEXT," INCREMENTS THE VARIABLE BY THE AMOUNT "STEP," AND REPEATS THE LOOP UNTIL THE UPPER LIMIT IS EXCEEDED.

A SIMPLE EXAMPLE:

```
10 FOR I = 1 TO 4
20 PRINT I*I
30 NEXT I
40 END
RUN
1
4
9
16
```

OMITTING "STEP" AUTOMATICALLY MAKES INCREMENT = 1.



NEXT
PRINT
I

PROBLEMS

PROBLEMS?
WHO HAS
PROBLEMS?



1. WHAT DOES THIS PROGRAM DO?

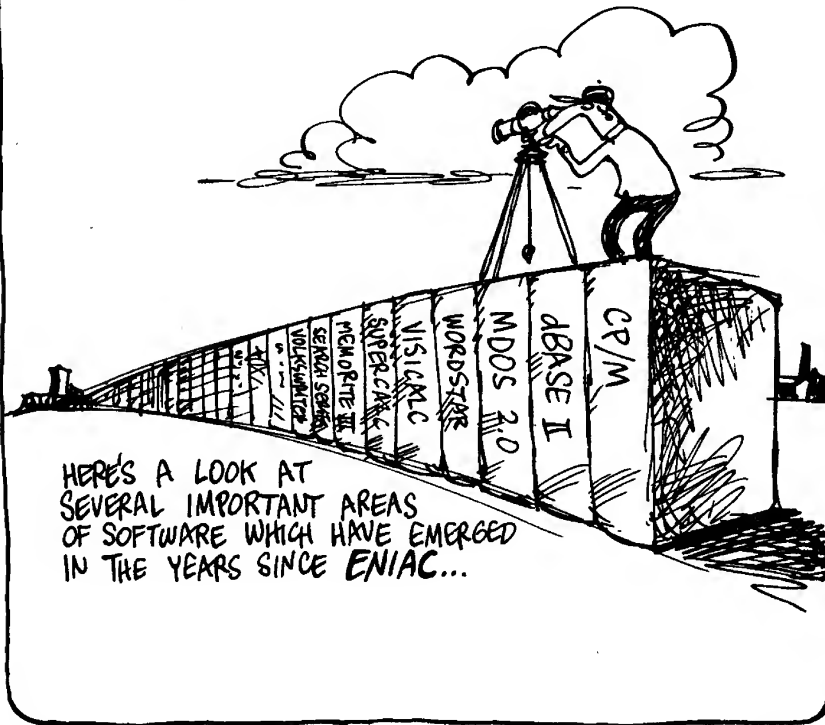
```
10 INPUT N
20 FOR I=1 TO N
30 PRINT I*I
40 NEXT I
50 END
```

2. REWRITE THE "MULTIPLE PLUG-INS" PROGRAM USING THE "FOR NEXT" STATEMENT.
3. WRITE A PROGRAM WHICH ADDS THE INTEGERS (WHOLE NUMBERS) FROM 1 TO 1,000,000. DITTO FROM 1 TO N, FOR ANY N.
4. IN THE FIBONACCI SEQUENCE 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ... EACH NUMBER IS THE SUM OF THE PREVIOUS TWO NUMBERS. WRITE A PROGRAM WHICH GENERATES THIS SEQUENCE.
5. READ ENOUGH OF A BASIC TEXTBOOK TO WRITE A "ROOMMATE RECEIPTS" PROGRAM FOR ANY NUMBER OF ROOMMATES.

THERE ARE PLENTY OF OTHER BASIC FEATURES, ENOUGH TO FILL ENTIRE BOOKS — AND IN FACT TONS OF BOOKS ON BASIC HAVE BEEN PUBLISHED.



SOFTWARE SURVEY



HERE'S A LOOK AT SEVERAL IMPORTANT AREAS OF SOFTWARE WHICH HAVE EMERGED IN THE YEARS SINCE ENIAC...

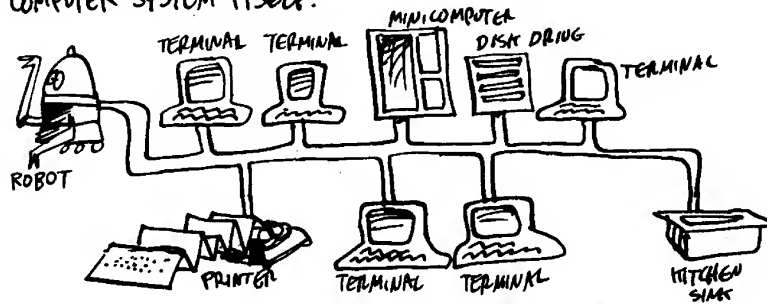
SYSTEMS SOFTWARE

PROGRAMS ARE COMMONLY DIVIDED INTO SYSTEMS SOFTWARE AND APPLICATIONS SOFTWARE.

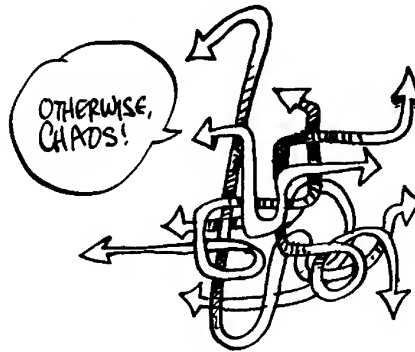


(WITH SOME GRAY AREA OF OVERLAP!)

APPLICATIONS SOFTWARE DOES "REAL WORLD" JOBS, WHILE SYSTEMS SOFTWARE EXISTS PURELY TO REGULATE THE COMPUTER SYSTEM ITSELF.



A SYSTEM TYPICALLY CONSISTS OF ONE OR MORE INPUT/OUTPUT DEVICES (TERMINALS, PRINTERS, CARD READERS, COMMUNICATIONS PORTS), PROCESSORS, MEMORY UNITS (MAIN AND MASS), AND WHO KNOWS WHAT ELSE. SOMETHING HAS TO COORDINATE IT ALL!



THE PROGRAM
THAT DOES IT
IS CALLED THE
**OPERATING
SYSTEM.**

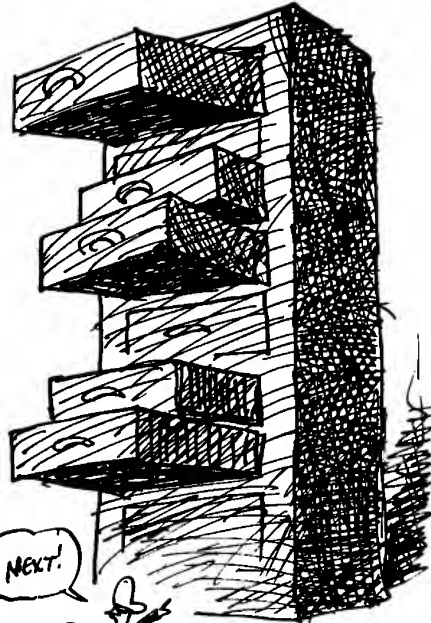
IF YOU THINK OF THE
COMPUTER'S CORE AS A
GIANT ELECTRONIC
FILING CABINET (WITH
A CALCULATOR ATTACHED),
THEN THE OPERATING
SYSTEM

★ CREATES THE STRUCTURE
OF THE FILES

★ MANAGES MEMORY
SO THAT DIFFERENT
FILES DON'T BUMP
INTO EACH OTHER

★ REGULATES ACCESS
TO THE FILES AND
THE MOVEMENT OF
INFORMATION TO AND FROM
OTHER PARTS OF THE
SYSTEM...

ETC!



BESIDES THE OPERATING
SYSTEM, SYSTEM SOFTWARE
INCLUDES OTHER PROGRAMS
"IN THE SYSTEM," SUCH AS
LOADERS (WHICH LOAD PROGRAMS
INTO MEMORY) AND COMPILERS
(WHICH TRANSLATE HIGHER-LEVEL
LANGUAGE INTO MACHINE CODE).


ALL INVISIBLE
TO THE USER!



DATA BASE MANAGEMENT



A DATA BASE IS JUST A BIG PILE OF INFORMATION: A LIBRARY'S CARD CATALOG, A BANK'S TRANSACTION RECORDS AND ACCOUNT BALANCES, AN AIRLINE'S FLIGHT SCHEDULES AND RESERVATIONS, POLICE FILES, STOCK EXCHANGE DATA — ALL ARE DATA BASES.

 A DATA BASE MANAGEMENT PROGRAM ORGANIZES, UPDATES, AND PROVIDES ACCESS TO THE DATA BASE.

IN THE CASE OF AN AIRLINE, FOR EXAMPLE, THE COMPUTER HAS TO BOOK RESERVATIONS, ASSIGN SEATS, ERASE RESERVATIONS WHEN THE CUSTOMER CANCELS, MAKE REASSIGNMENTS IF A FLIGHT IS CANCELED, PRINT THE TICKETS, AND PROVIDE ALL THE FLIGHT INFORMATION TO TRAVEL AGENTS — WORLDWIDE!!



WORD PROCESSING

A "PERSONAL" USE FOR COMPUTERS...

WORD PROCESSING SOFTWARE ALLOWS YOU TO WRITE, EDIT, AND FORMAT TEXT — ALL FROM THE SAME KEYBOARD. YOU CAN GO FROM FIRST TO FINAL DRAFT ELECTRONICALLY, BEFORE EVER PRINTING A WORD.

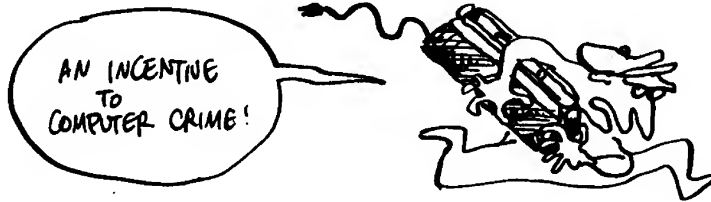


THERE ARE ALSO PROGRAMS TO CORRECT SPELLING — AND EVEN TO FIX SYNTAX AND GRAMMAR. SOON ILLITERATES WILL BE CREATING MASTERPIECES!

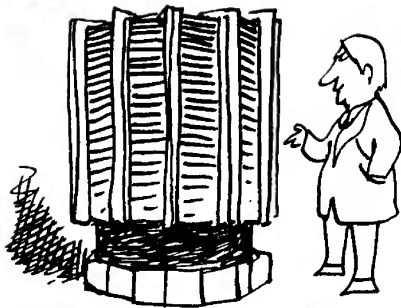
"HAMLET" BY J. FRED SHAKESPEARE...

A cartoon illustration of a man with a long, pointed beard and a simple tunic, sitting at a typewriter. He is looking at a speech bubble that says "HAMLET" BY J. FRED SHAKESPEARE... The scene is drawn in a simple, sketchy style.

A SMALL COMPUTER WITH WORD PROCESSING CAN BE QUITE INEXPENSIVE... THE CATCH IS THAT A "LETTER QUALITY" PRINTER CAN COST TEN TIMES THE PRICE OF A TYPEWRITER!



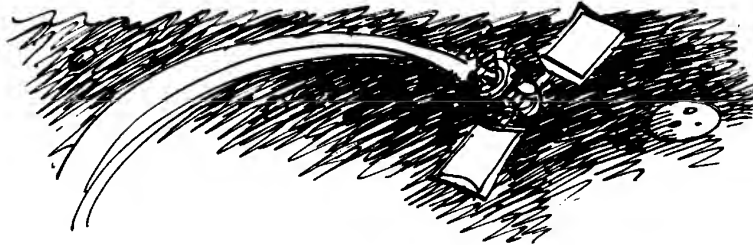
SCIENCE



CRAY-1 COMPUTER, CAPABLE OF 100 MILLION OPERATIONS PER SECOND !!

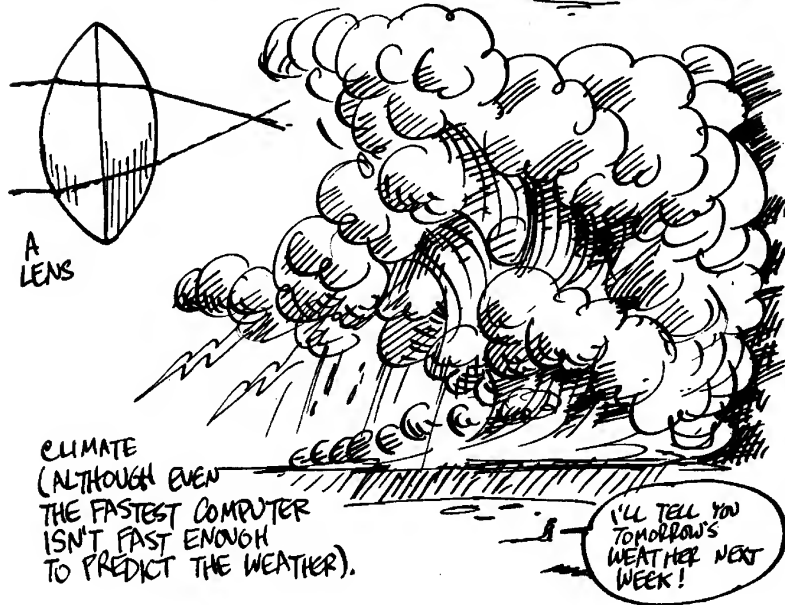
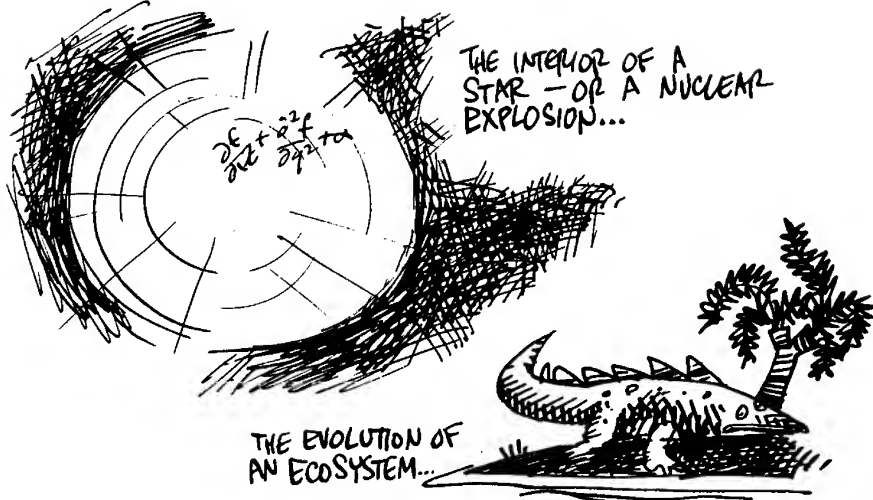
SCIENCE DEPENDS ON MATHEMATICS, AND COMPUTERS ARE SUPER MATH MACHINES. THE FASTEST, MOST POWERFUL COMPUTERS ARE MAINLY APPLIED TO SCIENTIFIC PROBLEMS.

THESE "SUPERCOMPUTERS" EXCEL AT SIMULATION. THE IDEA BEHIND SIMULATION IS TO FEED THE COMPUTER THE EQUATIONS GOVERNING A PHYSICAL SYSTEM AND THEN MATHEMATICALLY "MOVE" THE SYSTEM ACCORDING TO THOSE LAWS.



TIME SPACE TRAVEL: A COMPUTER CAN GUIDE A CRAFT TO THE MOON, BECAUSE IT CAN INTERNALLY SIMULATE THE ENTIRE FLIGHT !!

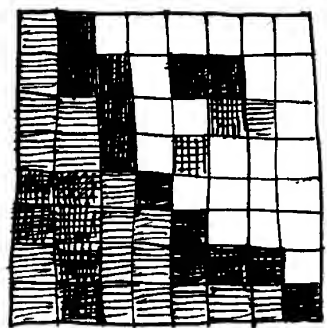
COMPUTERS CAN SIMULATE:



GRAPHICS



FROM THE SIMPLEST 'PONG' SCREEN TO THE MOST SOPHISTICATED FLIGHT SIMULATOR, THE IDEA IS THE SAME:



DIVIDE THE SCREEN AREA INTO A LARGE NUMBER OF TINY RECTANGLES ("PIXELS") AND ASSIGN EACH ONE A COLOR AND BRIGHTNESS.

THAT'S WHY COMPUTER PICTURES HAVE CORNERS!

BUT THERE ARE ALSO ALGORITHMS FOR SMOOTHING CORNERS!

UNFORTUNATELY, IT TAKES A LOT OF COMPUTER POWER TO DO FANCY GRAPHICS. SMALL COMPUTERS MOSTLY DO THINGS LIKE MAKE PIE CHARTS...

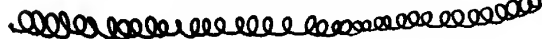


COMMUNICATION

THE BIGGEST COMPUTER SYSTEM OUTSIDE GOVERNMENT BELONGS TO THE TELEPHONE COMPANY.

A VOICE (OR ANY OTHER SIGNAL) CAN BE DIGITALLY ENCODED, TRANSMITTED, AND DECODED.

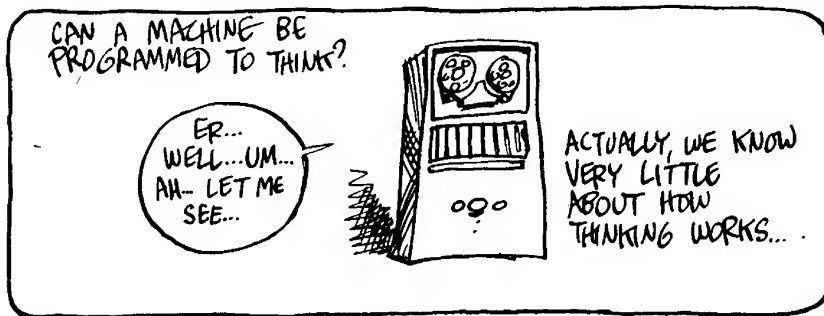
COMPUTERS ALSO CONTROL THE ROUTING AND SWITCHING OF CALLS THROUGH THE NETWORK — AND KEEP TRACK OF EVERYONE'S BILL!



COMPUTERS CAN BE PROGRAMMED TO RECOGNIZE PARTICULAR WORDS OR GROUPS OF WORDS — A CAPABILITY NOT LOST ON THE INTELLIGENCE COMMUNITY..

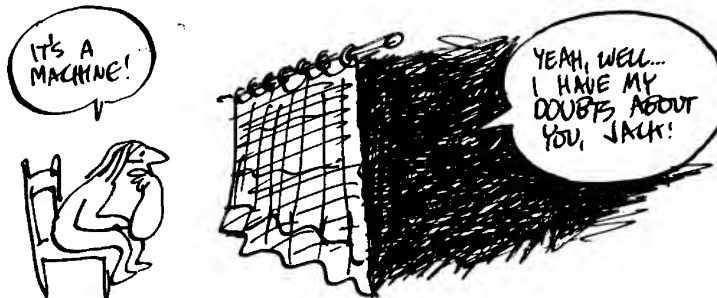
ARTIFICIAL INTELLIGENCE

DESPITE THEIR INCREDIBLE SPEED AND ACCURACY, COMPUTERS ARE LOUSY AT PATTERN RECOGNITION, ANALYSIS, HUNCH-PLAYING, AND UNDERSTANDING HUMAN LANGUAGE!

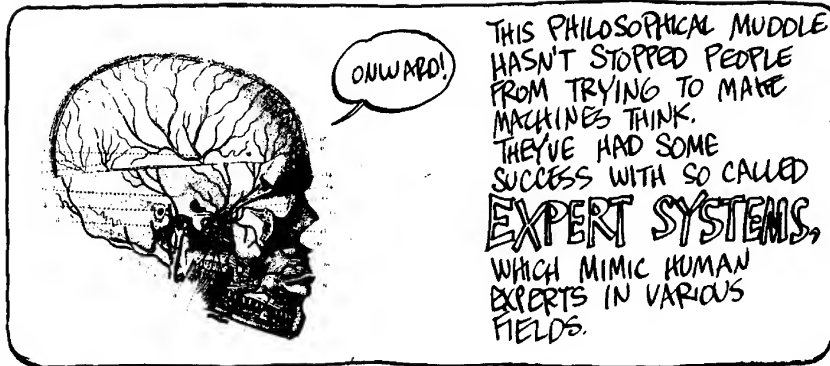


SO A BETTER QUESTION IS: HOW CAN YOU TELL IF A MACHINE IS THINKING?

ALAN TURING SUGGESTED THIS TEST: SUPPOSE YOU COULD COMMUNICATE WITH SOMETHING, OR SOMEONE, CONCEALED FROM VIEW. IF, ON THE BASIS OF THE CONVERSATION, YOU COULDN'T SAY WHETHER IT WAS MACHINE OR HUMAN, YOU WOULD HAVE TO SAY IT WAS THINKING!

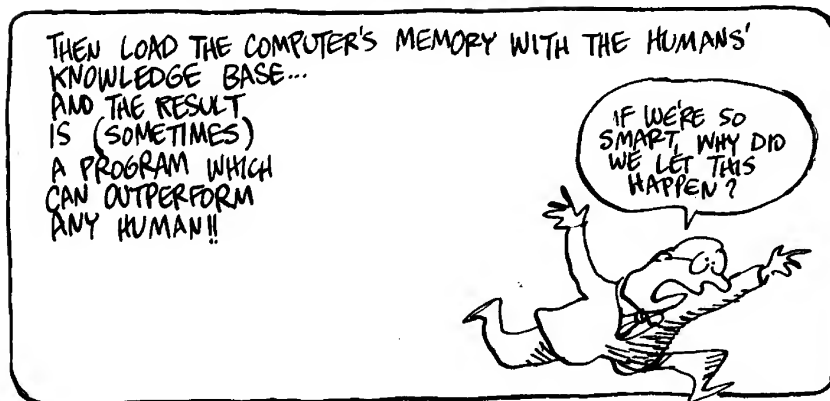
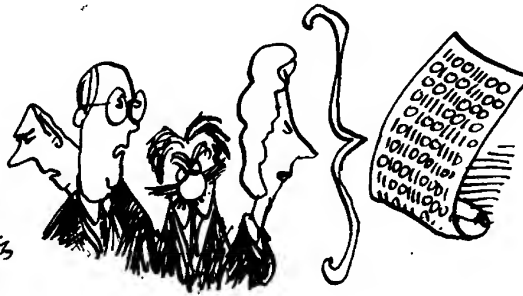


I PERSONALLY DISLIKE THIS CRITERION, ON THE GROUNDS THAT A SIMULATION ISN'T THE REAL THING...



HOW DO YOU CREATE AN EXPERT SYSTEM?

FIRST, INTERVIEW A BUNCH OF EXPERTS — GEOLOGISTS, FOR EXAMPLE — AND FORCE THEM TO SPELL OUT THE ALGORITHMS BEHIND THEIR SKILLS, HUNCHES AND BRAINSTORMS.



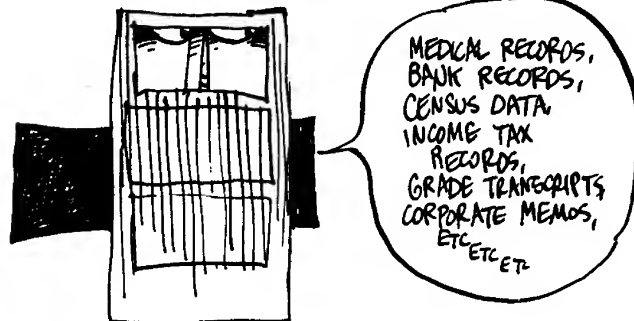
CRYPTOGRAPHY

SHHH!

THERE ARE STANDARD
CODES LIKE ASCII
(P. 128) FOR
CONVERTING WRITTEN
TEXT INTO BINARY...
BUT WHAT ABOUT USING
COMPUTERS FOR
SECRET
CODES??

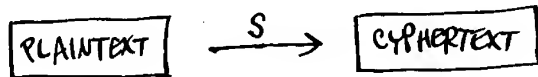


SECRET CODES USED TO BE STRICTLY MILITARY AND
SPY STUFF, BUT NOW MORE AND MORE SENSITIVE
INFORMATION IS STORED IN COMPUTER SYSTEMS:



SCRAMBLING DATA HAS BECOME AN IMPORTANT
WAY OF PROTECTING PRIVACY !!

ORDINARILY, INFORMATION IS STORED AS A BINARY STRING ANY COMPUTER CAN READ: THE **PLAINTEXT**, IN CRYPTOGRAPHIC JARGON. TO ENCRYPT IT YOU APPLY SOME ALGORITHM **S**, WHICH CONVERTS IT TO A SCRAMBLED MESSAGE CALLED THE **CYPHERTEXT**.



THEORETICALLY, IT'S IMPOSSIBLE TO RECONSTRUCT THE PLAINTEXT FROM THE CYPHERTEXT WITHOUT KNOWING SOMETHING ABOUT **S** ... HOWEVER, A POTENTIAL CODE-BREAKER COULD PUT A COMPUTER TO WORK SEARCHING FOR **S**.

➔ TO BE SECURE, **S** HAS TO BE SO COMPLICATED THAT EVEN THE FASTEST COMPUTER WOULD TAKE, SAY, A FEW MILLION YEARS TO FIGURE IT OUT!



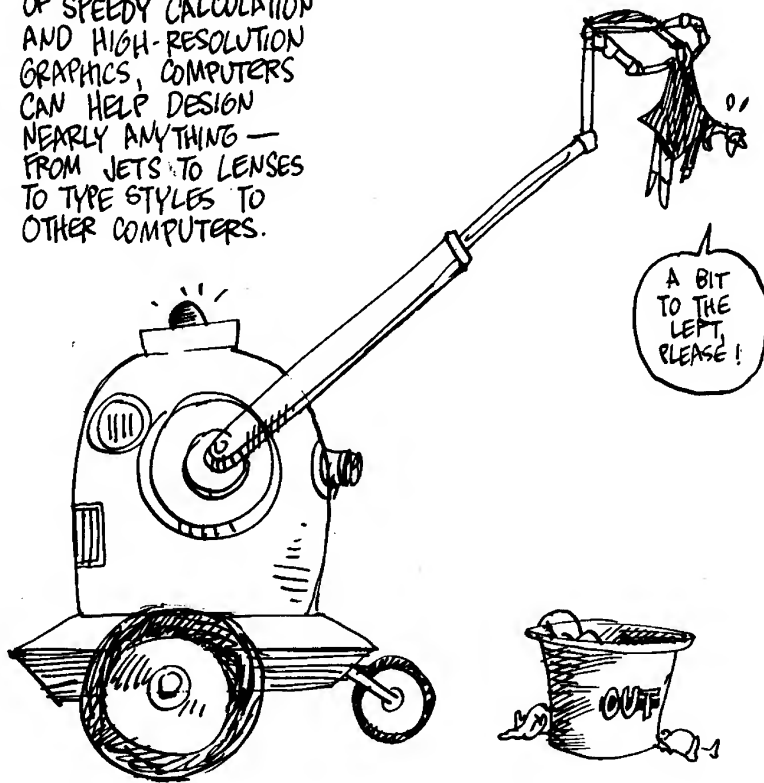
RECENTLY, THE NATIONAL BUREAU OF STANDARDS APPROVED A FAMILY OF ALGORITHMS AS A DATA ENCRYPTION STANDARD FOR THE NATION. SEVERAL SCIENTISTS SUSPECT THAT THIS STANDARD IS JUST COMPLEX ENOUGH TO STYMIE ORDINARY COMPUTERS, BUT NOT TOO TOUGH FOR THE NINE ACRES OF COMPUTERS OF THE NATIONAL SECURITY AGENCY!

A dark, dense forest scene with three speech bubbles containing the letters 'SH'.

CAD/CAM

- o COMPUTER-AIDED DESIGN /
- o COMPUTER-AIDED MANUFACTURE

USING A COMBINATION OF SPEEDY CALCULATION AND HIGH-RESOLUTION GRAPHICS, COMPUTERS CAN HELP DESIGN NEARLY ANYTHING — FROM JETS TO LENSES TO TYPE STYLES TO OTHER COMPUTERS.



THEN THEY CAN GO ON TO CONTROL AUTOMATIC MANUFACTURING PROCESSES AS WELL. YES, **ROBOTS** ARE ALREADY HERE!!

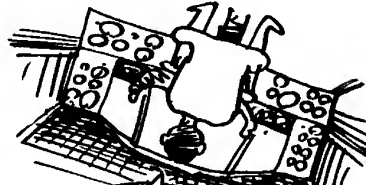
WAR

THE MILITARY CAN USE JUST ABOUT EVERY TYPE OF SOFTWARE WE'VE MENTIONED—AND THEN SOME!

ENIAC WAS BUILT FOR CALCULATING BALLISTICS... NOW WE HAVE BALLISTIC MISSILES!



FLIGHT SIMULATORS CAN TRAIN PILOTS RIGHT ON THE GROUND...



GREAT GRAPHICS ON THESE!

THEN THERE ARE THE FAMOUS "SMART" MISSILES, WHICH CAN FOLLOW A MOVING TARGET...



I'M ABOUT TO BLOW MYSELF UP... HOW SMART IS THAT?

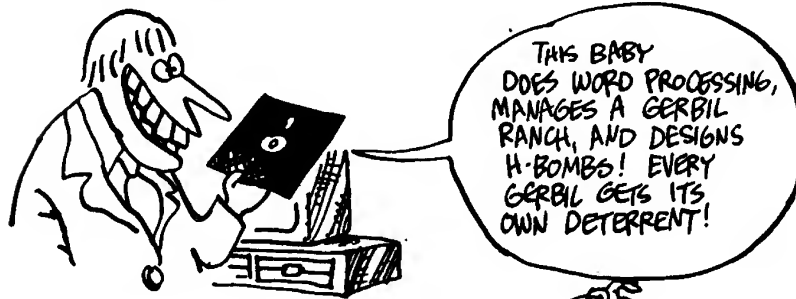
SUPERCOMPUTERS HELP DESIGN NUKES...



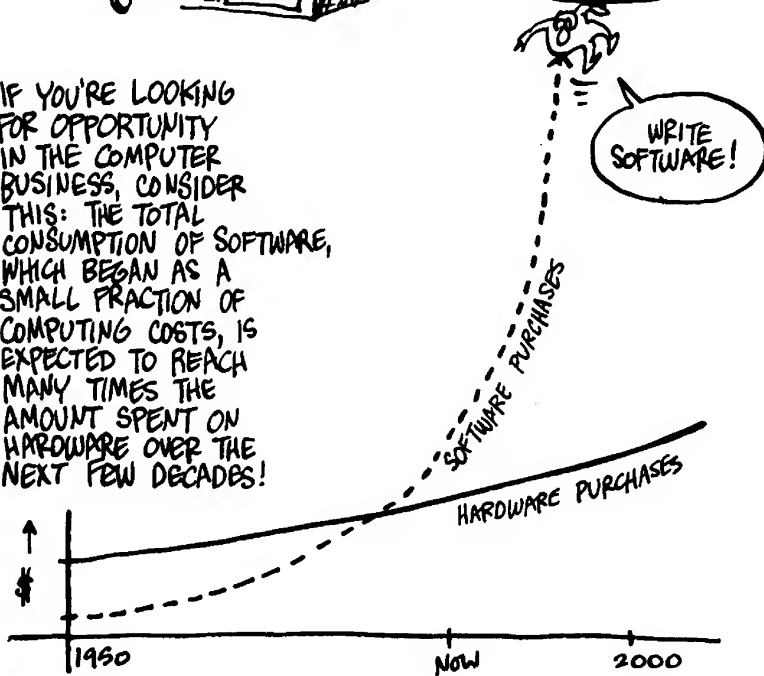
...NOT TO MENTION DATA PROCESSING AND CRYPTOGRAPHY... SO GREAT IS THE DEFENSE DEPARTMENT'S SOFTWARE NEED THAT THEY HAVE THEIR OWN PROGRAMMING LANGUAGE: ADA NAMED AFTER THE UNFORTUNATE LADY LOVELACE.



THIS LITTLE SURVEY ONLY BEGINS TO SUGGEST THE RANGE OF SOFTWARE CURRENTLY AVAILABLE. EVERY DAY THERE'S MORE... SOME PROGRAMS MOVE INTO NEW AREAS, WHILE OTHERS INTEGRATE EXISTING ROUTINES INTO NEW, MORE POWERFUL PACKAGES.



IF YOU'RE LOOKING FOR OPPORTUNITY IN THE COMPUTER BUSINESS, CONSIDER THIS: THE TOTAL CONSUMPTION OF SOFTWARE, WHICH BEGAN AS A SMALL FRACTION OF COMPUTING COSTS, IS EXPECTED TO REACH MANY TIMES THE AMOUNT SPENT ON HARDWARE OVER THE NEXT FEW DECADES!



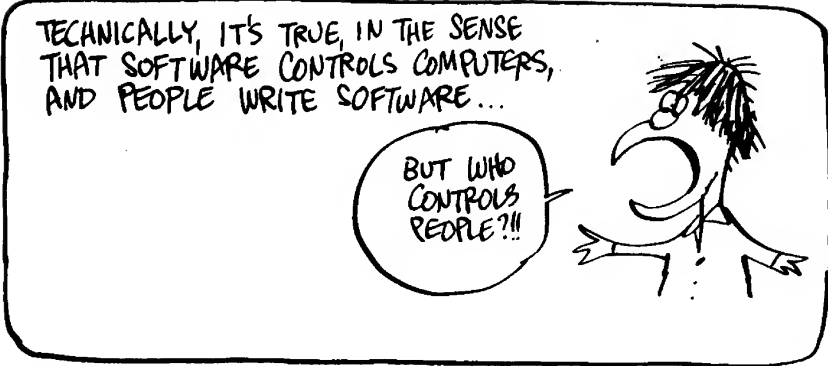
IN CONCLUSION,

A FEW WORDS
ABOUT THIS FAMILIAR
SENTENCE:

COMPUTERS
ONLY DO WHAT
PEOPLE TELL
THEM TO DO!



(WHICH IS WHAT
COMPUTER SCIENTISTS
SAY WHEN THEY WANT
TO BE REASSURING...)

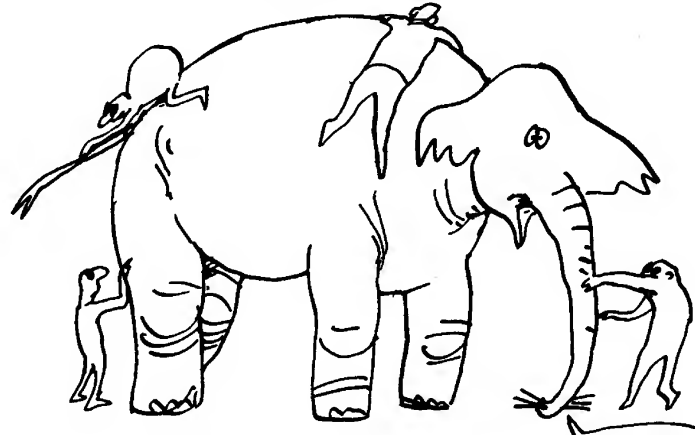


MY CALCULATOR
SAYS
 $2^{16} = 65,536.001$
(REALLY!)



ANOTHER PROBLEM IS THAT
ALGORITHMS DON'T ALWAYS
DO EXACTLY WHAT THEY
ARE SUPPOSED TO.

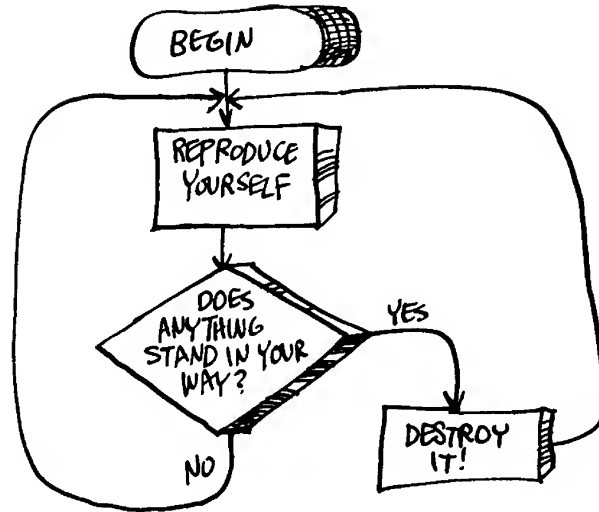
LARGE SOFTWARE SYSTEMS ARE WRITTEN BY TEAMS
OF PROGRAMMERS. LIKE THE ELEPHANT, NO ONE
UNDERSTANDS THE WHOLE THING!



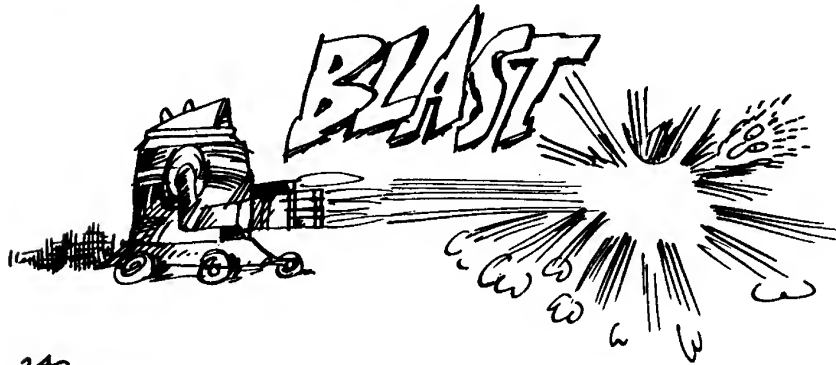
COMPUTERS ROUTINELY DO BIZARRE AND
UNEXPECTED THINGS, ESPECIALLY WHEN
RUNNING NEW, UNTESTED SOFTWARE!

I'M GETTING TO
BE A TIRED
METAPHOR!

FINALLY, CONSIDER THIS OMINOUS ALGORITHM:



WHILE NO COMPUTER IS INTELLIGENT, MOBILE, OR WELL EQUIPPED ENOUGH — YET — TO EXECUTE THESE INSTRUCTIONS, SUCH A MACHINE REMAINS A THEORETICAL POSSIBILITY. THIS PROGRAM WOULD MAKE IT SOMETHING VERY MUCH LIKE A COMPETING LIFE FORM!!!



AND IF YOU THINK THAT BECAUSE "IT'S ONLY A MACHINE," YOU CAN ALWAYS TURN IT OFF, PONDER THE WORDS OF NORBERT WIENER, A SCIENTIST WHO THOUGHT DEEPLY ABOUT THESE THINGS:

"TO TURN A MACHINE OFF EFFECTIVELY, WE MUST BE IN POSSESSION OF INFORMATION AS TO WHETHER THE DANGER POINT HAS COME. THE MERE FACT THAT WE HAVE MADE THE MACHINE DOES NOT GUARANTEE THAT WE SHALL HAVE THE PROPER INFORMATION TO DO THIS.... THE VERY SPEED OF... MODERN DIGITAL MACHINES STANDS IN THE WAY OF OUR ABILITY TO PERCEIVE AND THINK THROUGH THE THE INDICATIONS OF DANGER." *



* CYBERNETICS, SECOND EDITION, P. 176

SO WELCOME TO THE INFORMATION AGE, AND HAPPY COMPUTING !!



SOME FURTHER READING:



MEDIEVAL AND EARLY MODERN SCIENCE BY A. C. CROMBIE. TELLS HOW ISLAMIC SCIENCE CAME TO EUROPE.

THE MAKING OF THE MICRO BY C. EVANS. NICE DIAGRAMS OF OLD ADDING MACHINES

HISTORY OF MATHEMATICS BY A. GOTTLEMAN. DON'T MISS THE STORY OF NAPIER'S "PSYCHIC" CHICKEN!

THE COMPUTER FROM PASCAL TO VON NEUMANN BY H. GOLDSTINE. THE DEFINITIVE ACCOUNT OF ENIAC.

CHARLES BABBAGE, FATHER OF THE COMPUTER BY D. WALACE. AN EASY READ.

CHARLES BABBAGE AND HIS CALCULATING ENGINES, ED. BY P. & E. MORRISON. IN HIS OWN WORDS!



SYMBOLIC LOGIC AND THE GAME OF LOGIC BY LEWIS CARROLL. MILLIONS OF SILLY SYLLOGISMS

THE MATHEMATICAL THEORY OF COMMUNICATION BY C. SHANNON. CONTAINS TWO PARTS, ONE WITH AND ONE WITHOUT MATH

CYBERNETICS, 2ND EDITION, BY N. WEINER. THE THEORY OF AUTOMATIC CONTROL

UNDERSTANDING DIGITAL ELECTRONICS, BY D. McWHORTER. BOOLEAN CIRCUITS

UNDERSTANDING DIGITAL COMPUTERS, BY P. MIMS. A PERSONAL FAVORITE, BUT LOOK OUT FOR MISPRINTS!

INTRODUCTION TO MICROCOMPUTERS, BY A. OSBORNE (4 VOLUMES). VERY DETAILED

UNDERSTANDING COMPUTER SCIENCE BY R.S. WALKER. MORE ADVANCED TOPICS



ILLUSTRATING BASIC BY D. ALCOCK. A QUICK COURSE, USING QUASI-CARTOONS

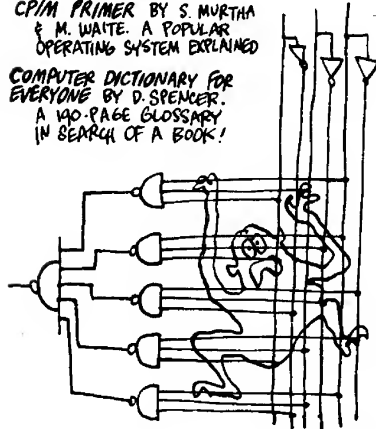
USING BASIC, BY R. DIDDAY & R. PAGE. A GENTLER, BUT WORDIER, APPROACH

PASCAL PRIMER BY D. FOX & M. WAITE. IT HELPS TO KNOW BASIC BEFORE READING THIS

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INDEX

- Abacus, 32-34, 43
- Abbreviations, mnemonic, 175
- Abstract symbol-manipulation, 42
- Accumulator, 173
- Ada programming language, 235
- Adder, 123
 - 1-bit, 125-126
- Adding machines, 59
- Address register, 173
- Addresses, 155
 - possible, 156
- Aiken, Howard, 72
- Algebra, 40
 - Boolean, 101-105
- Algorithm, 41-42
- Algorithms, 41, 195
 - examples of, 196
 - examples of non-algorithms, 197
 - flow of, 198
- Al-Khwarismi, 40
- Alphabet, 30-31
- Alphabetical order, 31
- Alphanumeric information, 130
- ALU (arithmetic logic unit), 130-132
- American Standard Code for Information Interchange (ASCII), 128
- Analysis, 200
- Analytical Engine, 53-55
- AND-gate, 107
 - multiple-input, 111
 - seatbelt buzzer in, 109
- AND logical operator, 103
- Applications software, 222
- Arabs, 40-43
- Arithmetic logic unit (ALU), 130-132
- Arithmetic on paper, 39
- Arithmetic table, Chinese, 29
- Artificial intelligence, 230-231
- ASCII (American Standard Code for Information Interchange), 128
- Assemblers, 205
- Assembly language, 174-176
- Assembly language statements, 175
- Asynchronous ripple counter, 148
- Automated type-setter, 54
- Automatic switches, 106-109
- B register, 173
- Babbage, Charles, 51-56, 58
- Babbage's Law, 58
- Ballistic tables, 74
- Ballistics, 73
- BASIC language, 162, 206
 - basic, 207-209
- Begin box, 198
- Binary calculation, 121
- Binary code, 205
- Binary coded decimal, 127
- Binary numbers, 115
 - adding two, 121
 - counting in, 120
 - multiplying, 122
 - subtracting, 122
 - translated into decimal numbers, 119
- Binary strings, 130
- Bits, 123
 - carry, 125
- Bone, tally, 20, 23
- Boole, George, 101
- Boolean algebra, 101-105
- Boxes, specially shaped, 198
- Branch instructions, 183
- Bubble memories, 166
- Bugs, 210, 211
- Bus architecture, 170
- Bytes, 123
- C register, 173
- CAD/CAM, 234

- Calculation, 34
 - binary, 121
- Calculators, 36
 - mechanical, 59
- Calculus, 34
- Cards, punched, *see* Punched cards
- Carroll, Lewis, 100
- Carry bit, 125
- Cash registers, 59
- Census tabulators, 60
- Charge-coupled devices, 166
- Chinese arithmetic table, 29
- Chinese number system, 27-29
- Chinese writing, 22
- Circuits
 - electrical, completing, 62
 - integrated, 84
- Clocks, 142-143
- Codes, secret, 232-233
- Combinational logic, 142
- Communication, 229
- Compilers, 205, 223
- Computer-aided design/computer-aided
 - manufacture, 234
- Computers, 5-6
 - clocks and, 142-143
 - control of, 237-241
 - Cray-1, 226
 - describing, 89
 - evolution of, 14
 - IBM Personal, 96
 - information and, 6
 - problems with, 237-241
 - research in, 153
 - size of, 152
- Conditional branch box, 198
- Conditional jumps, 57, 183
 - "smart," 215
- Control, 170
 - transfer of, 183
- Control bus, 170
- Control flow, 93-94
- Control unit, 92
- Core memories, 158
- Counters, 146-148
 - asynchronous ripple, 148
 - synchronous, 148
- Counting
 - binary and decimal, 120
 - digital, 19
- Cray-1 computer, 226
- Crusaders, 44
- Cryptography, 232-233
- Cycles, 142
- Cyphertext, 233
- Data base management, 224
- Data bases, 224
- Data processing, 63
- Decimal system, 115
 - binary numbers translated into, 119
- Decoders, 129
- Deductive logic, 99
- "Difference Engine, The," 51-52
- Digital counting, 19
- Disallowed input, 137
- Disks, magnetic and floppy, 166-167
- DNA, 12
- DNA-protein system, 13
- DNA technology, recombinant, 86
- Duns Scotus, 99
- EBCDIC, 128
- Eckert, J. Presper, 74
- Egyptian number system, 26
- Egyptians, 22
- Electric "mouse," 7
- Electrical circuits, completing, 62
- Electricity, 65
- Electromechanical memories, 154
- Electromechanical switches, 71
- Electronic memories, 154
- Electronic Numerical Integrator and
 - Calculator (ENIAC), 75-76
- Encoders, 129
- Encoding instructions, 79
- End box, 198
- ENIAC (Electronic Numerical Integrator
 - and Calculator), 75-76
- EPRM, 161
- Expert systems, 231
- External storage of information, 20
- Facts, 7
- Fetching instructions, 178
- Fibonacci sequence, 219
- Fields, 176
- Fives, counting by, 23
- Flight simulators, 235
- Flip-flop inputs, 137

- Flip-flops, 133-137
 - master-slave, 144
- Floating point representation, 127
- Floppy disks, 166-167
- Flow charts, 198-200
 - examples of, 199
 - for multiple plug-ins, 203
 - for roommate receipts, 202
- For-next commands, 218
- FORTTRAN, 206

- Gating network, 140
- Gene, 12
- Glitches, 143
- Go-to statement, 215
- Grammar, laws of, 18
- Graphics, 228
- Gravitation, theory of, 45-46
- Greek mathematicians, 33

- Handfuls, counting by, 24-25
- Hardware, 187
- Hertz (one cycle per second), 142
- Hexadecimal numerals, 157
- Higher-level programming languages, 205
- Hindus, 37-38
- Hollerith, Herman, 60, 64

- IBM, 64
- IBM Personal Computer, 96
- If-then statement, 215
- Incas, 22
- Increments, 146
- Inductive logic, 99
- Industrial Revolution, 49
- Information
 - ages of, 1-86
 - alphanumeric, 130
 - computers and, 6
 - defined, 7-8
 - excess, 3-5, 86
 - external storage of, 20
 - forms of, 8-9
 - power of, 12
 - stored, 10
- Information flow, 93-94
- Information processing, 11
 - understanding, 90
- Information theory, 7

- Input, 48, 92, 95
 - card-reading device, 54
 - disallowed, 137
 - flip-flop, 137
- Input box, 198
- Input-output (I/O) tables, 110, 112-113
- Input statement, 214
- Input wire, 106
- Instruction register, 173
- Instruction set, Motorola 6800, 182
- Instructions, 48
 - branch or jump, 183
 - 8-bit, 156
 - encoding, 79
 - fetching, 178
 - machine, 176
 - microinstructions, 178
 - to mill, 53
- Integers, 127
- Integrated circuits, 84
- Integration, large-scale and very
 - large-scale, 84
- Intelligence, artificial, 230-231
- Internal memory, 155
- Interpreters, 205
- Inverters, 108
- I/O (input-output) tables, 110, 112-113

- Jacquard, Joseph Marie, 50
- Japanese calculation of π , 29
- Jump, conditional, *see* Conditional jumps
- Jump instructions, 183

- K (kilo), 163

- Language
 - assembly, *see* Assembly language
 - entries
 - BASIC, *see* BASIC language
 - expressive, 17
 - higher-level programming, 205
 - machine, 177
- Large-scale integration (LSI), 84
- Latches, 138
 - gated, 140
- Leibniz, Gottfried Wilhelm, 47
- Let statement, 211
- Life form, 13
 - competing, 240
- Line numbers, 209

Loaders, 223
 Logic, 99
 combinational, 142
 laws of, 18
 sequential, 142
 simple, 150
 symbolic, 101
 Logic gates, multiple-input, III
 Logic unit, arithmetic (ALU), 130-132
 Logical operations, 98
 Logical operators, 103-104
 Logical spaghetti, 87-184
 Loom, Jacquard, 50
 Looping, value of, 57
 Lovelace, Ada, 56-58

 Machine instruction, 176
 Machine language, 177
 Magnetic disks, 166
 Magnetic tape, 165
 Mainframes, 85
 Mark I, 72
 Mass storage, 165
 uses of, 168
 Master-slave flip-flop, 144
 Mauchly, John, 74
 Mechanical calculators, 59
 Megaflops, 85
 Memory, 94, 95
 bubble, 166
 core, 158
 electromechanical, 154
 electronic, 154
 internal, 155
 random access (RAM), 159
 read-only, *see* Read-only memory
 Memory unit, 54
 Merge program, 81
 Messages, form of, 16
 Messenger RNA, 12
 "Method of the Celestial Element, The,"
 29
 Microcomputer, 85
 Microinstructions, 178
 Microprogram, 181
 Military software, 235
 Mill of the Analytical Engine, 53-55
 Minicomputer, 85
 Mnemonic abbreviations, 175

 Modem, 96
 Motorola 6800 instruction set, 182
 "Mouse," electric, 7
 Multiple-input logic gates, III
 Multiplication, binary, 122
 Music, 9

 NAND-gate, 134
 Napier, John, 47
 "Napier's bones," 47
 Newton, Isaac, 45
 Nibbles, 124
 NOR-gate, 138
 NOT logical operator, 104
 Nucleotide pairs, 12
 Number system
 Chinese, 27-29
 Egyptian, 26
 Numbers, 18
 binary, *see* Binary numbers
 hexadecimal, 157
 Numerical variables, 210-211

 Object code, 205
 Op-code, 176
 Operand, 175
 Operating system, 223
 Operations, logical, 98
 Operator, 175
 logical, 103-104
 Optic nerve, 8
 Optical disks, 166
 OR-gate, 108
 multiple-input, III
 OR logical operator, 103
 Order, alphabetical, 31
 Output, 54, 92, 95
 Output box, 198
 Output wire, 106

 Paper, 37
 arithmetic on, 39
 Paper tape, 165
 Papermaking, 43
 Parallel registers, 141
 Pascal, 206
 Pascal, Blaise, 47
 Personal Computer, IBM, 96
Pi, Japanese calculation of, 29

Pictograms, 30-31
 Pictures, 8
 Pixels, 228
 Plaintext, 233
 Powers of two, 118
 Print statement, 212
 Printer, 96
 Procedure box, 198
 Processing unit, 92, 95, 97
 Program counter, 173
 Programmable ROMs (PROMS), 161
 Programmer, first, 57
 Programming languages, higher-level, 205
 Programs, *see also* Software microprogram, 181
 for multiple plug-ins, 217
 for roommate receipts, 216
 self-modification by, 80
 sort and merge, 81
 stored, 78-80, 82
 Prompts, 208
 PROMS (programmable ROMs), 161
 Punched cards, 49-50, 165
 functions of, 55
 input device for, 54
 responses on, 61
 Punctuation, 209
 Pushbutton switches, 67

 Random access memory (RAM), 158-160
 Read-data statement, 210
 Read-only memory (ROM), 158-159, 161
 programmable (PROMS), 161
 uses of, 162
 Recombinant DNA technology, 86
 Recorded signals, 10
 Registers, 139
 parallel, 141
 shift, 144
 Relay
 telephone, 69
 Relay, automatic, 68
 Remarks, 209
 Renaissance, 44
 Return key, 208
 Rings on floppy disks, 167
 Ripple counter, asynchronous, 148
 RNA, messenger, 12

 Robots, 234
 ROM, *see* Read-only memory
 Romans, 34-35
 Rotary switches, 67
 Run statement, 208

 Schickard, Wilhelm, 47
 Scientific problems, 226-227
 Secret codes, 232-233
 Sectors on floppy disks, 167
 Self-modification, program, 80
 Self-reproducing machines, 193
 Semicolons, 213
 Semiconductors, 83
 Senses, 15
 Sensory impressions, 15
 Sequential logic, 142
 Shannon, Claude, 7
 Shift register, 145
 Signals, 8-9
 recorded, 10
 Simulations, 226-227
 Software, 185-236
 applications, 222
 defined, 187
 growth of, 236
 military, 235
 range of, 236
 survey of, 221-236
 systems, 222
 Sort and merge program, 81
 Source code, 205
 Spaghetti, logical, 87-184
 Statements, 209
 Storage
 external, of information, 20
 mass, *see* Mass storage
 Store, memory, 54
 Stored information, 10
 Stored programs, 78-80, 82
 Subroutines, 57
 Subtraction, binary, 122
 Sumerians, 21-22
 Supercomputers, 85
 Superminicomputers, 86
 Swan-pan, Chinese arithmetic table, 29
 Switchboard, 68
 Switches, 66-68
 automatic, 106-109

Switches (*cont'd*)
 electromechanical, 71
 patterns of, 70
Symbol-manipulation, abstract, 42
Symbolic logic, 100
Synchronous counters, 148
Systems, 222
 expert, 231
 operating, 223
Systems software, 222

Tabulators, census, 60
Tally bone, 20, 23
Tape, paper and magnetic, 165
Tartaglia, Niccolo, 45
Telephone company, 229
Telephone relay, 69
Telephones, 68
Ten, 116-117
Tens, counting by, 23
Three-body problem, 46
Timing, 142
Toggle switches, 67
Toggling, 147
Transfer of control, 183
Transistors, 83-84
Transition rules, 191
Truth-values, 102

Tube, vacuum, 69
Turing, Alan, 190, 230
Turing machines, 191-192
Two, powers of, 118
"Two's complement" method, 122
Type-setter, automated, 54

Unconditional branching statement, 215
Universal Turing machine, 192

Vacuum tube, 69
Variables, 210
 numerical, 210-211
Very large-scale integration (VLSI), 84
von Neumann, John, 77, 193

Wiener, Norbert, 241
Word processing, 225
Words, 8
World War II, 72
Writing, 21-22
 Chinese, 22
Written zero, 37-38

Zero, 27-28
 written, 37-38
Zuse, Konrad, 71